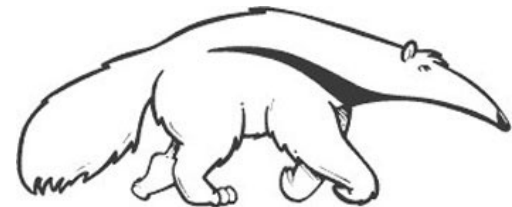


# Algorithms for Causal Probabilistic Graphical Models

Class 3:  
**Search**

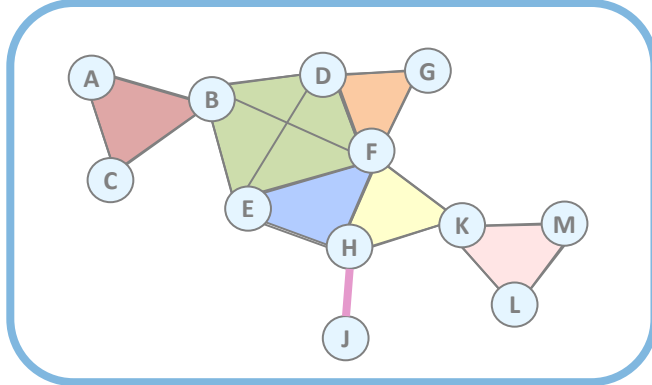
Athens Summer School on AI  
July 2024

Prof. Rina Dechter  
Prof. Alexander Ihler

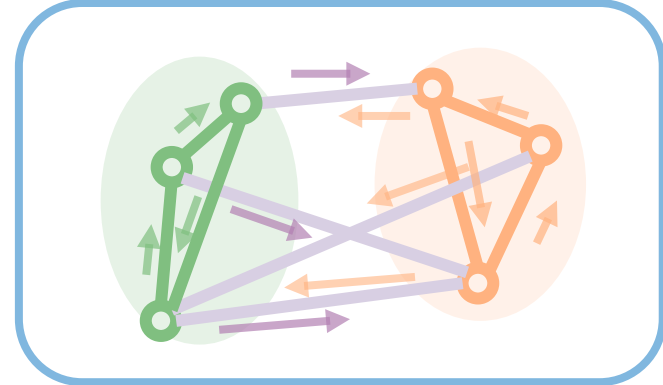


# Outline of Lectures

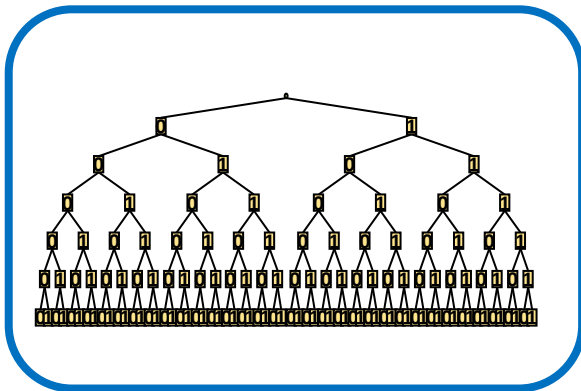
Class 1: Introduction & Inference



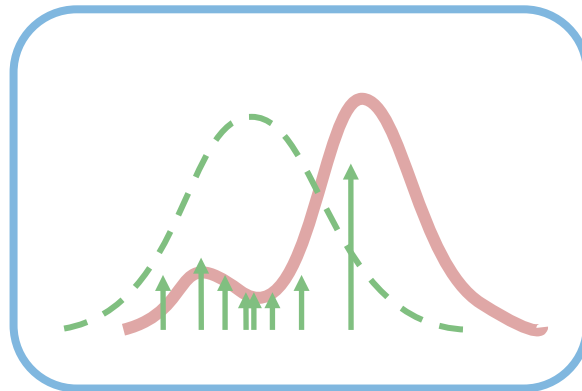
Class 2: Bounds & Variational Methods



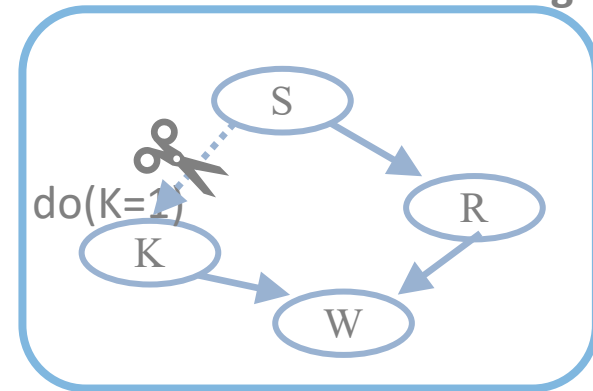
Class 3: Search Methods



Class 4: Monte Carlo Methods



Class 5: Causal Reasoning



# Graphical Models

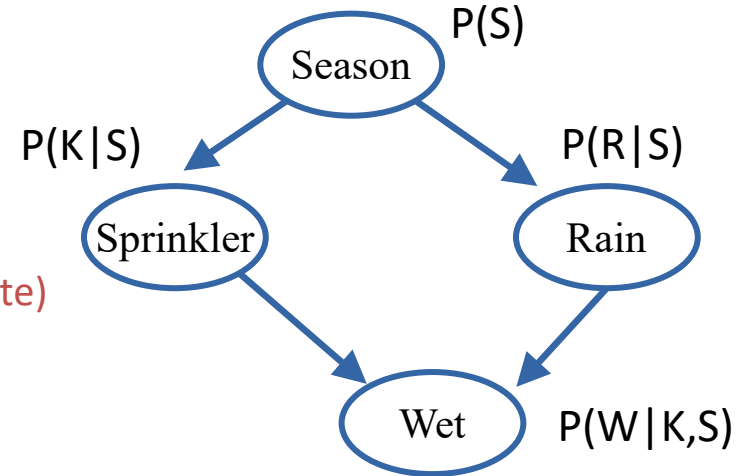
A *graphical model* consists of:

$X = \{X_1, \dots, X_n\}$  -- variables

$D = \{D_1, \dots, D_n\}$  -- domains (we'll assume discrete)

$F = \{f_{\alpha_1}, \dots, f_{\alpha_m}\}$  -- functions or CPTs

and a *combination operator*



The *combination operator* defines an overall function from the individual factors,

$$\text{e.g., "+" : } P(S, K, R, W) = P(S) \cdot P(K|S) \cdot P(R|S) \cdot P(W|K, S)$$

Notation:

Discrete  $X_i$  values called "states"

"Tuple" or "configuration": states taken by a set of variables

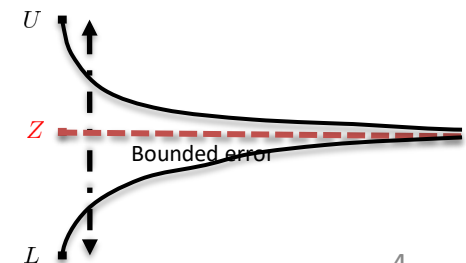
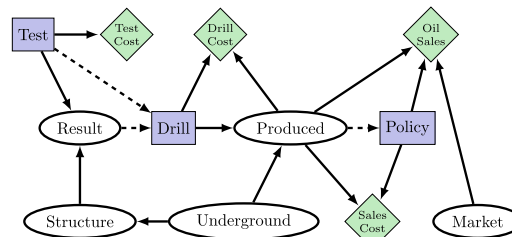
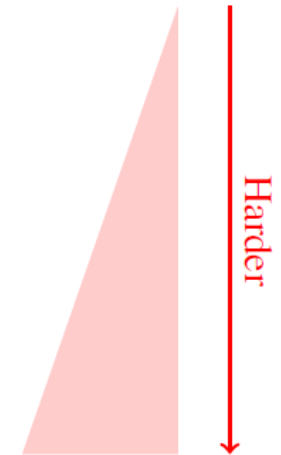
"Scope" of  $f$ : set of variables that are arguments to a factor  $f$

often index factors by their scope, e.g.,  $f_{\alpha}(X_{\alpha})$ ,  $X_{\alpha} \subseteq X$

# Probabilistic Reasoning Problems

- Exact inference time, space exponential in induced width
- Use **search** to trade memory for time and time for anytime bounds.

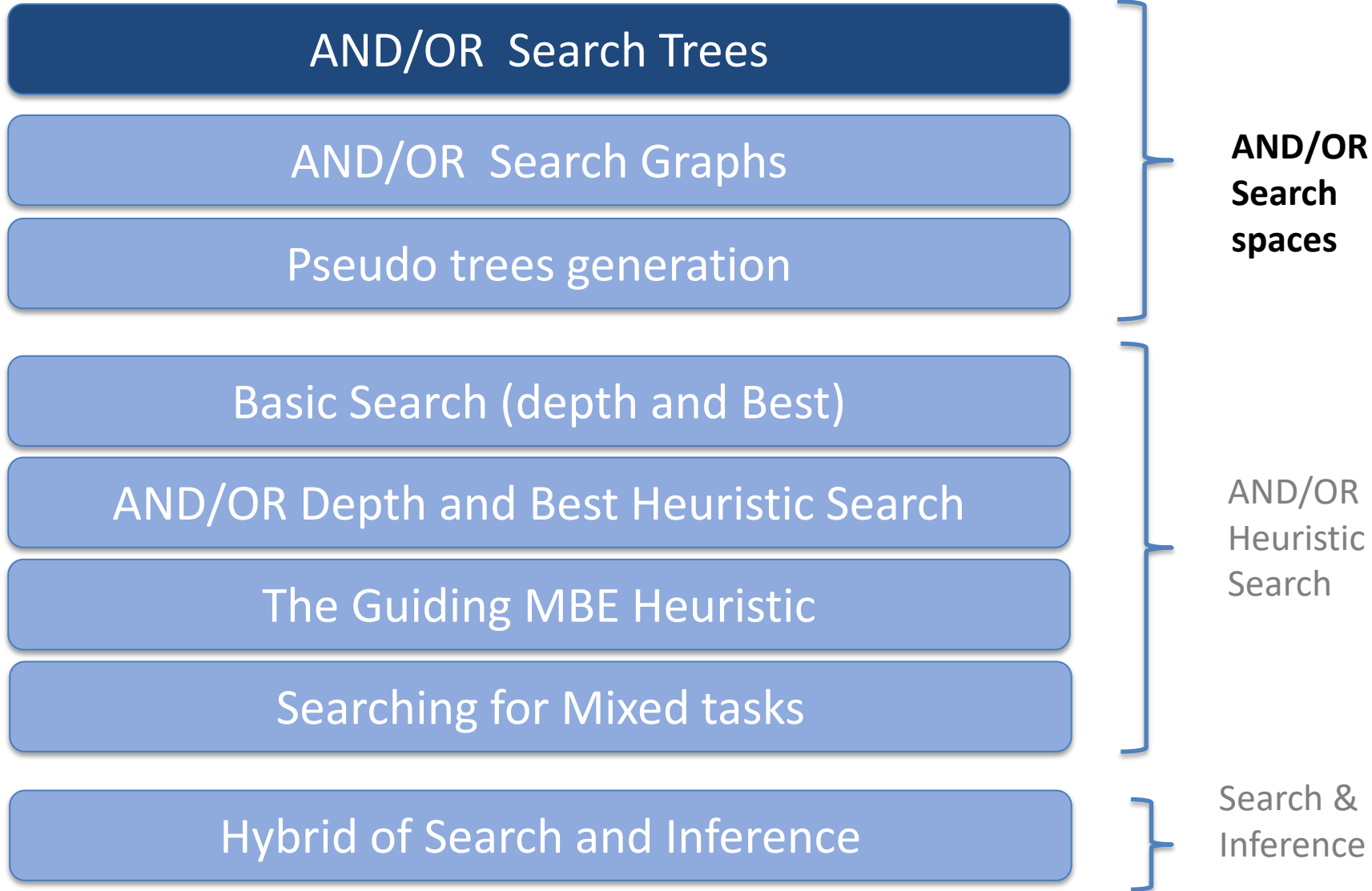
Max-Inference:	$f(x^*) = \max_x \prod_{\alpha} f_{\alpha}(x_{\alpha})$
Sum-Inference: (e.g., causal effects)	$Z = \sum_x \prod_{\alpha} f_{\alpha}(x_{\alpha})$
Mixed-Inference (MMAP):	$f_M(x_M^*) = \max_{x_M} \sum_{x_S} \prod_{\alpha} f_{\alpha}(x_{\alpha})$
Mixed-Inference (MEU): (e.g., decisions, planning)	$\text{MEU} = \max_{D_1, \dots, D_m} \sum_{X_1, \dots, X_n} \left( \prod_{P_i \in P} P_i \right) \times \left( \sum_{r_i \in R} r_i \right)$



# Outline: Search



# Outline: Search



# The Probability Tree

$$P(a, e = 0) = P(a) \sum_b P(b | a) \sum_c P(c | a) \sum_b P(d | a, b) \sum_{e=0} P(e | b, c)$$

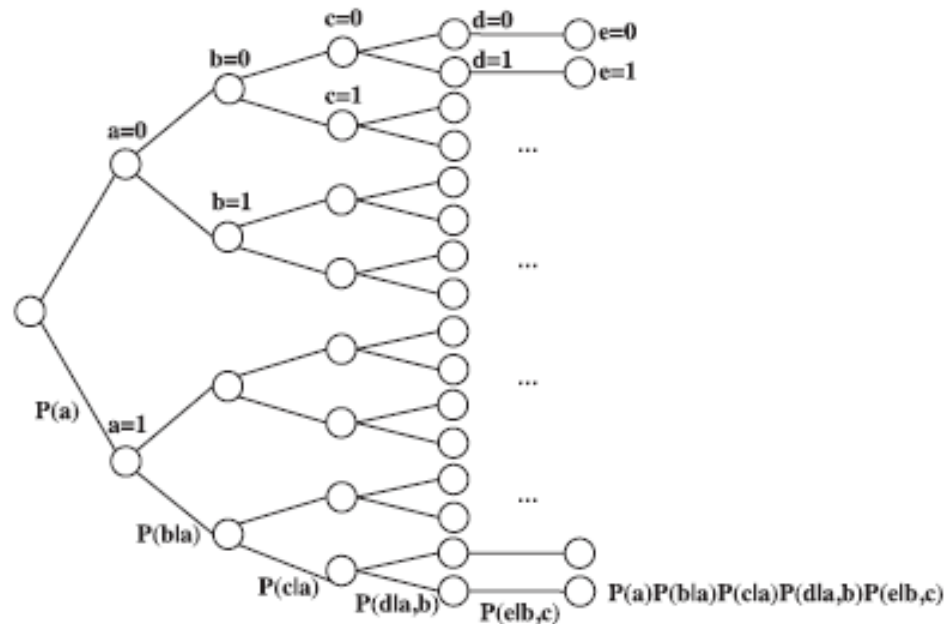
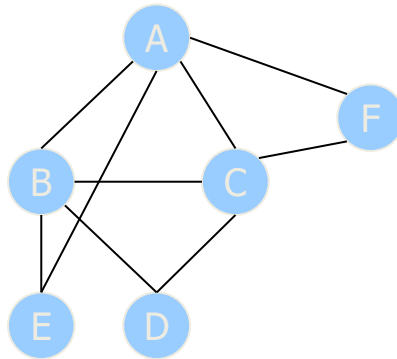


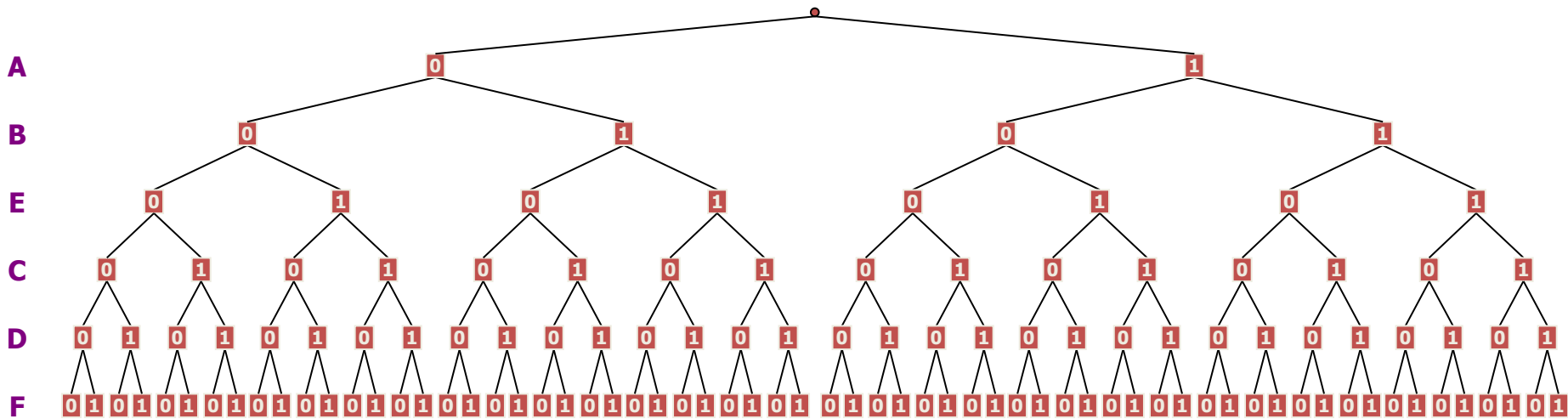
Figure 6.1: Probability tree for computing  $P(d=1, g=0)$ .

Complexity of conditioning: exponential time, linear space

# The Classic OR Search Space

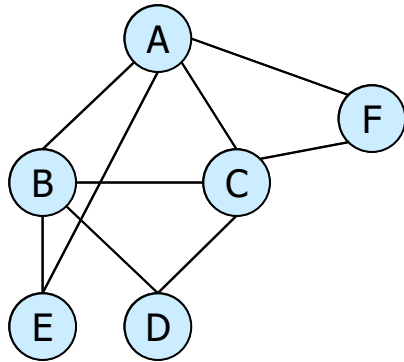


Ordering: A B E C D F

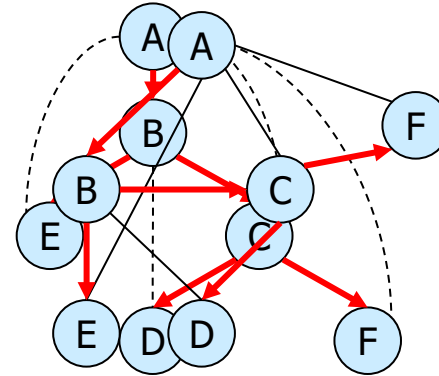




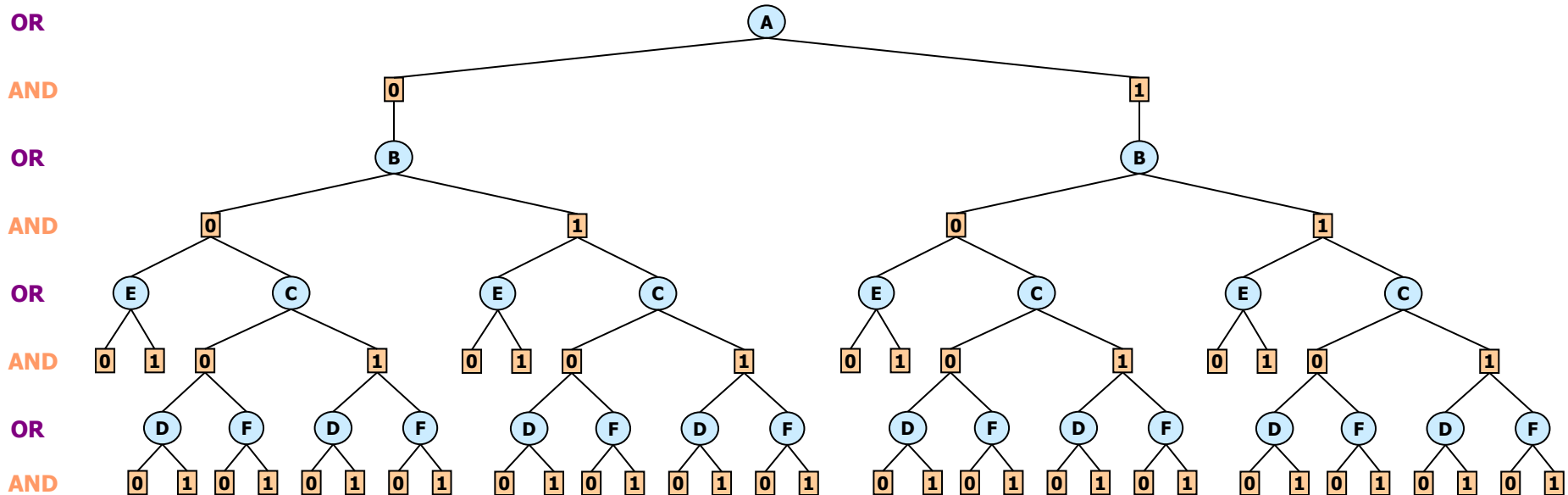
# AND/OR Search Space



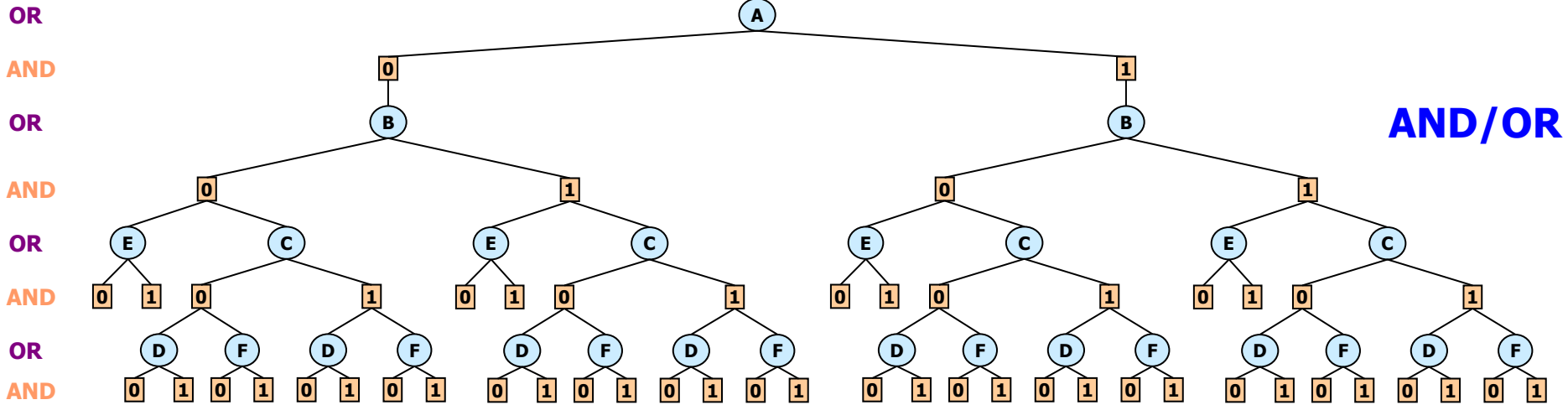
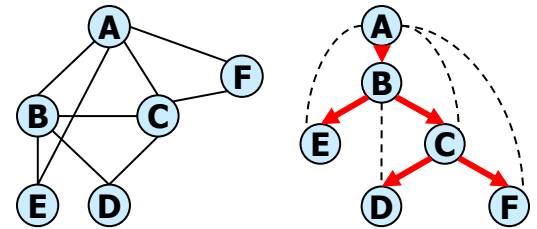
Primal graph



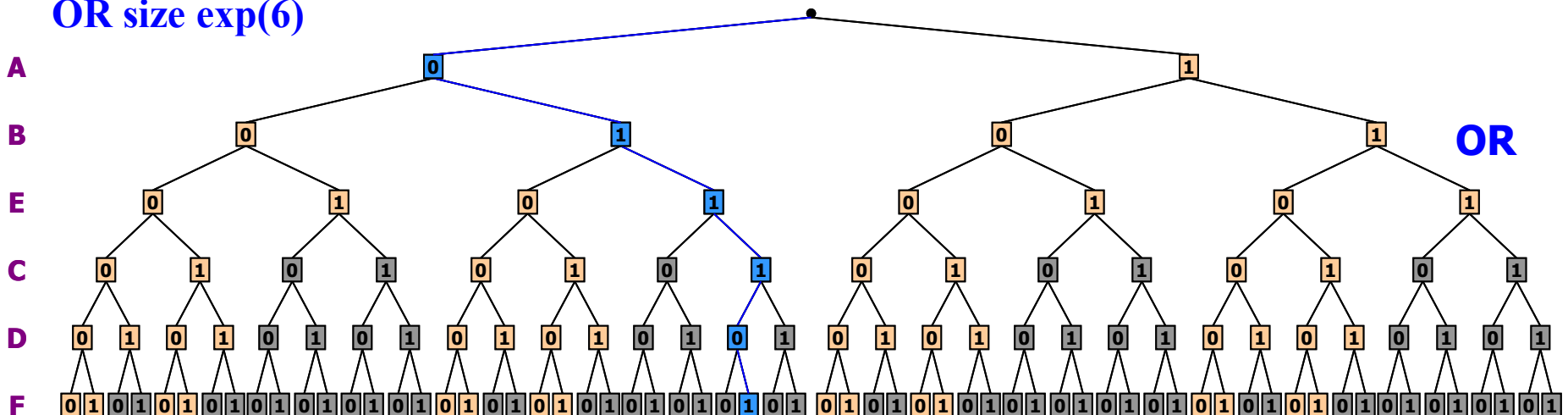
DFS tree



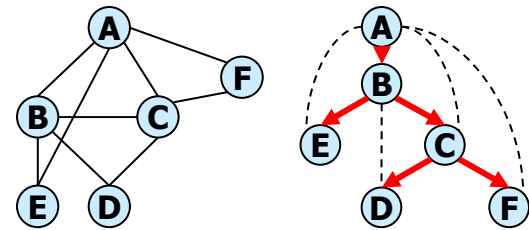
# AND/OR vs. OR



AND/OR size:  $\exp(4)$ ,  
OR size  $\exp(6)$



# AND/OR vs. OR



OR

AND

OR

AND

OR

AND

OR

AND

A

B

E

C

D

F

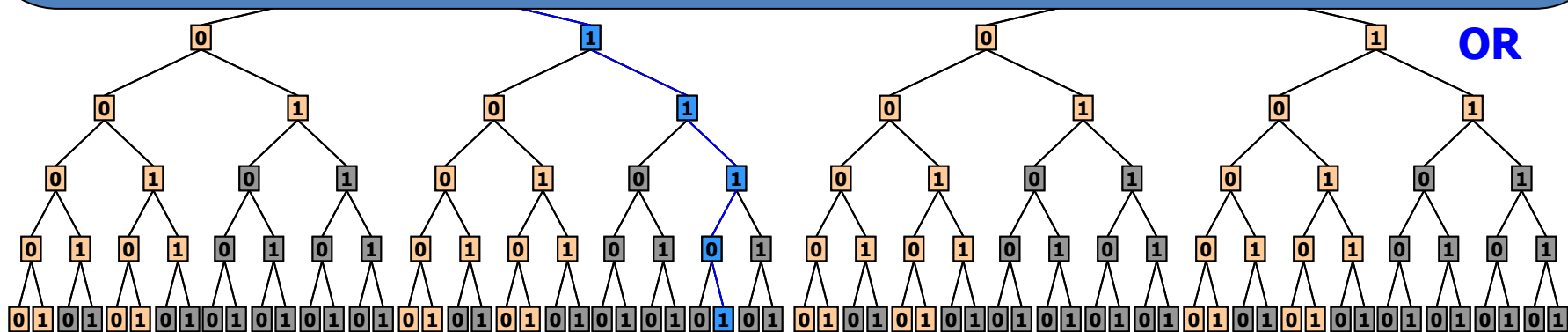
Dechter & Ihler


ESSA12024

AND/OR

- *Size of tree  $O(nk^h)$*
- *Can be traversed in*
- *Time  $O(nk^h)$ , Space  $O(n)$*
- *All solution trees = all configurations*

OR





Arc weights  
Cost of a solution tree  
The value function

# Arc Weights for AND/OR Trees

$P(E | A, B)$

A	B	E=0	E=1
0	0	.4	.6
0	1	.5	.5
1	0	.7	.3
1	1	.2	.8

Evidence: E=0

$P(B | A)$

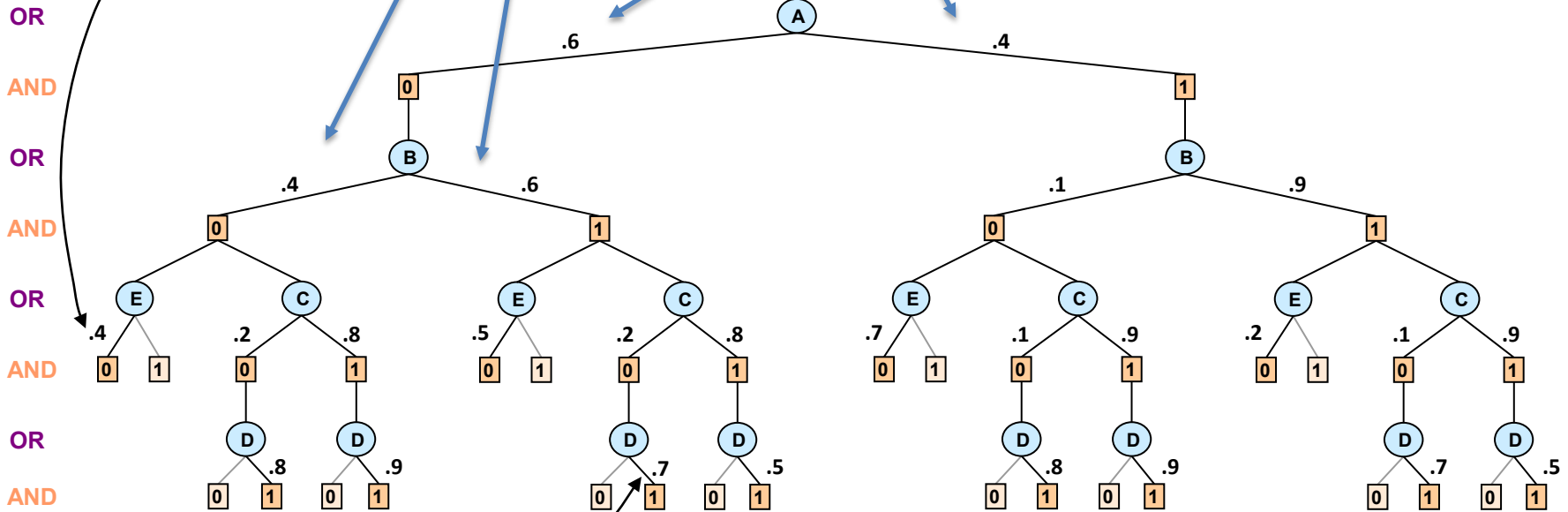
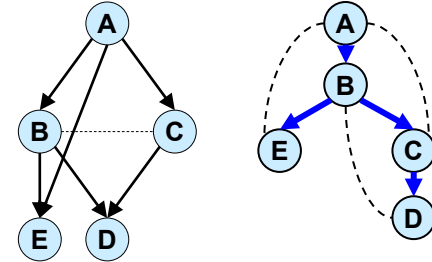
A	B=0	B=1
0	.4	.6
1	.1	.9

$P(C | A)$

A	C=0	C=1
0	.2	.8
1	.7	.3

$P(A)$

A	P(A)
0	.6
1	.4



$P(D | B, C)$

B	C	D=0	D=1
0	0	.2	.8
0	1	.1	.9
1	0	.3	.7
1	1	.5	.5

Evidence: D=1

OR to AND arc weight  $\langle X, x \rangle$  is the product of factors that all their arguments are just assigned at AND node  $X=x$  but not before

# Cost of a Solution Tree

$P(E | A, B)$

A	B	E=0	E=1
0	0	.4	.6
0	1	.5	.5
1	0	.7	.3
1	1	.2	.8

Evidence: E=0

$P(B | A)$

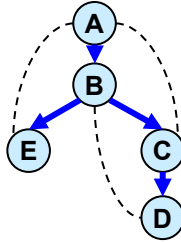
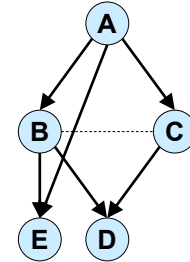
A	B=0	B=1
0	.4	.6
1	.1	.9

$P(C | A)$

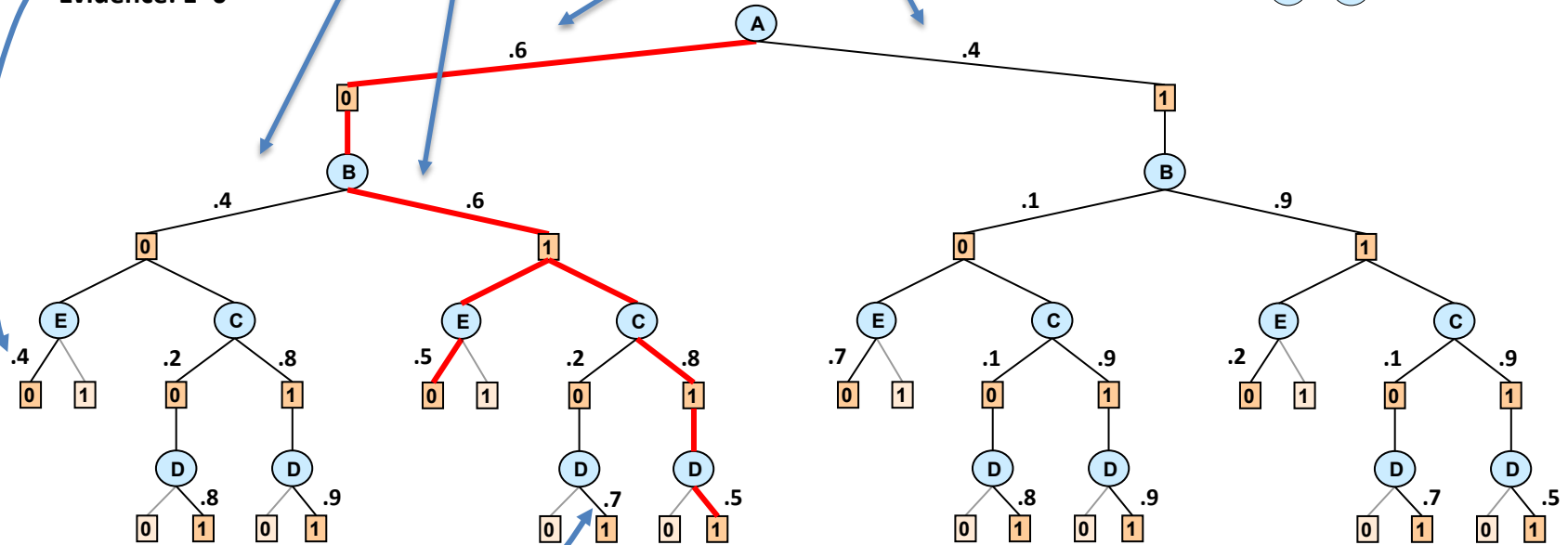
A	C=0	C=1
0	.2	.8
1	.7	.3

$P(A)$

A	P(A)
0	.6
1	.4



OR  
AND  
OR  
AND  
OR  
AND  
OR  
AND



$P(D | B, C)$

B	C	D=0	D=1
0	0	.2	.8
0	1	.1	.9
1	0	.3	.7
1	1	.5	.5

Evidence: D=1

A solution tree includes the root and has a single child for any OR node, and all children of any of its AND nodes

Cost of the solution tree: the product of weights on its arcs

Cost of  $(A=0, B=1, C=1, D=1, E=0) = 0.6 \cdot 0.6 \cdot 0.5 \cdot 0.8 \cdot 0.5 = 0.0720$

# The Value Function for (Probability of Evidence)

$$P(E | A, B)$$

A	B	E=0	E=1
0	0	.4	.6
0	1	.5	.5
1	0	.7	.3
1	1	.2	.8

Evidence: E=0

$$P(B | A)$$

A	B=0	B=1
0	.4	.6
1	.1	.9

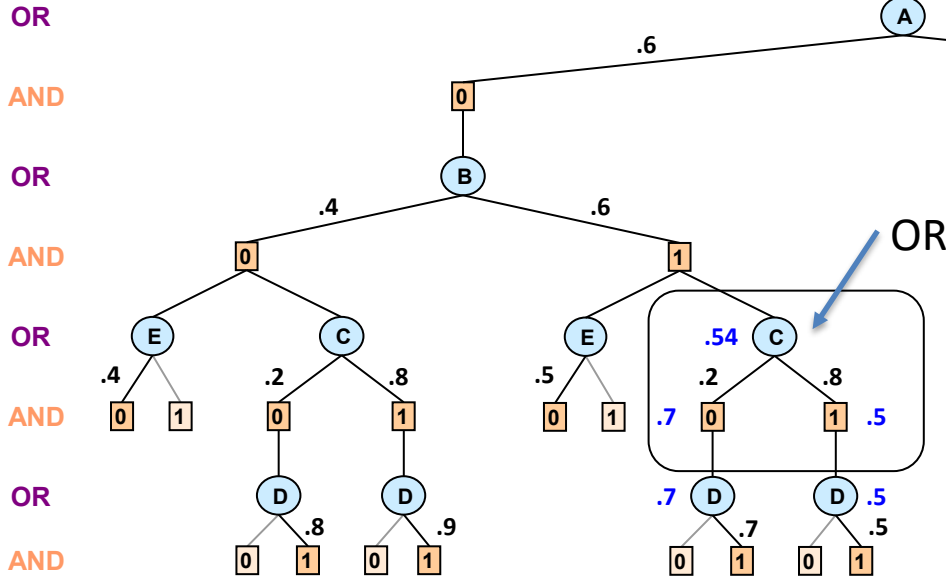
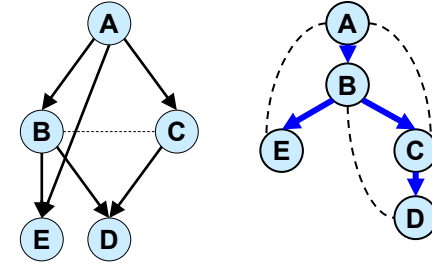
$$P(C | A)$$

A	C=0	C=1
0	.2	.8
1	.7	.3

$$P(A)$$

A	P(A)
0	.6
1	.4

$P(D=1, E=0)=?$



$$P(D | B, C)$$

B	C	D=0	D=1
0	0	.2	.8
0	1	.1	.9
1	0	.3	.7
1	1	.5	.5

Evidence: D=1

Value of node = updated belief for sub-problem below

AND node: product

$$\prod_{n' \in \text{children}(n)} v(n')$$

OR node: Marginalization by summation

$$\sum_{n' \in \text{children}(n)} w(n, n') v(n')$$

# The Value Function (Probability of Evidence)

$P(E | A, B)$

A	B	E=0	E=1
0	0	.4	.6
0	1	.5	.5
1	0	.7	.3
1	1	.2	.8

Evidence: E=0

$P(B | A)$

A	B=0	B=1
0	.4	.6
1	.1	.9

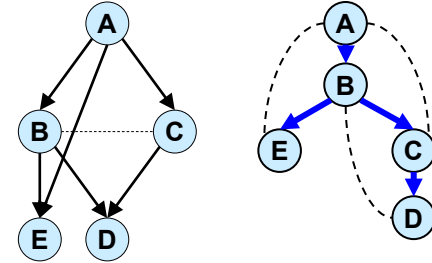
$P(C | A)$

A	C=0	C=1
0	.2	.8
1	.7	.3

$P(A)$

A	P(A)
0	.6
1	.4

$P(D=1, E=0) = ?$



OR

AND

OR

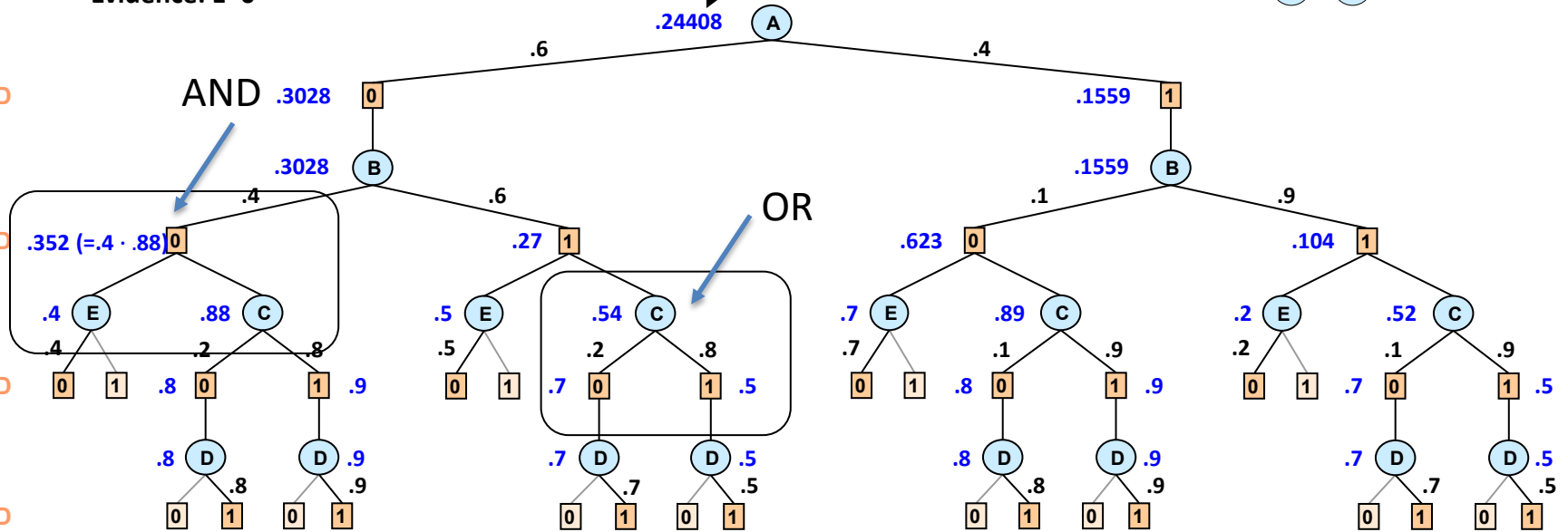
AND

OR

AND

OR

AND



$P(D | B, C)$

B	C	D=0	D=1
0	0	.2	.8
0	1	.1	.9
1	0	.3	.7
1	1	.5	.5

Evidence: D=1

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AND node: product

$$\prod_{n' \in \text{children}(n)} v(n')$$

OR node: Marginalization by summation

$$\sum_{n' \in \text{children}(n)} w(n, n') v(n')$$



# The Value Function

$$P(E | A, B)$$

A	B	E=0	E=1
0	0	.4	.6
0	1	.5	.5
1	0	.7	.3
1	1	.2	.8

Evidence: E=0

$$P(B | A)$$

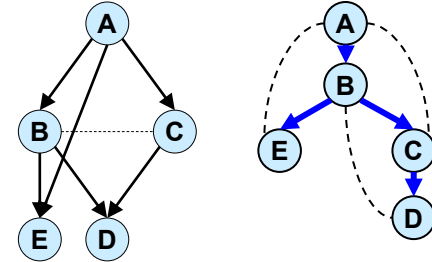
A	B=0	B=1
0	.4	.6
1	.1	.9

$$P(C | A)$$

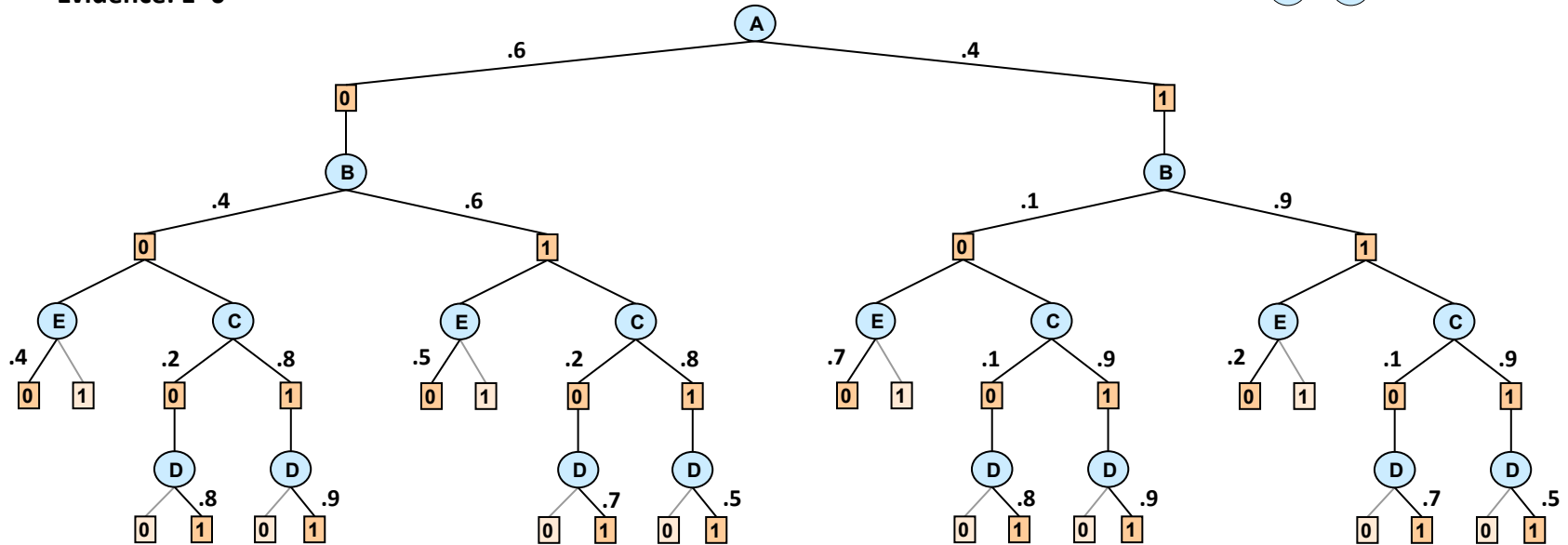
A	C=0	C=1
0	.2	.8
1	.7	.3

$$P(A)$$

A	P(A)
0	.6
1	.4



OR  
AND  
OR  
AND  
OR  
AND  
OR  
AND



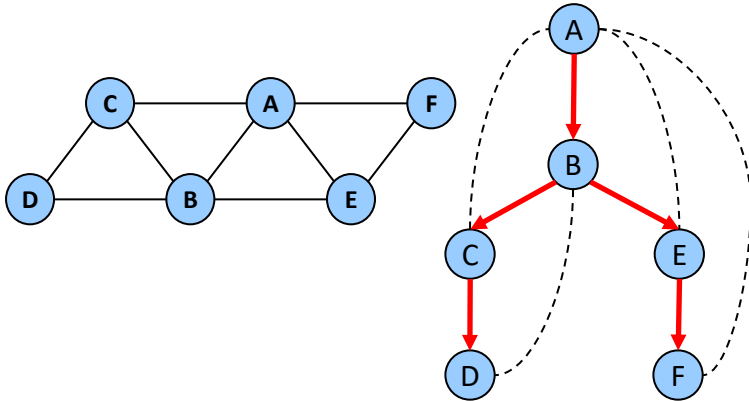
$$P(D | B, C)$$

B	C	D=0	D=1
0	0	.2	.8
0	1	.1	.9
1	0	.3	.7
1	1	.5	.5

Evidence: D=1

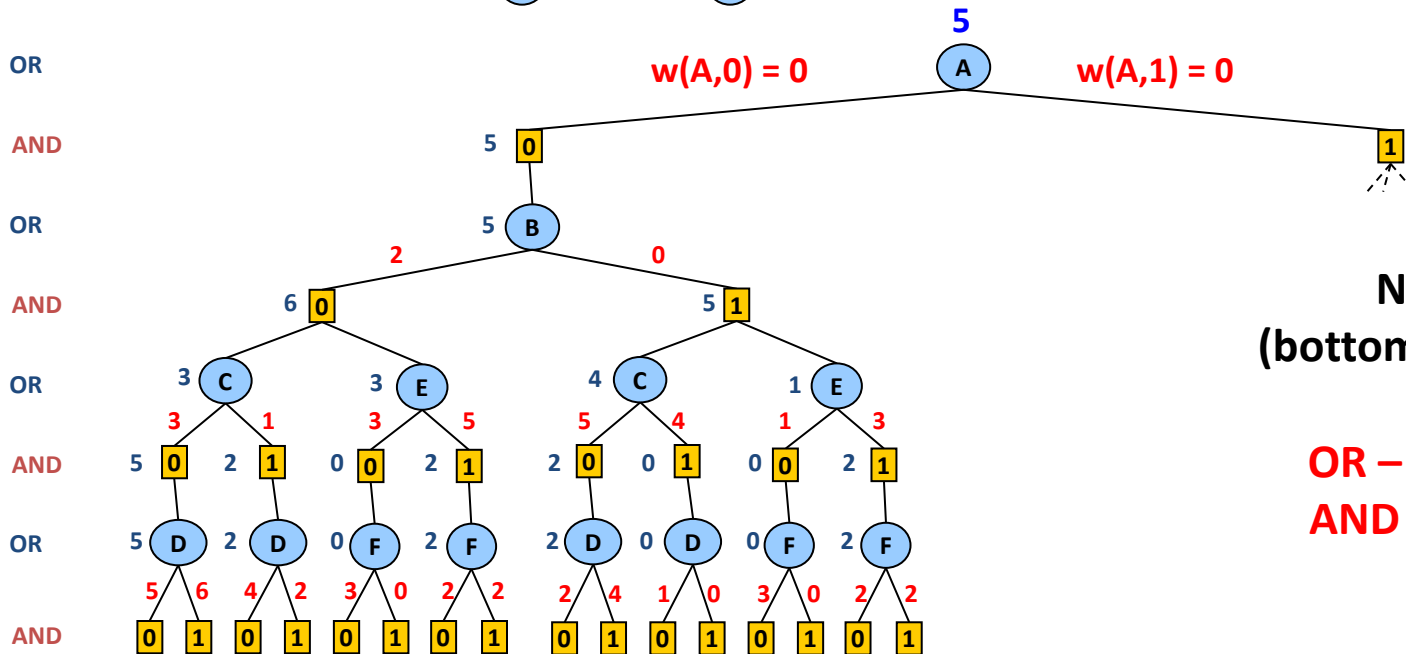
- $V(n)$  is dictated by the query of interest
- $V(n)$  the value of the sub-problem represented by  $T(n)$
- For sum-inference it is the probability mass below  $n$
- Can be computed recursively based on child values.

# The Value Function for Optimization



A	B	f <sub>1</sub>	A	C	f <sub>2</sub>	A	E	f <sub>3</sub>	A	F	f <sub>4</sub>	B	C	f <sub>5</sub>	B	D	f <sub>6</sub>	B	E	f <sub>7</sub>	C	D	f <sub>8</sub>	E	F	f <sub>9</sub>
0	0	2	0	0	3	0	0	0	0	0	2	0	0	0	0	0	4	0	0	3	0	0	1	0	0	1
0	1	0	0	1	0	0	1	3	0	1	0	0	1	1	0	1	2	0	1	2	0	1	4	0	1	0
1	0	1	1	0	0	1	0	2	1	0	0	1	0	2	1	0	1	1	0	1	1	0	0	1	0	0
1	1	4	1	1	1	1	1	0	1	1	2	1	1	4	1	1	0	1	1	0	1	1	0	1	1	2

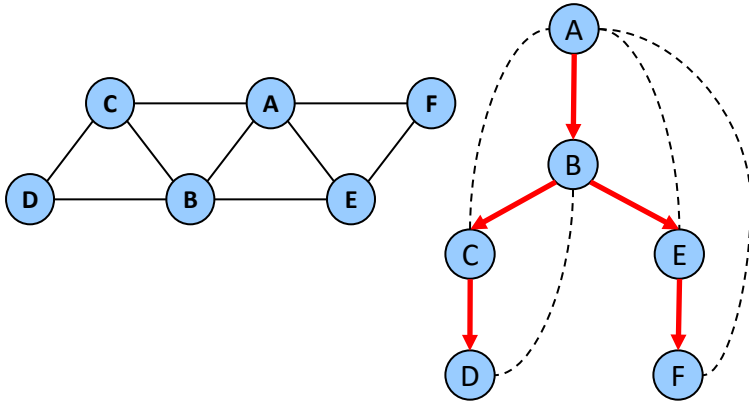
Objective function:  $F^* = \min_x \sum_{\alpha} f_{\alpha}(x_{\alpha})$



Node Value  
(bottom-up evaluation)

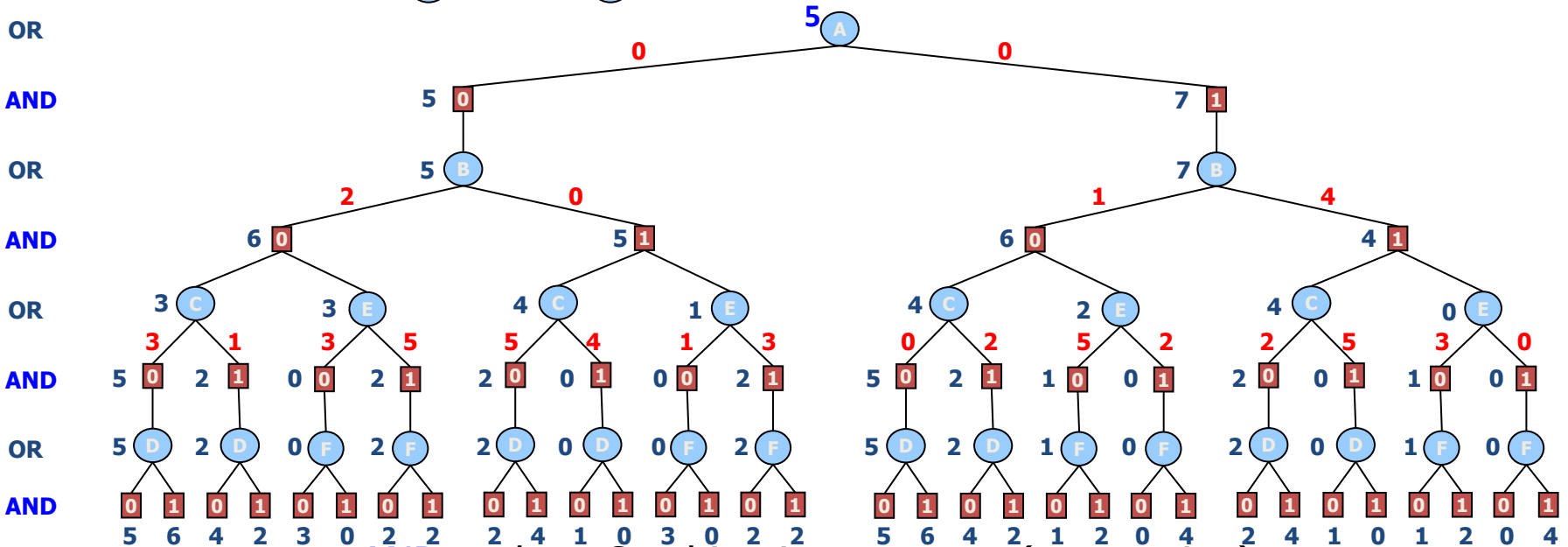
OR – minimization  
AND – summation

# The Value Function for Optimization



A	B	f <sub>1</sub>	A	C	f <sub>2</sub>	A	E	f <sub>3</sub>	A	F	f <sub>4</sub>	B	C	f <sub>5</sub>	B	D	f <sub>6</sub>	B	E	f <sub>7</sub>	C	D	f <sub>8</sub>	E	F	f <sub>9</sub>
0	0	2	0	0	3	0	0	0	0	0	2	0	0	0	0	0	4	0	0	3	0	0	1	0	0	1
0	1	0	0	1	0	0	1	3	0	1	0	0	1	1	0	1	2	0	1	2	0	1	4	0	1	0
1	0	1	1	0	0	1	0	2	1	0	0	1	0	2	1	0	1	1	0	1	1	0	0	1	0	0
1	1	4	1	1	1	1	1	0	1	1	2	1	1	4	1	1	0	1	1	0	1	1	0	1	1	2

Objective function:  $F^* = \min_x \sum_{\alpha} f_{\alpha}(x_{\alpha})$

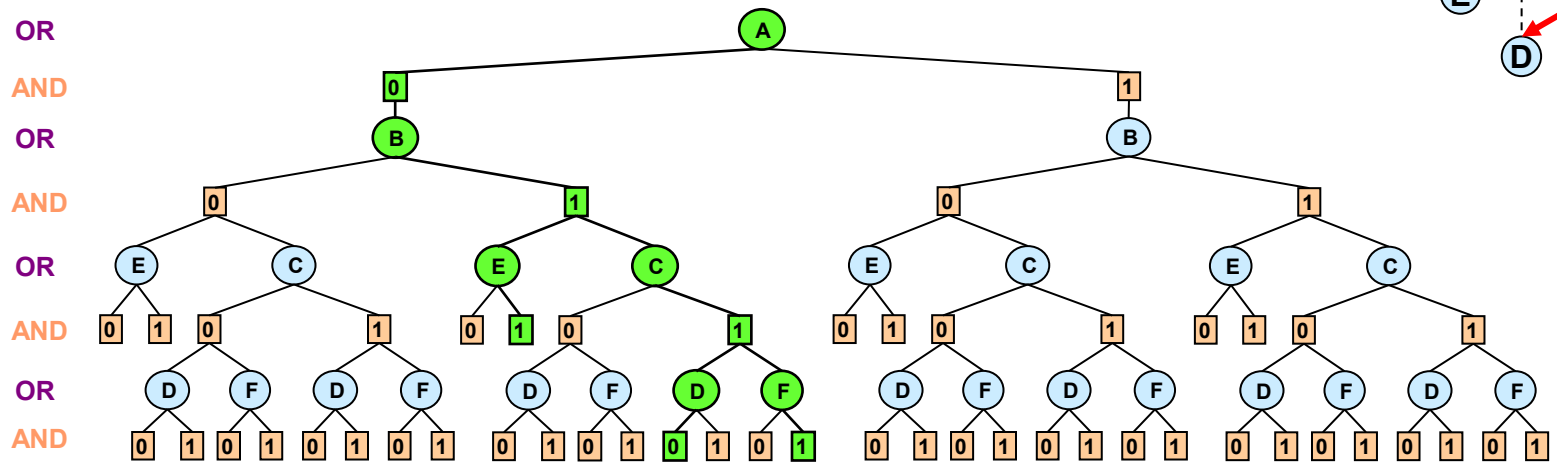
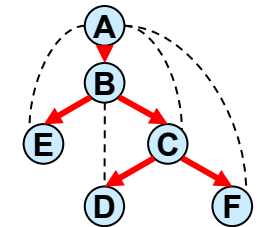
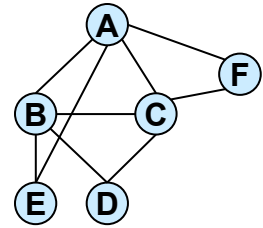


AND node = Combination operator (summation)

OR node = Marginalization operator (minimization)

# Summary: AND/OR Search Trees for GMs

- The AND/OR search tree of R relative to a pseudo-tree, T, has:
  - Alternating levels of: **OR** nodes (variables) and **AND** nodes (values)
- Successor function:
  - The successors of **OR nodes X** are all its consistent values along its path
  - The successors of **AND  $\langle X, v \rangle$**  are all X child variables in T
  - Arc-weight are assigned from the model factors
- A **solution** is a consistent subtree. Its cost, the product of the weights.
- **Query:** compute the value of the root node



# Size and Traversal of AND/OR Search Tree

	AND/OR tree	OR tree
Space	$O(n)$	$O(n)$
Size= Time	$O(n k^h)$ $O(n k^{w^* \log n})$ <small>(Freuder &amp; Quinn85), (Collin, Dechter &amp; Katz91), (Bayardo &amp; Miranker95), (Darwiche01)</small>	$O(k^n)$

$k$  = domain size

$h$  = height of pseudo-tree

$n$  = number of variables

$w^*$  = treewidth

$$h \leq w^* \log n$$

# AND/OR vs. OR Spaces

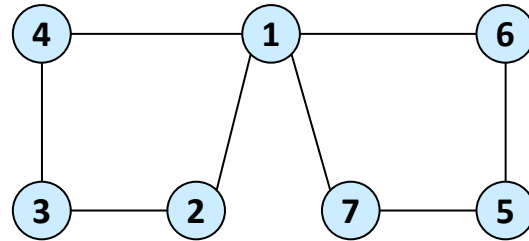
width	height	OR space		AND/OR space		
		Time (sec.)	Nodes	Time (sec.)	AND nodes	OR nodes
5	10	3.15	2,097,150	0.03	<b>10,494</b>	5,247
4	9	3.13	2,097,150	0.01	<b>5,102</b>	2,551
5	10	3.12	2,097,150	0.03	<b>8,926</b>	4,463
4	10	3.12	2,097,150	0.02	<b>7,806</b>	3,903
5	13	3.11	2,097,150	0.10	<b>36,510</b>	18,255

Random graphs with 20 nodes, 20 edges and 2 values per node

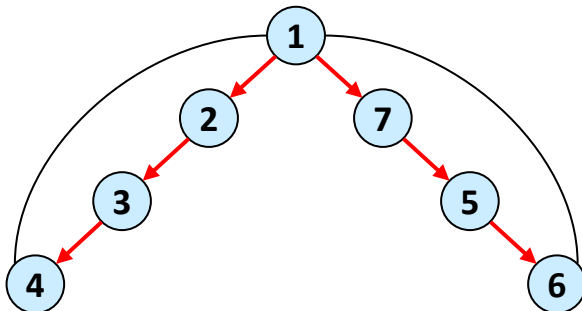
# Pseudo Trees

A **pseudo-tree** of a graph is a tree spanning its nodes, where all arcs in the graph not in the tree are back-arcs

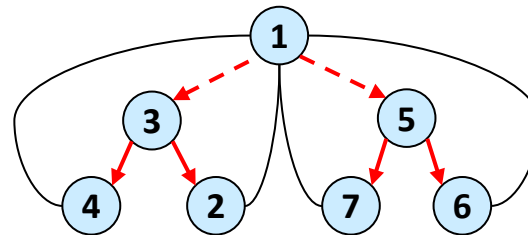
$$h \leq w^* \log n$$



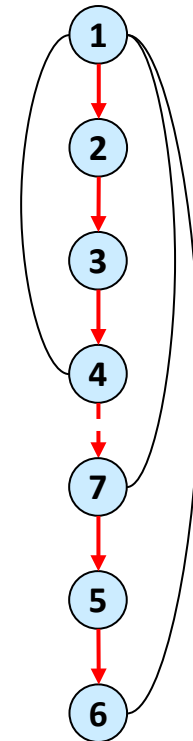
(a) Graph



(b) DFS tree  
height=3

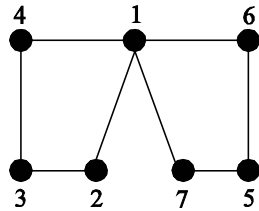


(c) Pseudo tree  
height=2

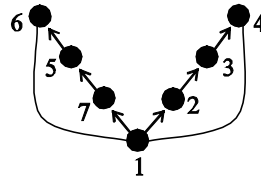


(d) Chain  
height=6

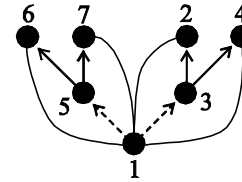
# From DFS-Trees to Pseudo-Trees



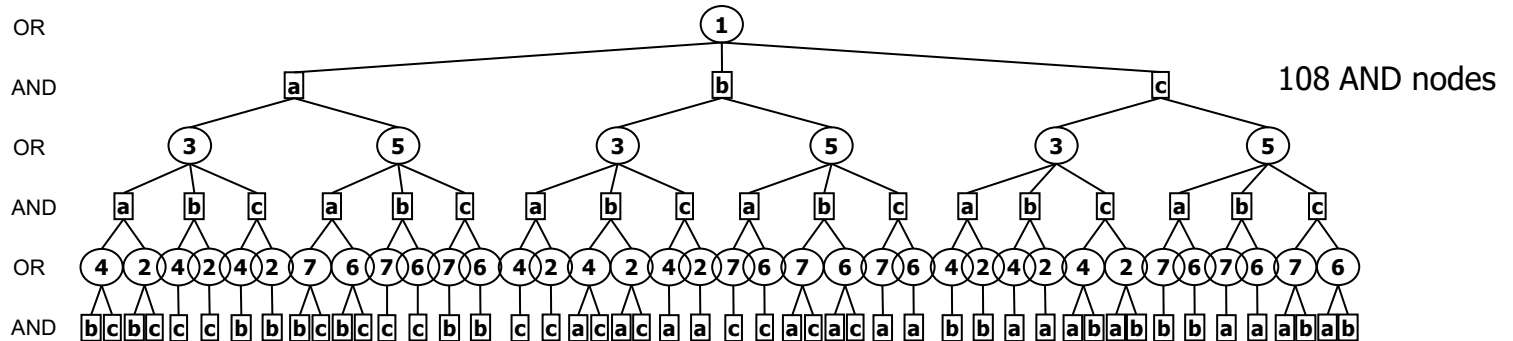
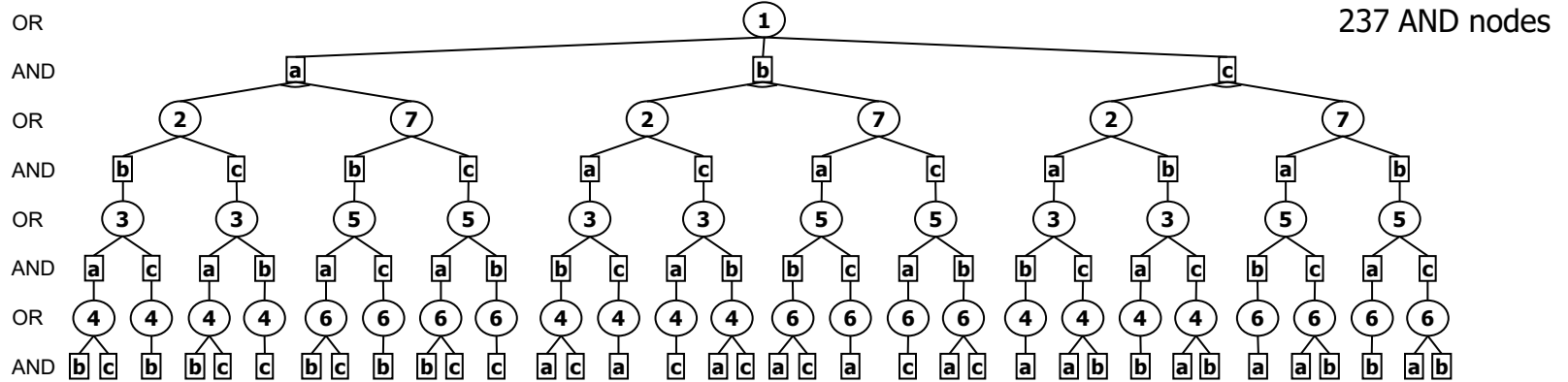
(a)



(b)



(c)





# AND/OR Search-Tree Properties

( $k$  = domain size,  $h$  = pseudo-tree height.  $n$  = number of variables)

- **Theorem:** Any AND/OR search tree based on a pseudo-tree is sound and complete (expresses all and only solutions)
- **Theorem:** Size of AND/OR search tree is  $O(n k^h)$   
Size of OR search tree is  $O(k^n)$
- **Theorem:** Size of AND/OR search tree can be bounded by  $O(\exp(w * \log n))$
- When the pseudo-tree is a chain we get an OR space

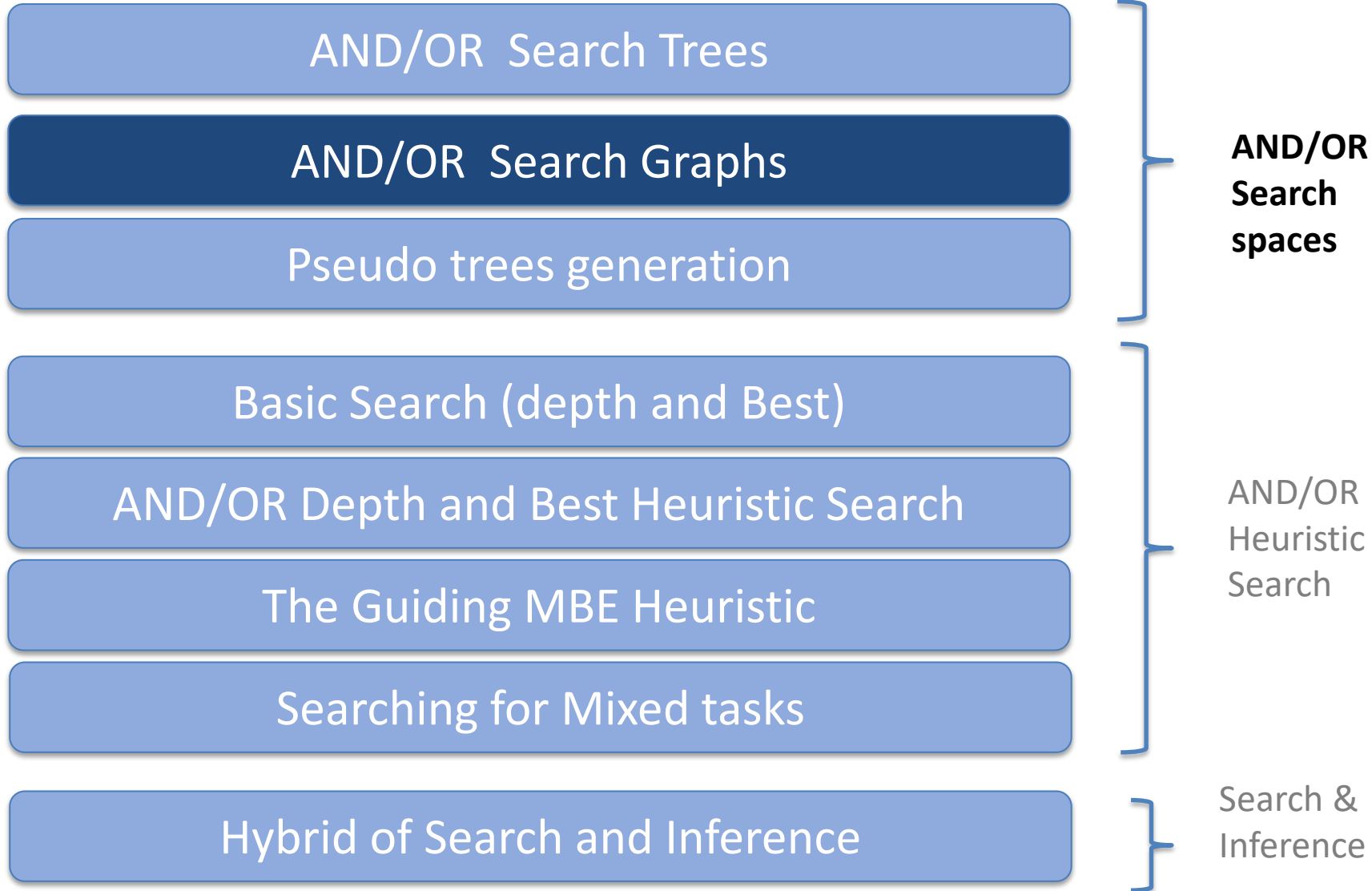
# Summary: Queries and Value of Nodes

- $V(n)$  is the value of the tree  $T(n)$  for the task:
  - **Max-Inference**:  $v(n)$  is the optimal solution in  $T(n)$
  - **Sum-Inference**:  $v(n)$  is probability of evidence in  $T(n)$ .
  - **Mixed-Inference**:  $v(n)$  is the marginal map in  $T(n)$ .
  - **Mixed-Inference**:  $v(n)$  is the max-expect utility in  $T(n)$  of ID.
- **Goal**: compute the value of the root node recursively traversing the AND/OR tree.

Complexity of searching depth-first is

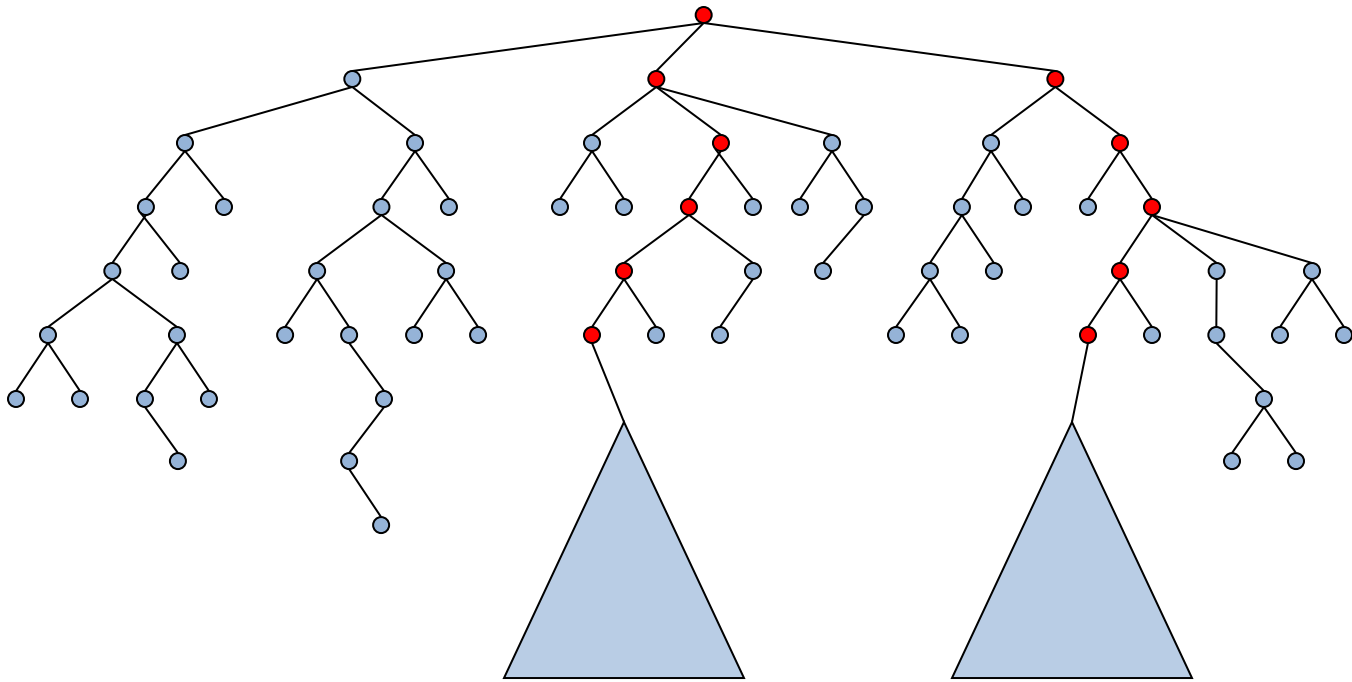
- Space:  $O(n)$
- Time:  $O(nk^h)$
- Time:  $O(k^{w \cdot \log n})$

# Outline: Search



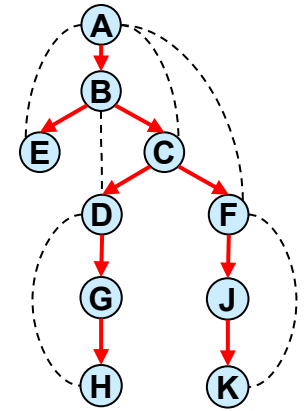
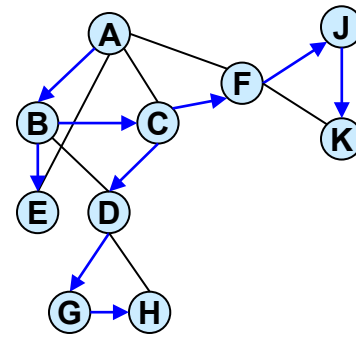
# From Search Trees to Search Graphs

- Any two nodes that root **identical** subtrees or subgraphs can be **merged**





# AND/OR Tree



OR

AND

OR

AND

OR

AND

OR

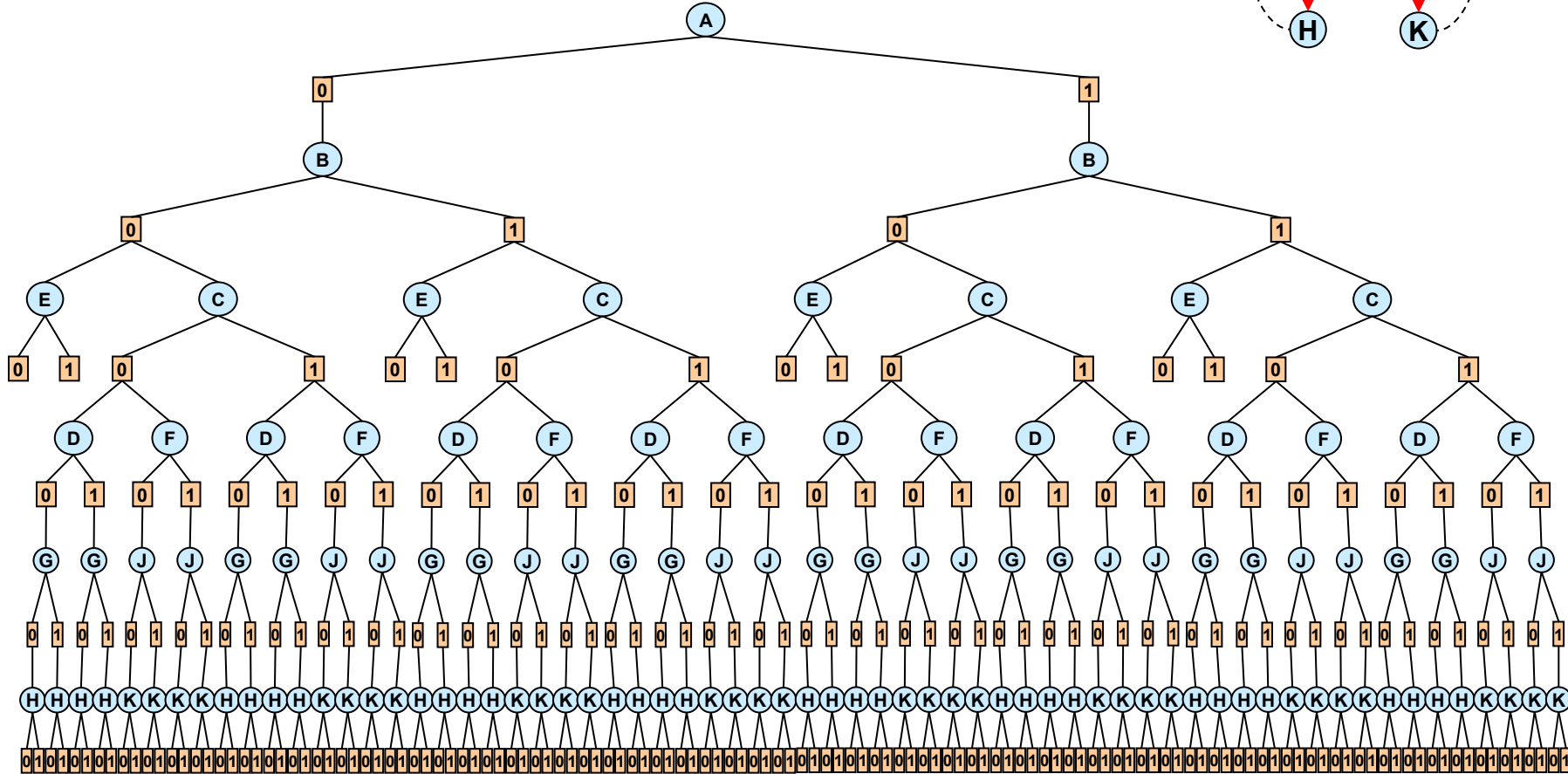
AND

OR

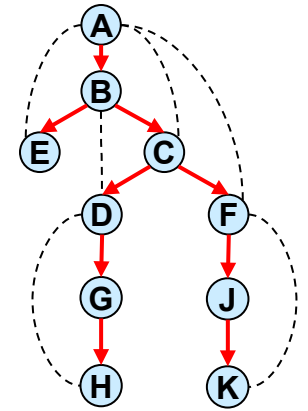
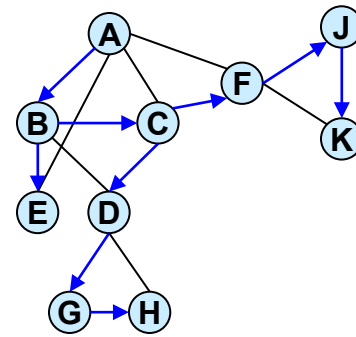
AND

OR

AND



# AND/OR Graph



OR

AND

OR

AND

OR

AND

OR

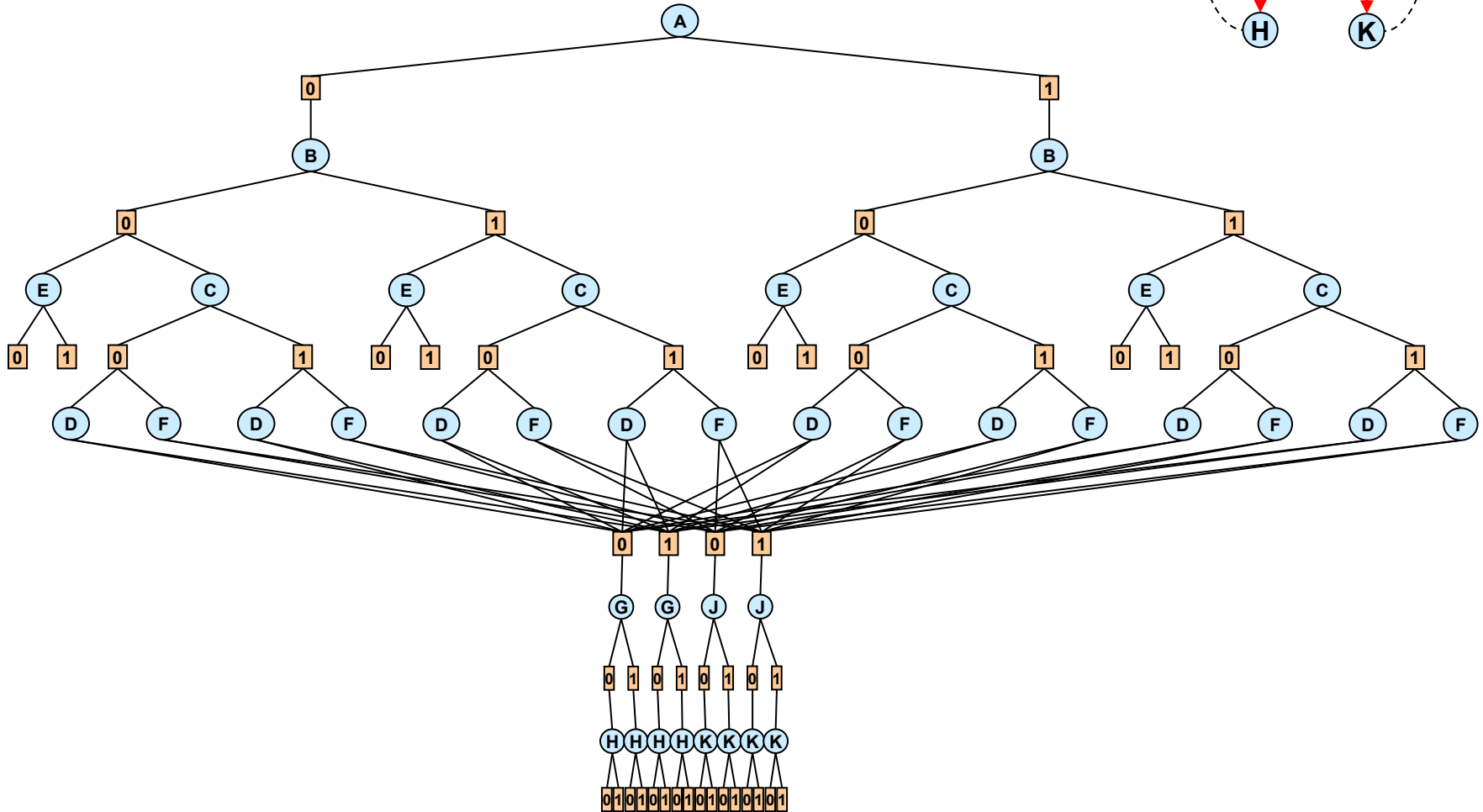
AND

OR

AND

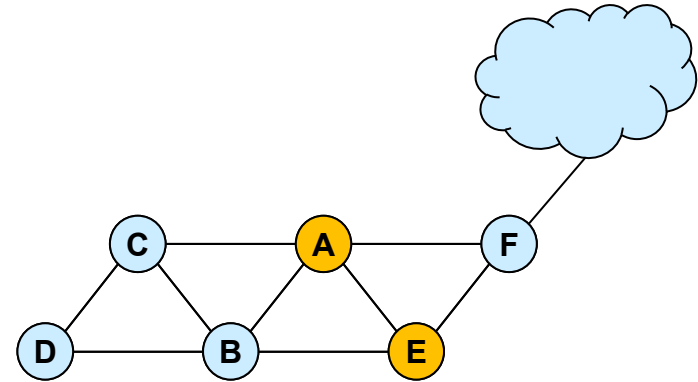
OR

AND

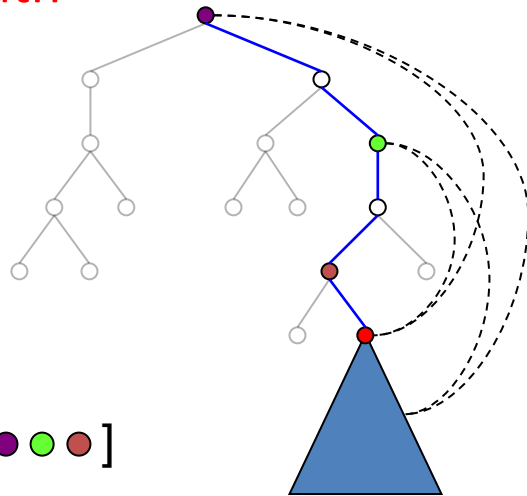


# Merging Based on Context

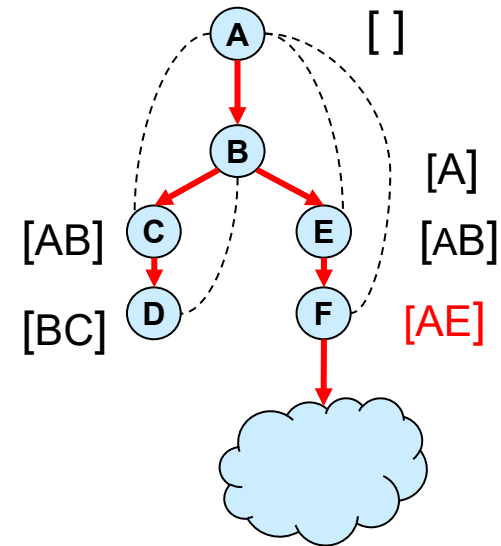
- context (X) = ancestors of X in pseudo tree, connected to X, or to descendants of X
- context (X) = parents in the induced graph
- max |context| = induced width = treewidth



pseudo tree



context(●) = [● ● ●]





# Context-Based Minimal AND/OR Search Graph

**Definition 7.2.13 (context minimal AND/OR search graph)** *The AND/OR search graph of  $M$  guided by a pseudo-tree  $T$  that is closed under context-based merge operator, is called the context minimal AND/OR search graph and is denoted by  $C_T(R)$ .*

# AND/OR Tree DFS Algorithm (Value=Sum-Product)

$$P(E | A, B)$$

A	B	E=0	E=1
0	0	.4	.6
0	1	.5	.5
1	0	.7	.3
1	1	.2	.8

Evidence: E=0

$$P(B | A)$$

A	B=0	B=1
0	.4	.6
1	.1	.9

$$P(C | A)$$

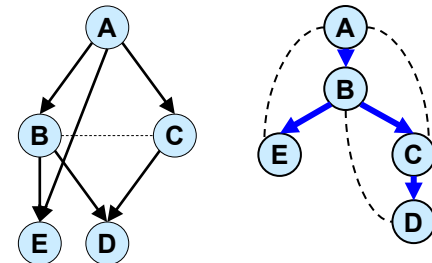
A	C=0	C=1
0	.2	.8
1	.7	.3

$$P(A)$$

A	P(A)
0	.6
1	.4

Result:  $P(D=1, E=0)$

.24408



OR

AND

OR

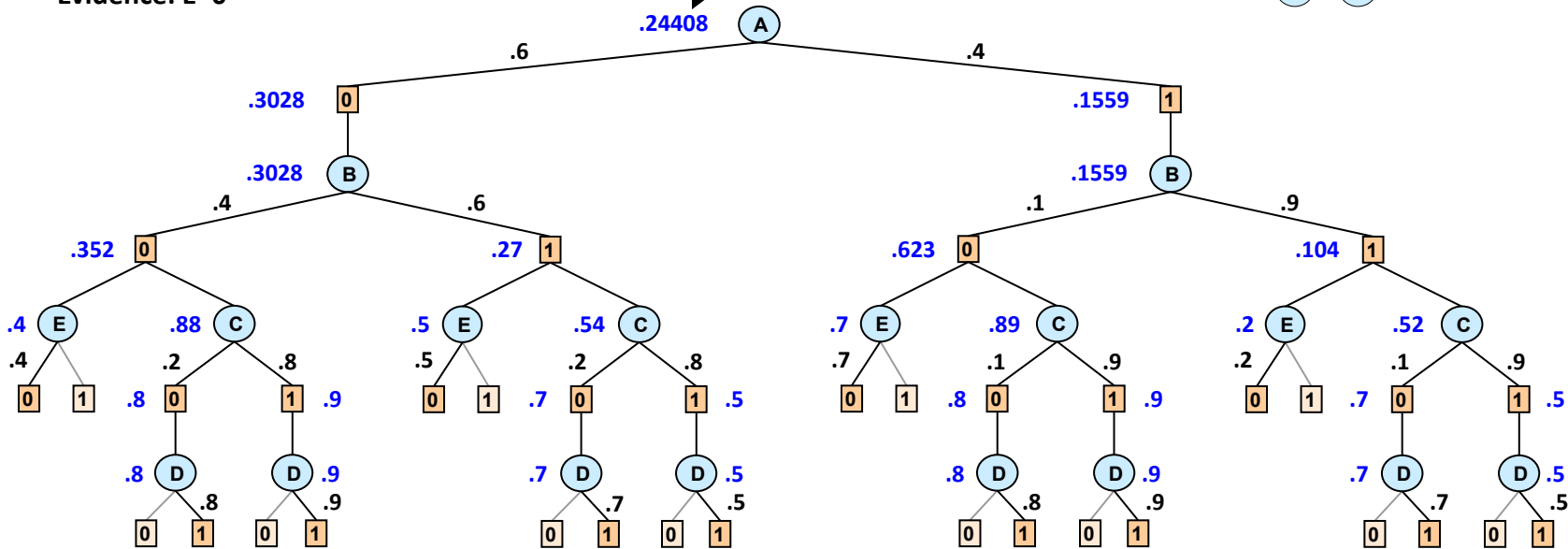
AND

OR

AND

OR

AND



$P(D | B, C)$

B	C	D=0	D=1
0	0	.2	.8
0	1	.1	.9
1	0	.3	.7
1	1	.5	.5

Evidence: D=1

# AND/OR Search Graph (Value=Sum-Product)

$P(E | A, B)$

A	B	E=0	E=1
0	0	.4	.6
0	1	.5	.5
1	0	.7	.3
1	1	.2	.8

Evidence: E=0

$P(B | A)$

A	B=0	B=1
0	.4	.6
1	.1	.9

$P(C | A)$

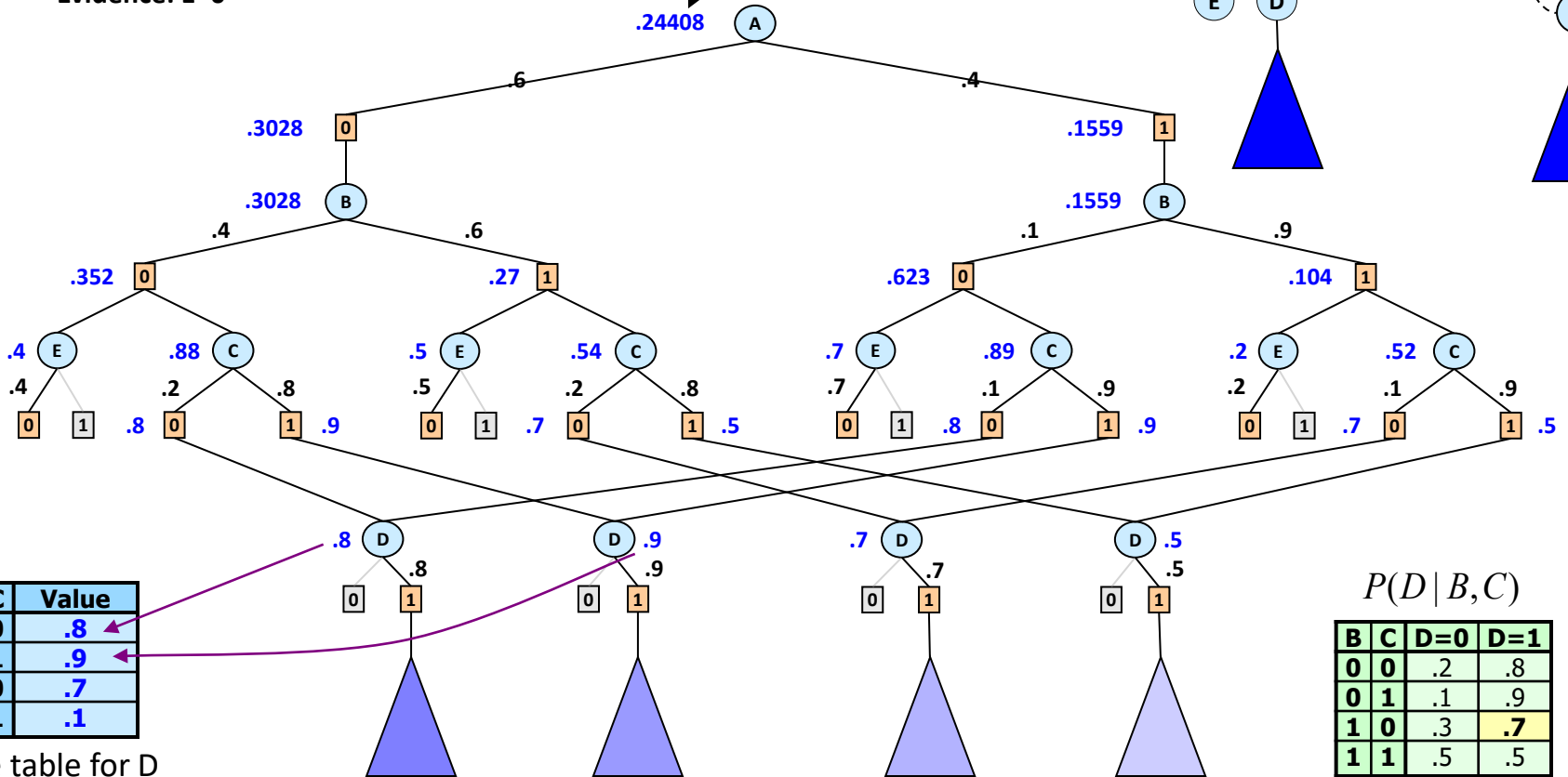
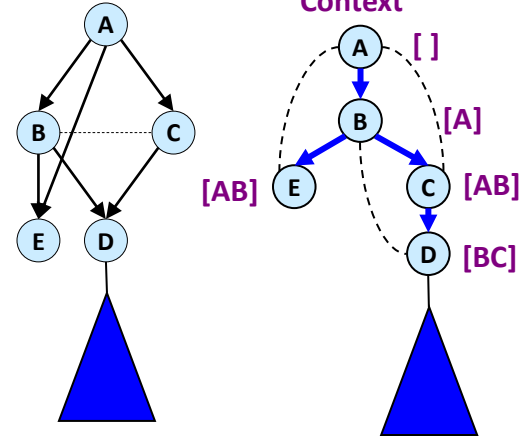
A	C=0	C=1
0	.2	.8
1	.7	.3

$P(A)$

A	P(A)
0	.6
1	.4

Result:  $P(D=1, E=0)$

.24408



B	C	Value
0	0	.8
0	1	.9
1	0	.7
1	1	.1

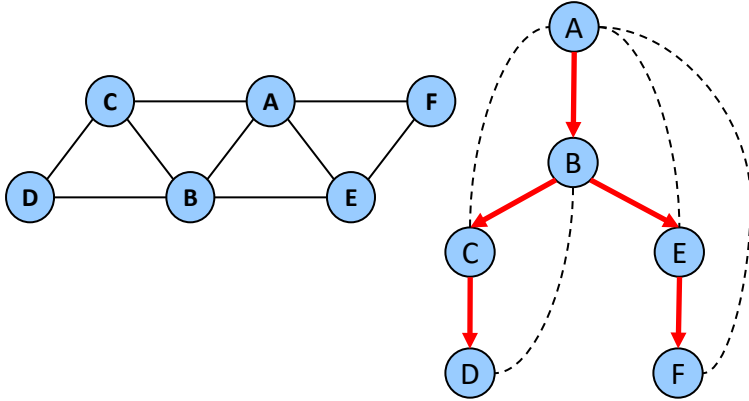
Cache table for D

$P(D | B, C)$

B	C	D=0	D=1
0	0	.2	.8
0	1	.1	.9
1	0	.3	.7
1	1	.5	.5

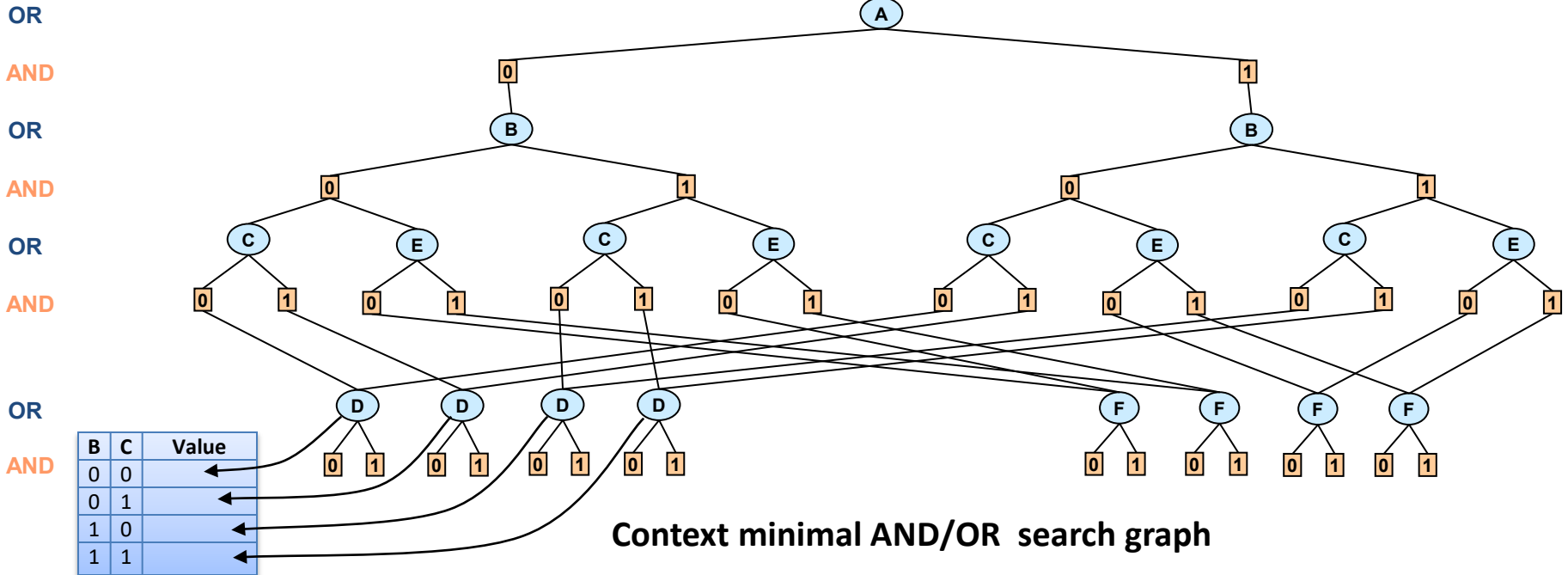
Evidence: D=1

# AND/OR Search Graph (Optimization)



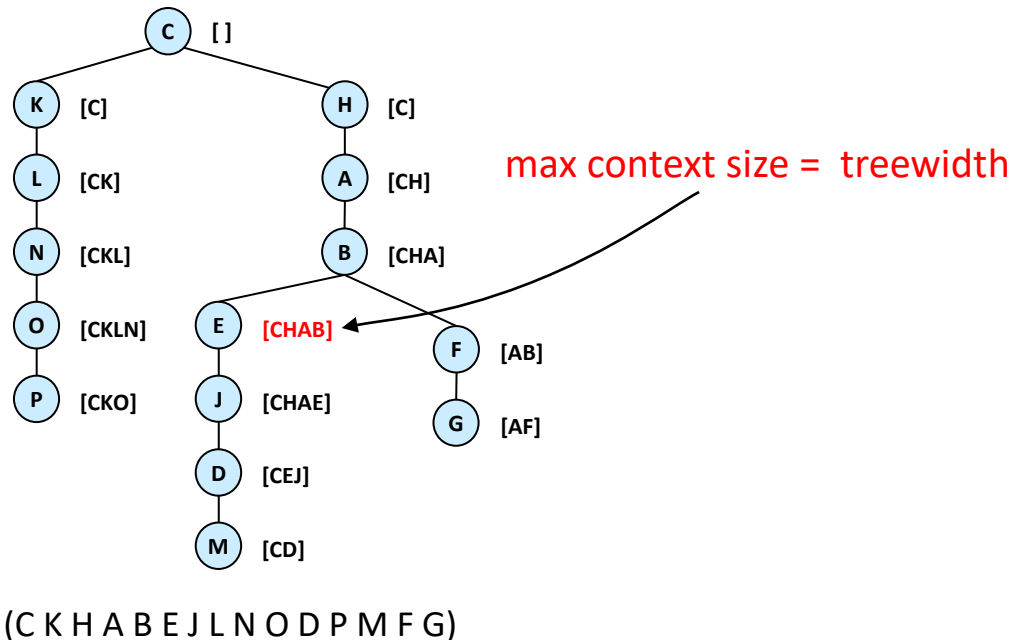
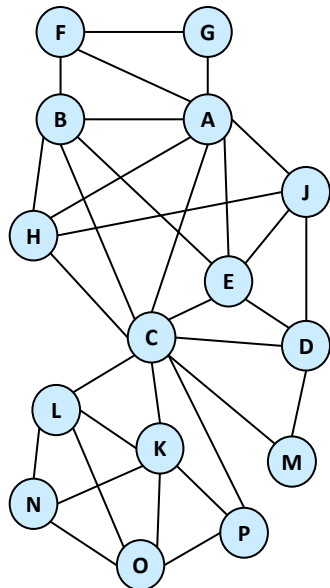
A	B	f <sub>1</sub>	A	C	f <sub>2</sub>	A	E	f <sub>3</sub>	A	F	f <sub>4</sub>	B	C	f <sub>5</sub>	B	D	f <sub>6</sub>	B	E	f <sub>7</sub>	C	D	f <sub>8</sub>	E	F	f <sub>9</sub>
0	0	2	0	0	3	0	0	0	0	0	2	0	0	0	0	0	4	0	0	3	0	0	1	0	0	1
0	1	0	0	1	0	0	1	3	0	1	0	0	1	1	0	1	2	0	1	2	0	1	4	0	1	0
1	0	1	1	0	0	1	0	2	1	0	0	1	0	2	1	0	1	1	0	1	1	0	0	1	0	0
1	1	4	1	1	1	1	1	0	1	1	2	1	1	4	1	1	0	1	1	0	1	1	0	1	1	2

Objective function:  $F^* = \min_x \sum_{\alpha} f_{\alpha}(x_{\alpha})$

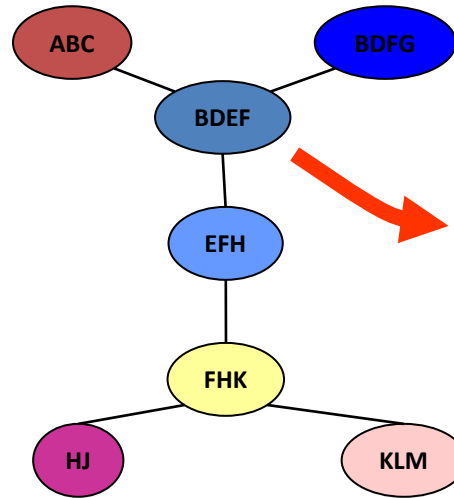
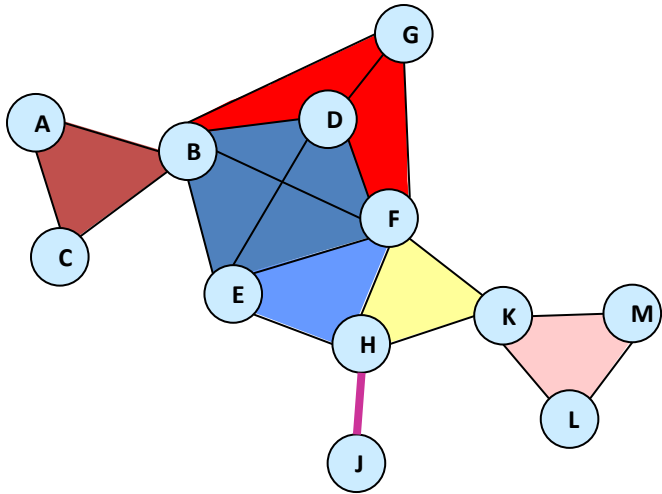


# How Big Is The Context?

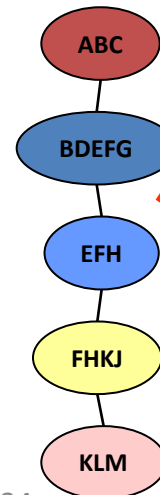
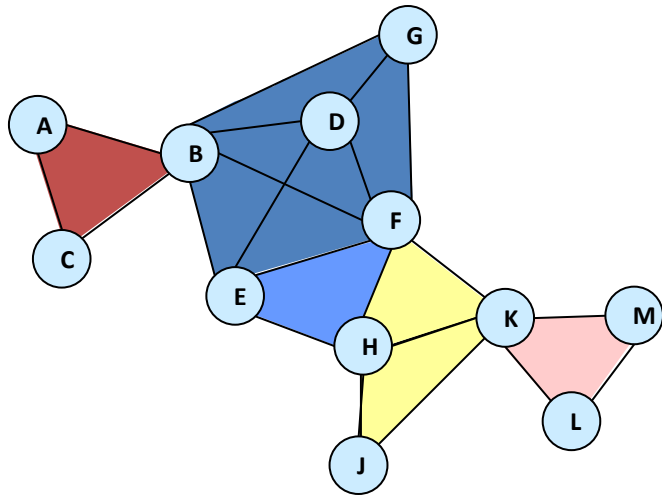
- Theorem:** The maximum context-size of a pseudo-tree equals the **treewidth** along the pseudo tree.



# Treewidth vs. Pathwidth

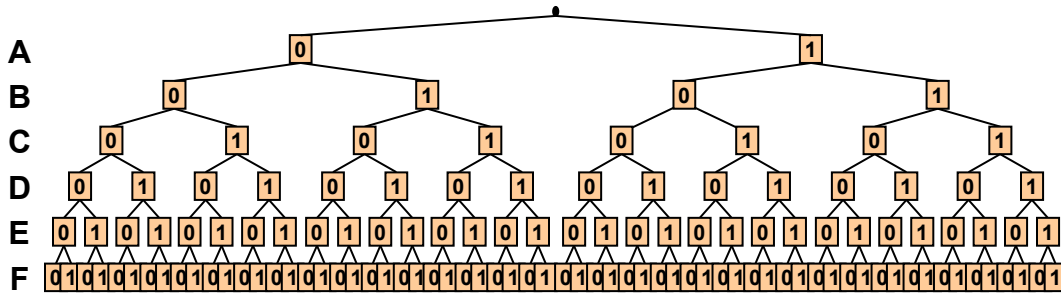
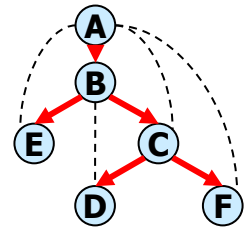


**treewidth = 3**  
 = (max cluster size) - 1



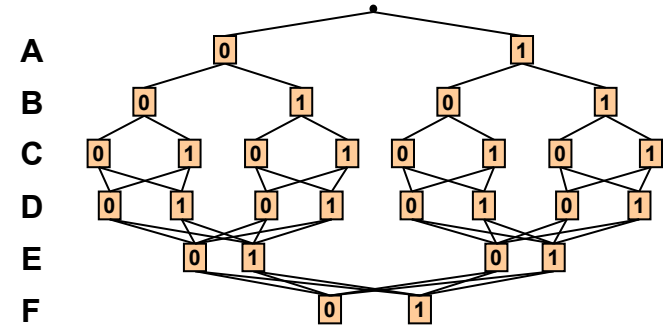
**pathwidth = 4**  
 = (max cluster size) - 1

# All Four Search Spaces



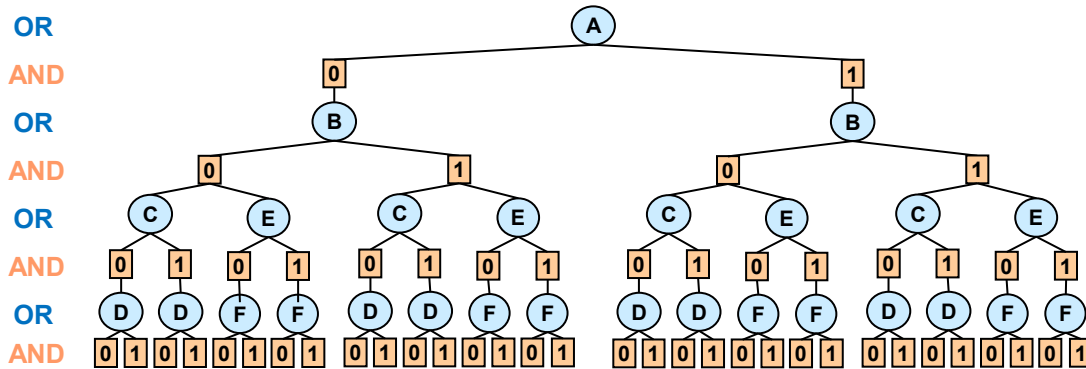
Full OR search tree

126 nodes



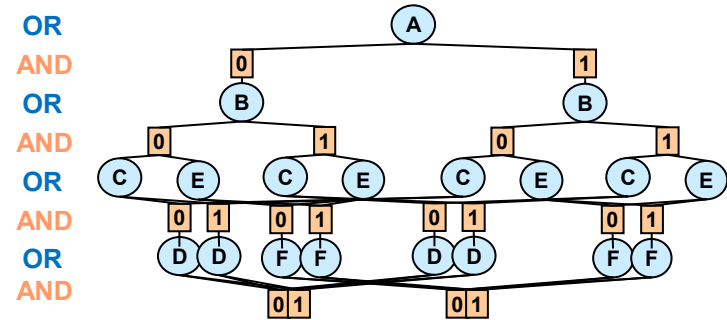
Context minimal OR search graph

28 nodes



Full AND/OR search tree

54 AND nodes



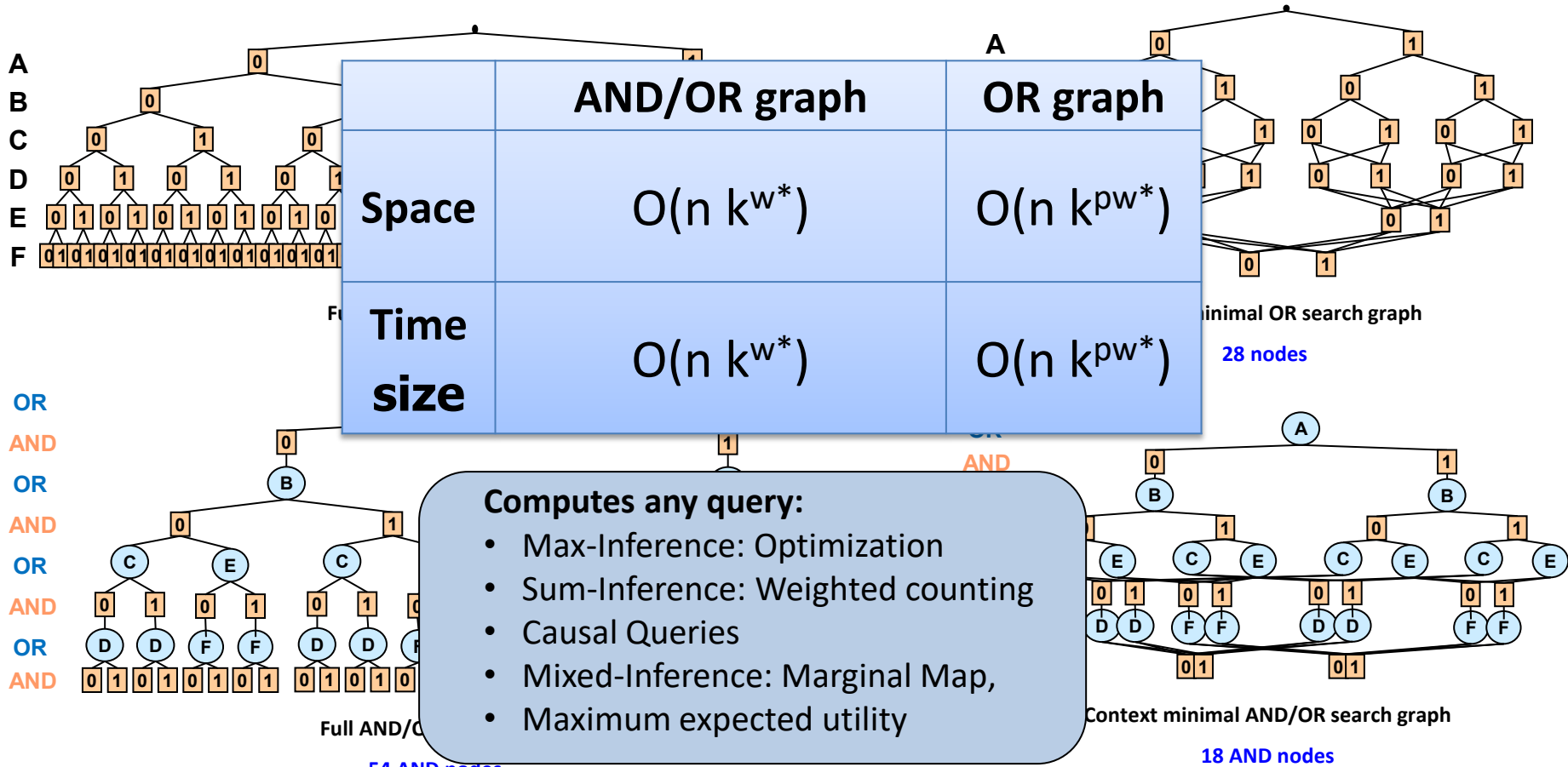
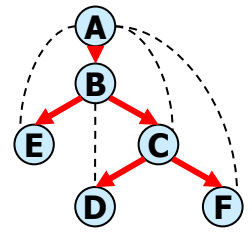
Context minimal AND/OR search graph

18 AND nodes

$k$  = domain size  
 $n$  = number of variables  
 $w^*$  = treewidth  
 $pw^*$  = pathwidth

Any query is best computed  
 over the context-minimal AND/OR space

# All Four Search Spaces



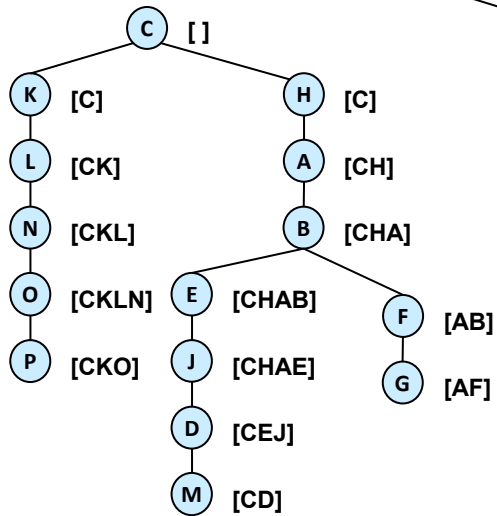
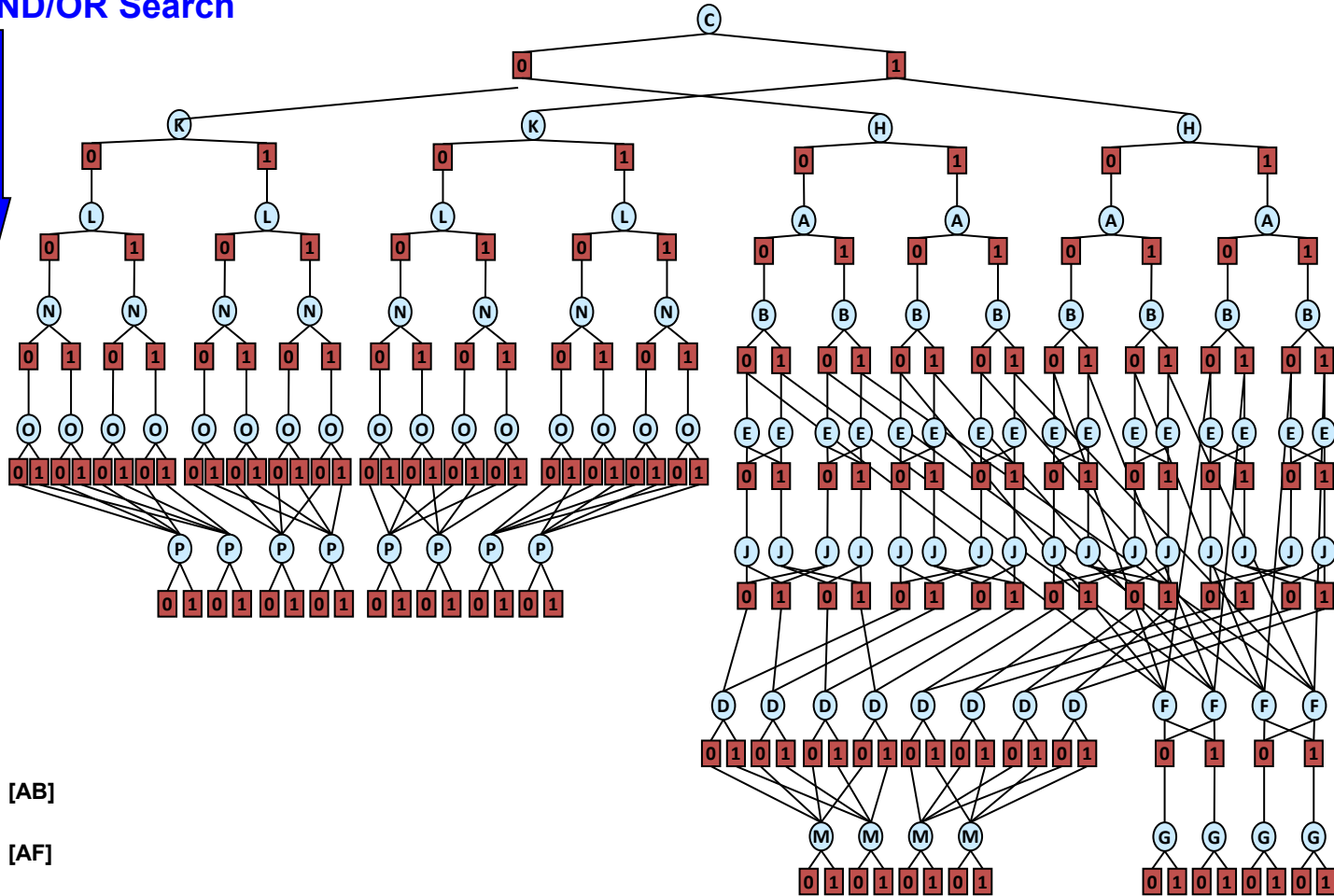
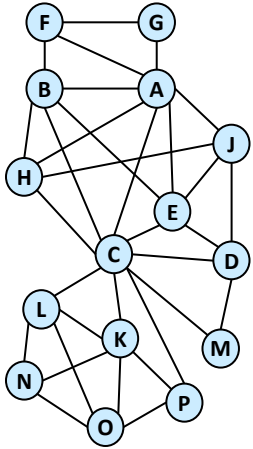
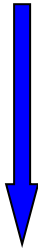
$k$  = domain size  
 $n$  = number of variables  
 $w^*$  = treewidth  
 $pw^*$  = pathwidth

**Any query is best computed over the context-minimal AND/OR space**



# AND/OR Search and Variable Elimination

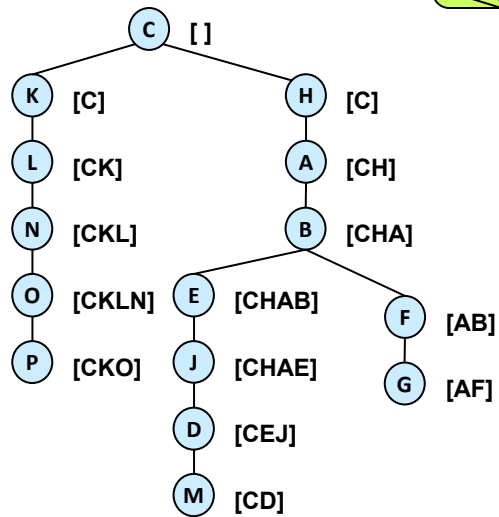
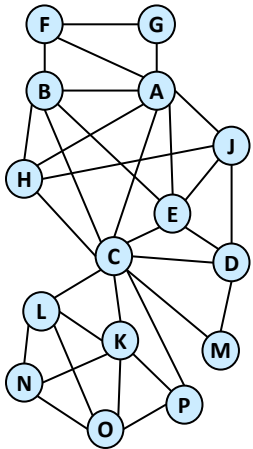
AND/OR Search



(CKHABEJLNODPMFG)

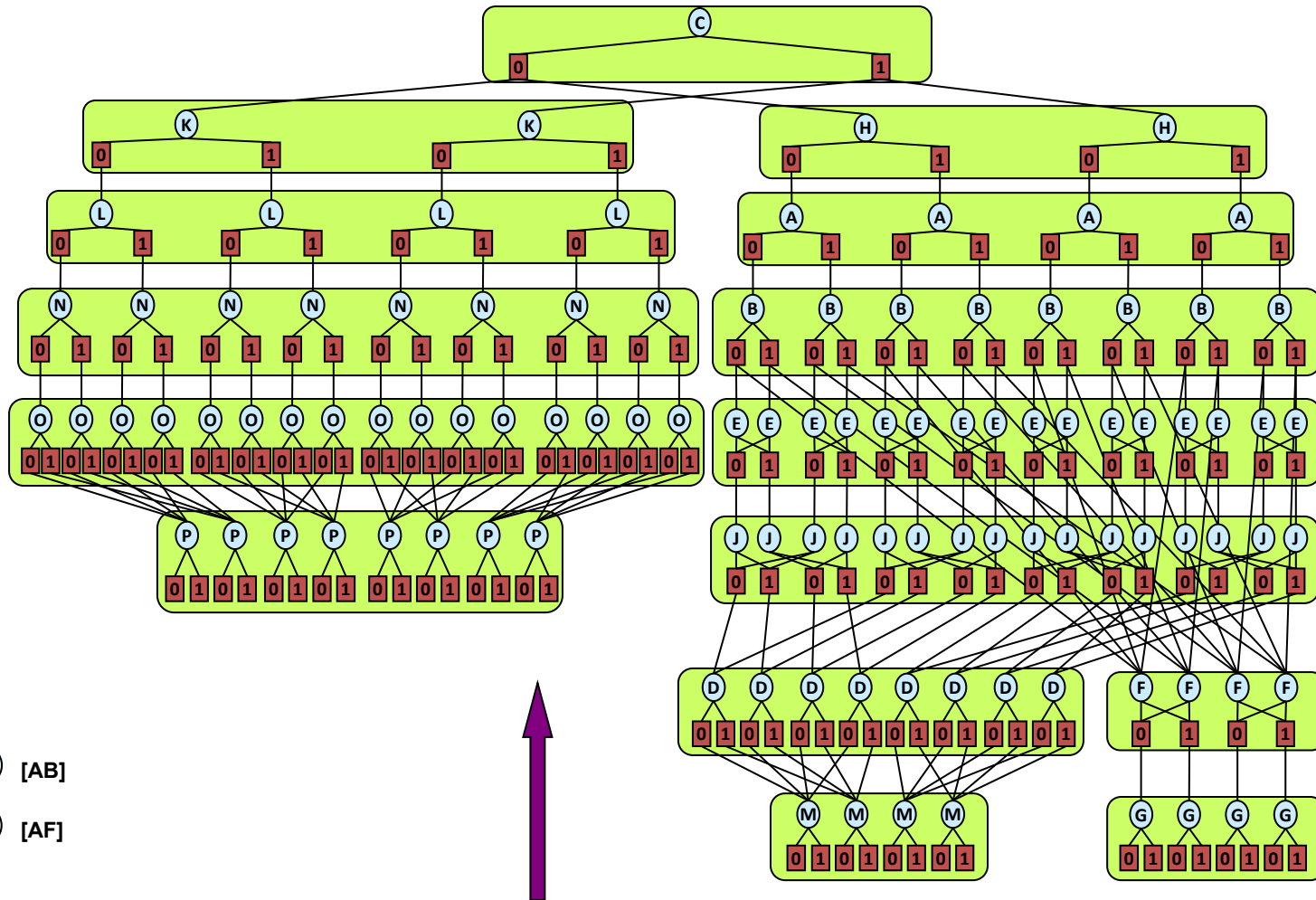
Dechter & Ihler

# AND/OR Search and Variable Elimination



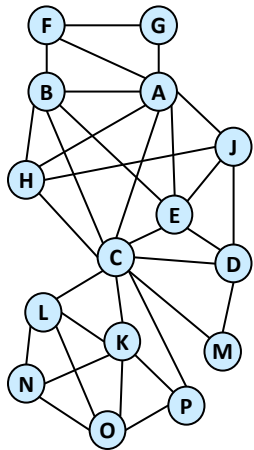
(CKHABEJLNODPMFG)

Dechter & Ihler

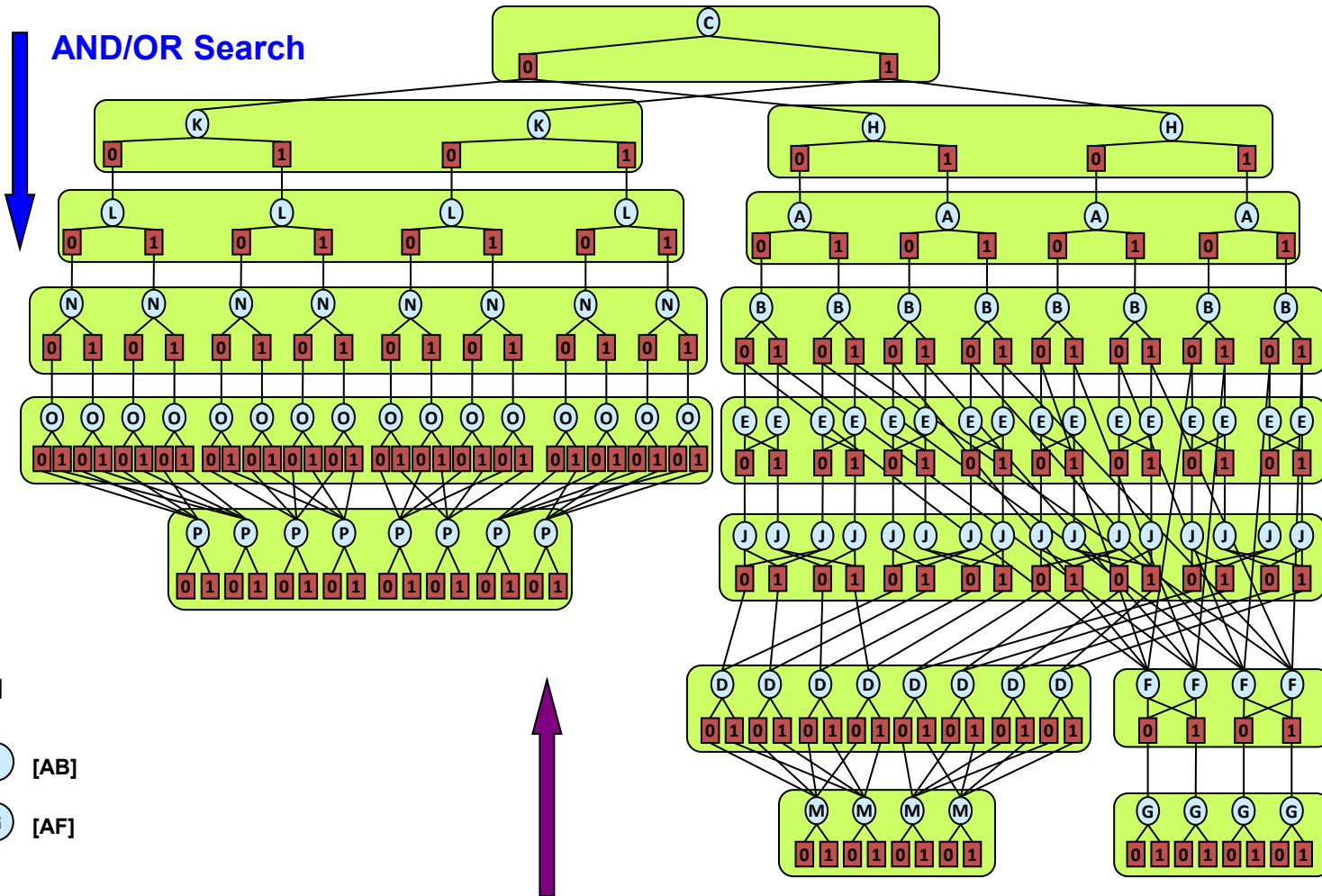


Variable Elimination

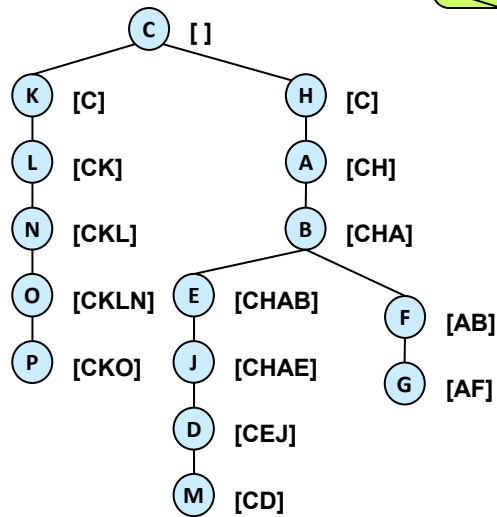
# AND/OR Search and Variable Elimination



AND/OR Search



Variable Elimination



(CKHABEJLNODPMFG)  
Dechter & Ihler

# Outline: Search



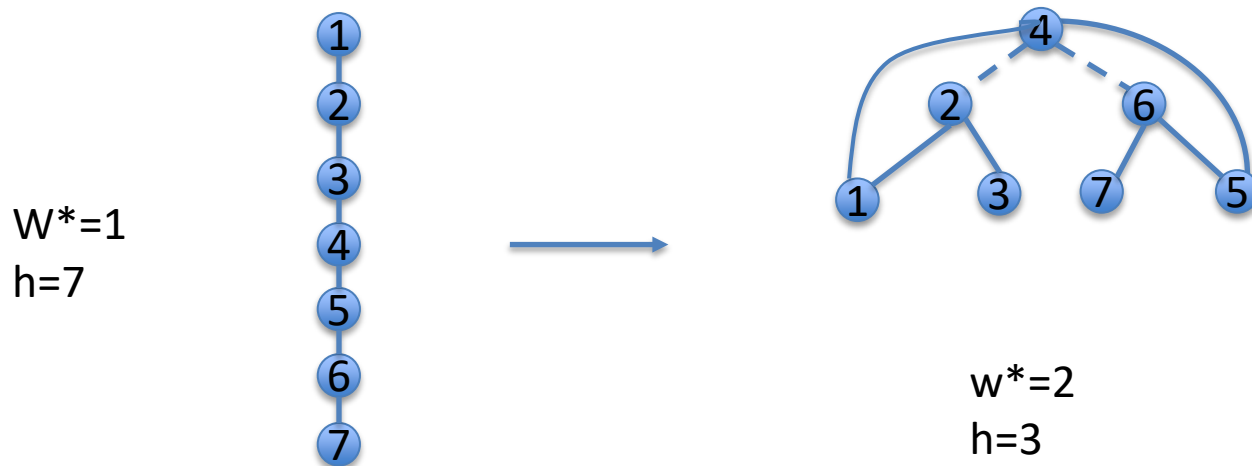
**AND/OR  
Search  
spaces**

AND/OR  
Heuristic  
Search

Search &  
Inference

# Finding Min-height Pseudo-Trees

- Finding a min height pseudo-tree is NP-complete, but:
- Given a tree-decomposition with treewidth  $w^*$ , there exists a pseudo-tree whose height satisfies
  - $h \leq w^* \log n$
- Optimality of  $h$  and  $w^*$  cannot be achieved at once.



# Constructing Pseudo-Trees

- **Min-Fill** [Kjaerulff, 1990]
  - Depth-first traversal of the induced graph obtained along the **min-fill** elimination order
  - Variables ordered according to the smallest “fill-set”
- **Hypergraph Partitioning** [Karypis and Kumar, 2000]
  - Functions are vertices in the hypergraph and variables are hyperedges
  - Recursive decomposition of the hypergraph while minimizing the separator size at each step
  - Using state-of-the-art software package **hMeTiS**

# Quality of Pseudo-Trees

Network	hypergraph		min-fill	
	width	depth	width	depth
barley	7	<b>13</b>	7	23
diabetes	7	<b>16</b>	4	77
link	21	<b>40</b>	15	53
mildew	5	<b>9</b>	4	13
munin1	12	<b>17</b>	12	29
munin2	9	<b>16</b>	9	32
munin3	9	<b>15</b>	9	30
munin4	9	<b>18</b>	9	30
water	11	<b>16</b>	10	15
pigs	11	<b>20</b>	11	26

Bayesian Networks Repository

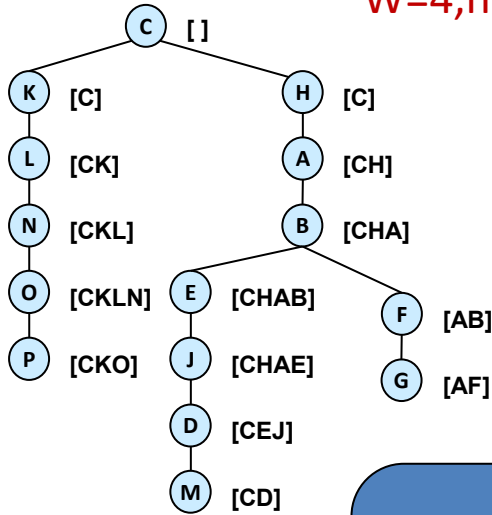
Network	hypergraph		min-fill	
	width	depth	width	depth
spot5	47	152	<b>39</b>	204
spot28	108	138	<b>79</b>	199
spot29	16	23	<b>14</b>	42
spot42	36	48	<b>33</b>	87
spot54	12	16	<b>11</b>	33
spot404	19	26	<b>19</b>	42
spot408	47	52	<b>35</b>	97
spot503	11	20	<b>9</b>	39
spot505	29	42	<b>23</b>	74
spot507	70	122	<b>59</b>	160

SPOT5 Benchmarks

For more see [Dechter 2003]

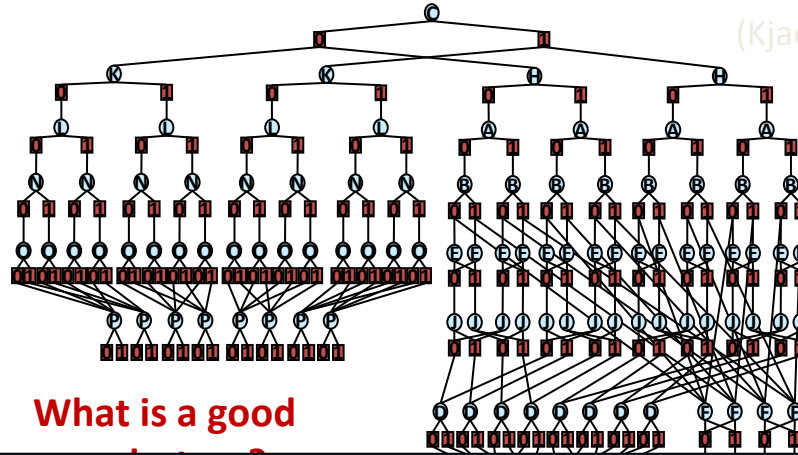
# The Impact of the Pseudo-Tree

$W=4, h=8$

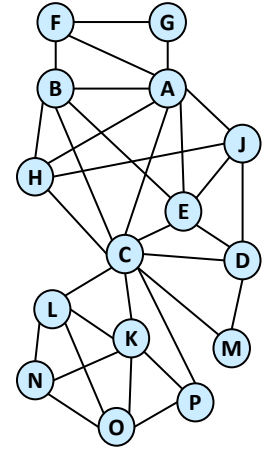


(CKHABEJLN O)

Min-Fill  
(Kjaerulff90)



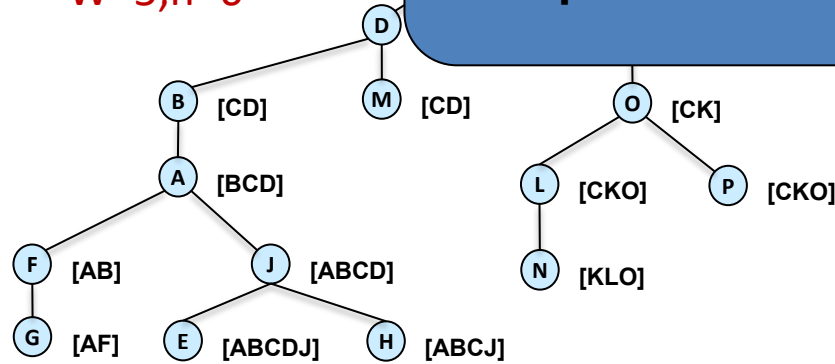
What is a good



graph  
oning  
(s)

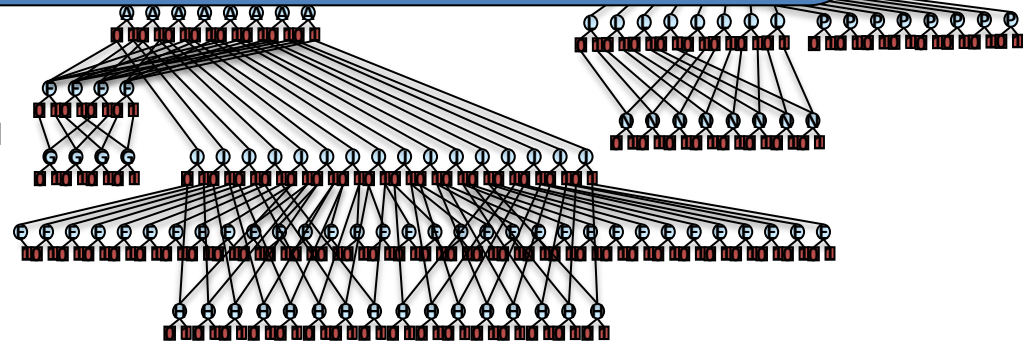
- Choose pseudo-tree with a minimal search graph
- But determinism and pruning for optimization is unpredictable

$W=5, h=6$



DCSA12084hler

(CDKBAOMLNPJHEFG)





# Outline: Search



# Probabilistic Reasoning Problems

- Exact Inference by elimination or search
- Complexity:

Causal effects

Max-Inference:	$f(x^*) = \max_x \prod_{\alpha} f_{\alpha}(x_{\alpha})$
Sum-Inference:	$Z = \sum_x \prod_{\alpha} f_{\alpha}(x_{\alpha})$
Mixed-Inference (MMAP):	$f_M(x_M^*) = \max_{x_M} \sum_{x_S} \prod_{\alpha} f_{\alpha}(x_{\alpha})$
Mixed-Inference (MEU):	$MEU = \max_{\tilde{x}} \sum_i ( \prod P_i ) \times ( \sum r_i )$

$e^{\text{tree-width}}$

Harder

- All solved by AND/OR Depth-first search,
  - Linear memory,  $\exp(h)$  time or
  - $\exp(w^*)$  memory and time
- But, we can do better by:
  - Pruning while searching
  - Generating upper and lower bounds anytime

# AND/OR Tree DFS Algorithm (Belief Updating)

$$P(E | A, B)$$

A	B	E=0	E=1
0	0	.4	.6
0	1	.5	.5
1	0	.7	.3
1	1	.2	.8

Evidence: E=0

$$P(B | A)$$

A	B=0	B=1
0	.4	.6
1	.1	.9

$$P(C | A)$$

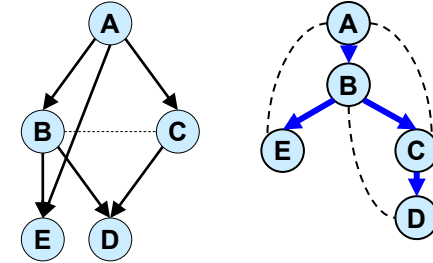
A	C=0	C=1
0	.2	.8
1	.7	.3

$$P(A)$$

A	P(A)
0	.6
1	.4

Result:  $P(D=1, E=0)$

.24408



OR

AND

OR

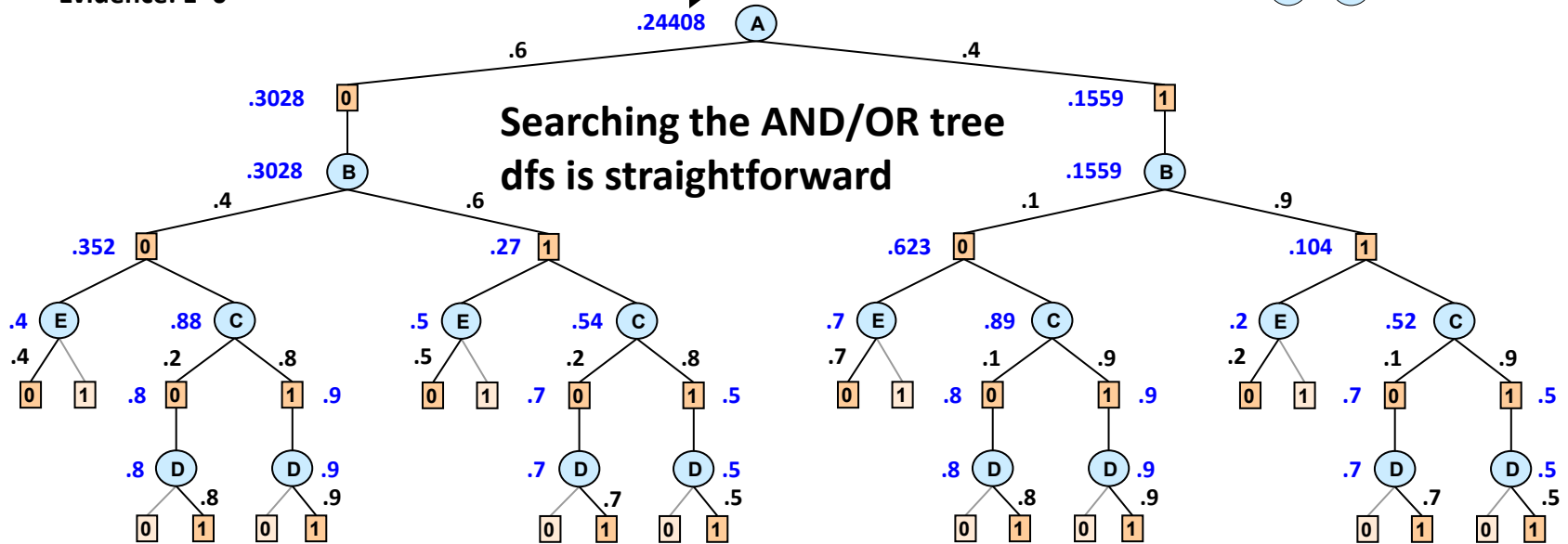
AND

OR

AND

OR

AND



$$P(D | B, C)$$

B	C	D=0	D=1
0	0	.2	.8
0	1	.1	.9
1	0	.3	.7
1	1	.5	.5

Evidence: D=1

OR node: Marginalization operator (summation)

AND node: Combination operator (product)

Value of node = updated belief for sub-problem below

# AND/OR Graph DFS Algorithm (Belief Updating)

$P(E | A, B)$

A	B	E=0	E=1
0	0	.4	.6
0	1	.5	.5
1	0	.7	.3
1	1	.2	.8

Evidence: E=0

$P(B | A)$

A	B=0	B=1
0	.4	.6
1	.1	.9

$P(C | A)$

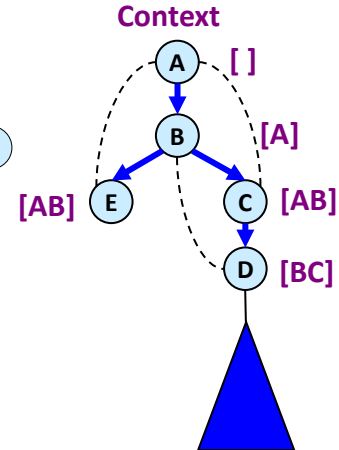
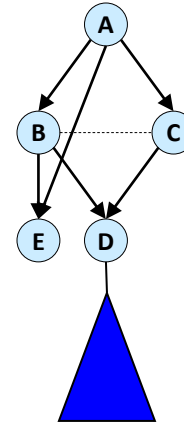
A	C=0	C=1
0	.2	.8
1	.7	.3

$P(A)$

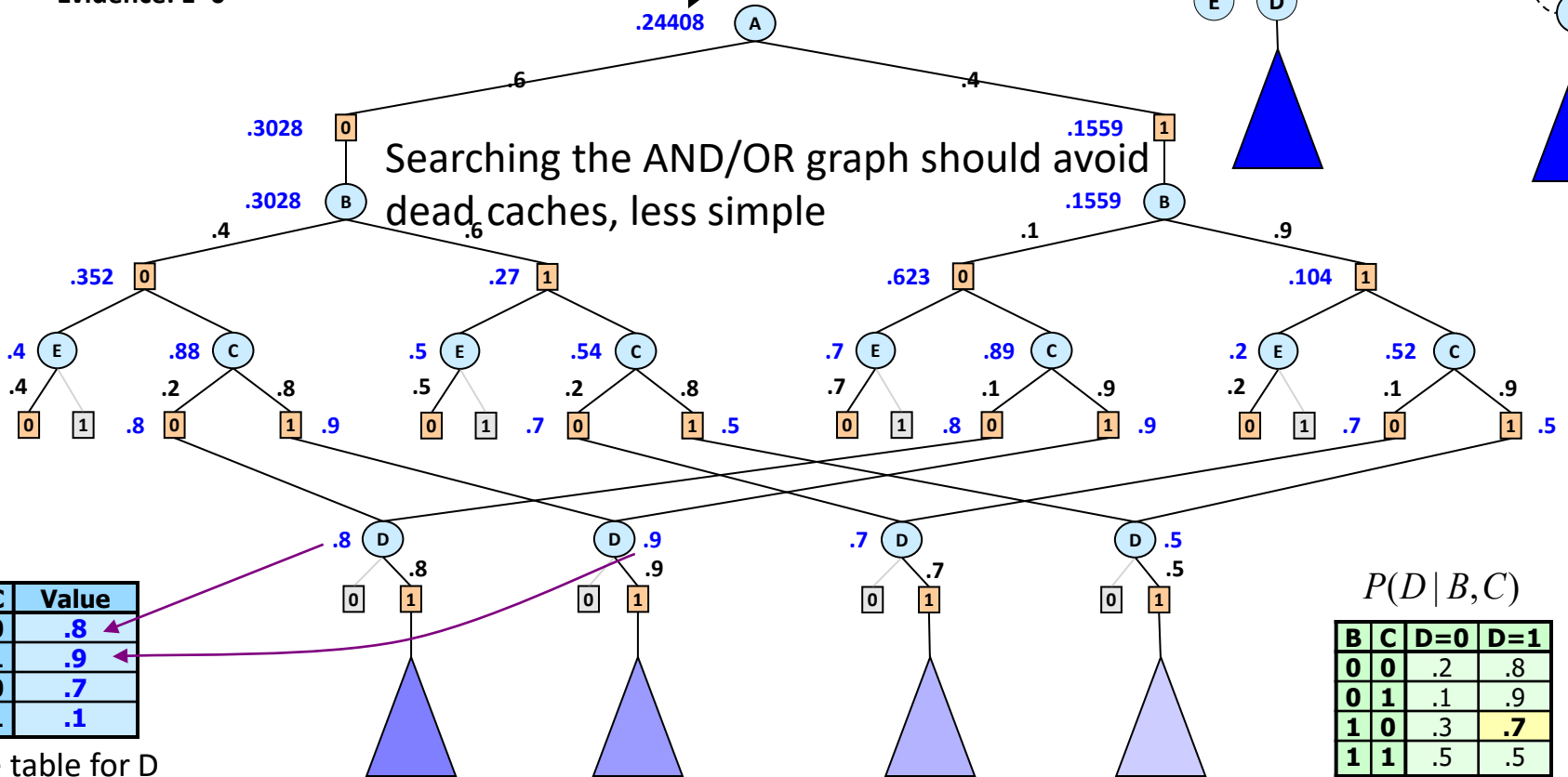
A	P(A)
0	.6
1	.4

Result:  $P(D=1, E=0)$

.24408



Searching the AND/OR graph should avoid dead caches, less simple



B	C	Value
0	0	.8
0	1	.9
1	0	.7
1	1	.1

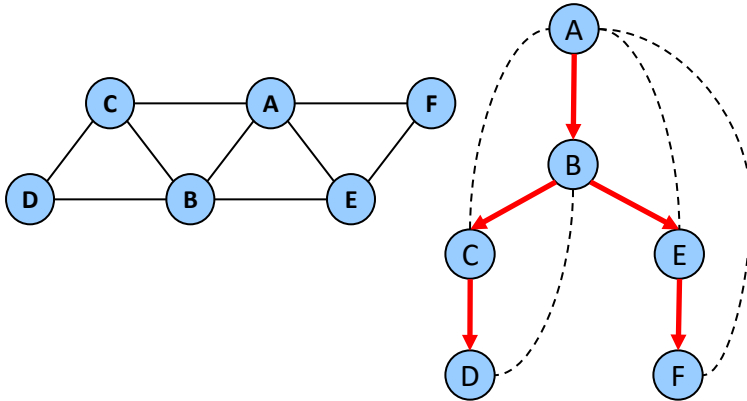
Cache table for D

$P(D | B, C)$

B	C	D=0	D=1
0	0	.2	.8
0	1	.1	.9
1	0	.3	.7
1	1	.5	.5

Evidence: D=1

# AND/OR Search Graph (Optimization)



A	B	f <sub>1</sub>	A	C	f <sub>2</sub>	A	E	f <sub>3</sub>	A	F	f <sub>4</sub>	B	C	f <sub>5</sub>	B	D	f <sub>6</sub>	B	E	f <sub>7</sub>	C	D	f <sub>8</sub>	E	F	f <sub>9</sub>
0	0	2	0	0	3	0	0	0	0	0	2	0	0	0	0	0	4	0	0	3	0	0	1	0	0	1
0	1	0	0	1	0	0	1	3	0	1	0	0	1	1	0	1	2	0	1	2	0	1	4	0	1	0
1	0	1	1	0	0	1	0	2	1	0	0	1	0	2	1	0	1	1	0	1	1	0	0	1	0	0
1	1	4	1	1	1	1	1	0	1	1	2	1	1	4	1	1	0	1	1	0	1	1	0	1	1	2

Objective function:  $F^* = \min_x \sum_{\alpha} f_{\alpha}(x_{\alpha})$

OR

AND

OR

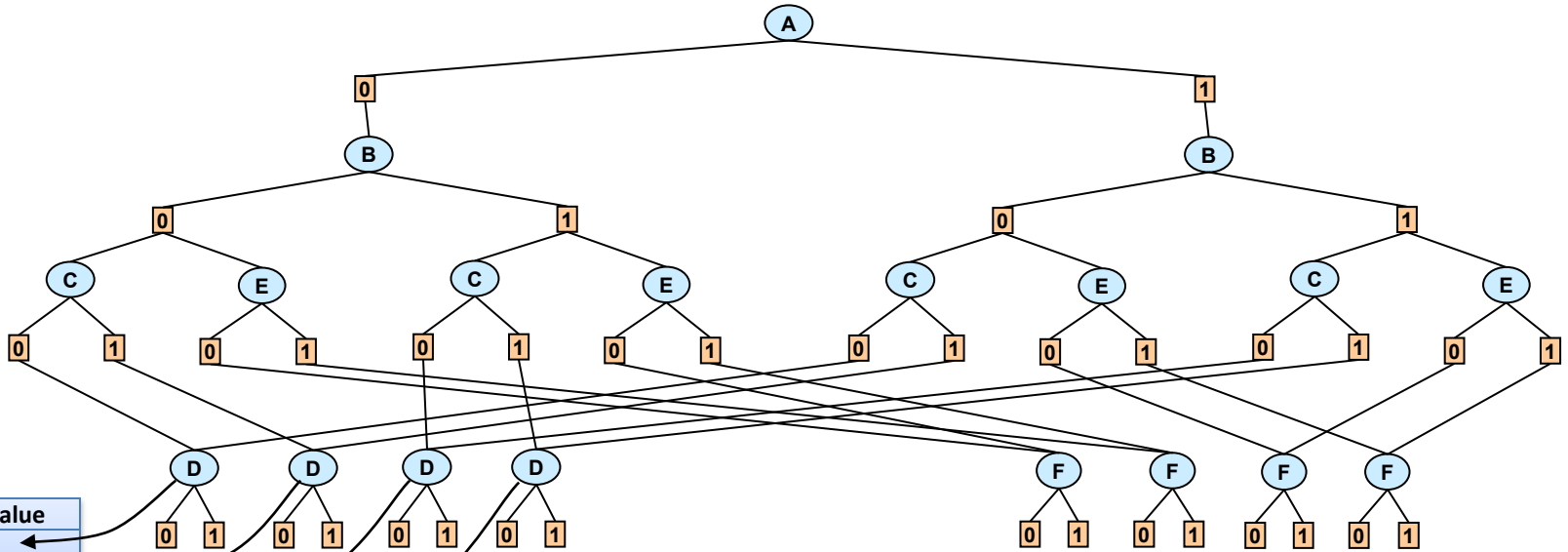
AND

OR

AND

OR

AND



B	C	Value
0	0	←
0	1	←
1	0	←
1	1	←

Context minimal AND/OR search graph

# Basic Heuristic Search

We assume min-sum problems in the following

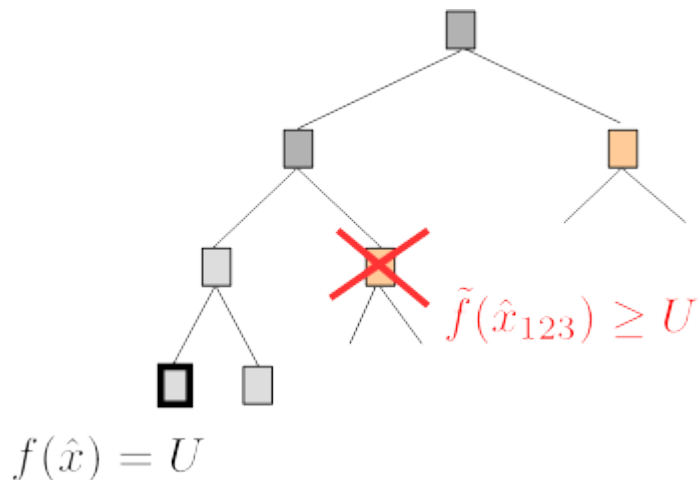
Heuristic function  $\tilde{f}(\hat{x}_p)$  computes a lower bound on the best extension of partial configuration  $\hat{x}_p$  and can be used to guide heuristic search.

We focus on:

## 1. Branch-and-Bound

Use heuristic function  $\tilde{f}(\hat{x}_p)$  to prune the depth-first search tree

Linear space

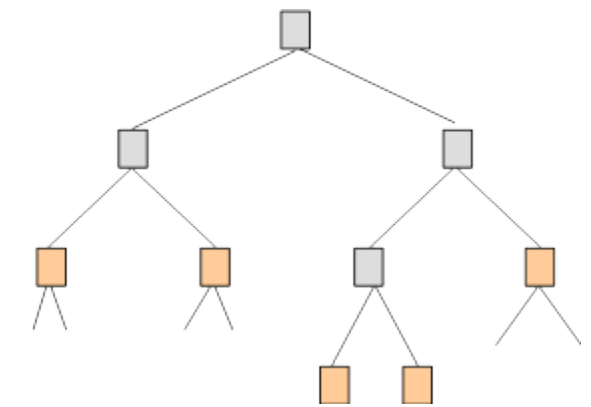
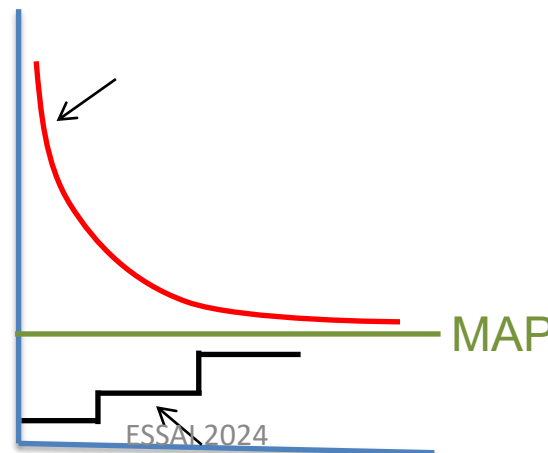


## 2. Best-First Search

Always expand the node with the lowest heuristic value  $\tilde{f}(\hat{x}_p)$

Needs lots of memory

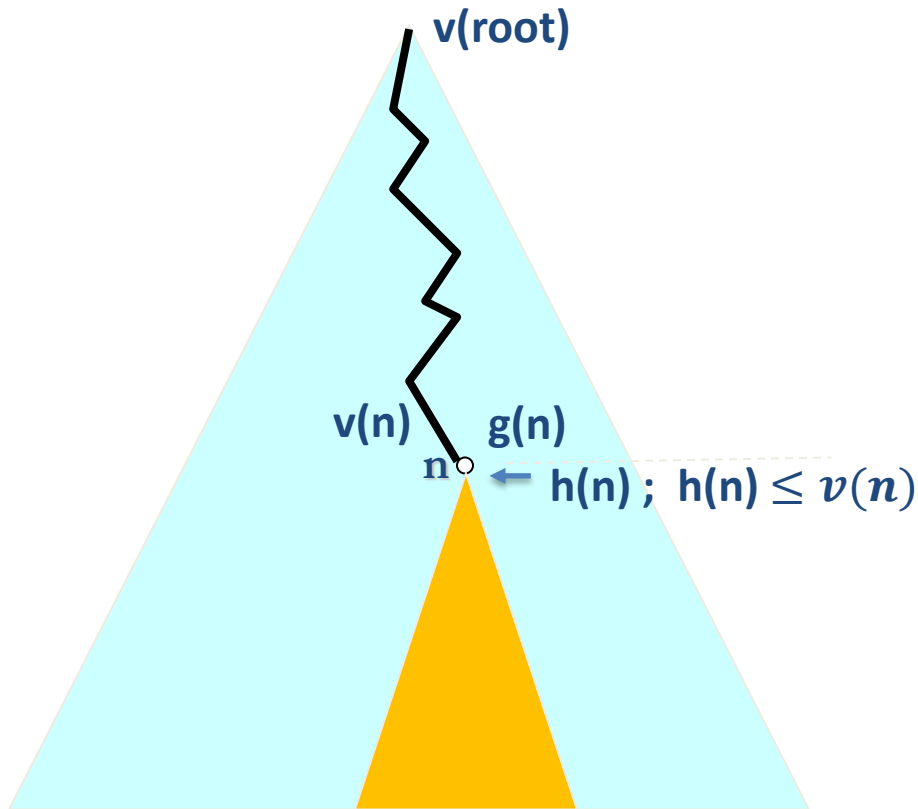
BnB is upper-bound anytime



# Basic Heuristic Search; Best-First

Task: compute  $v(\text{root})$ : MAP, Marginal, MMAP

Each node is a sub-problem  
(defined by current conditioning)

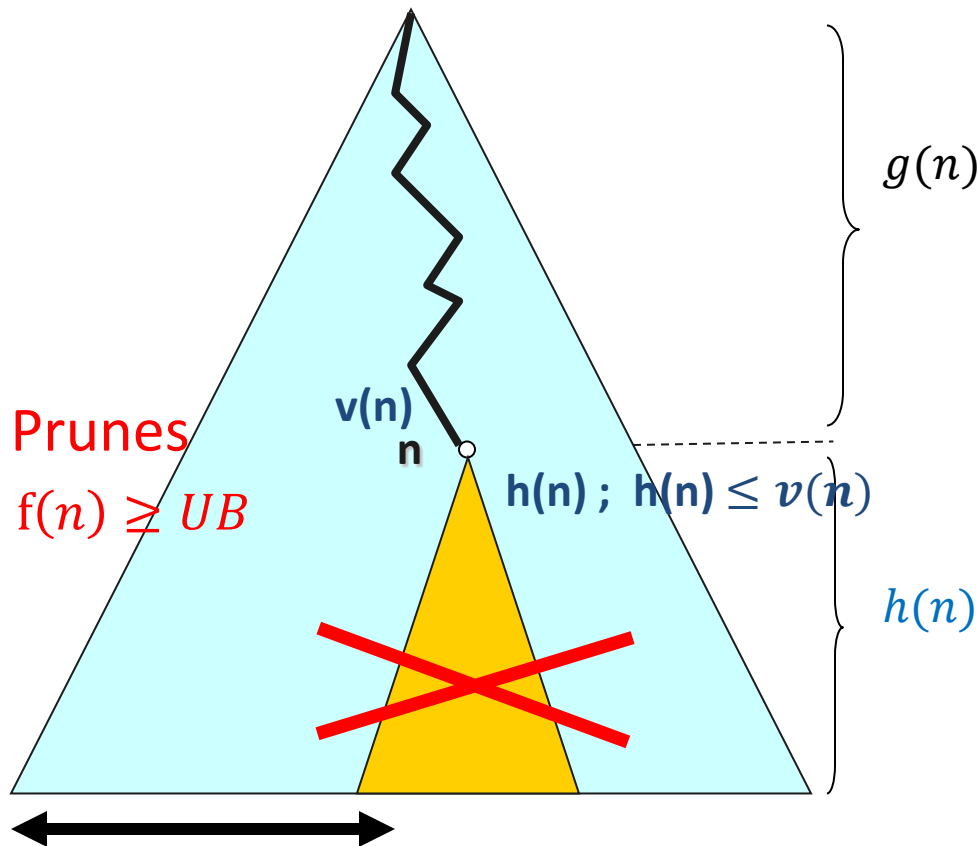


$$f(n) = g(n) + h(n) \leq g(n) + v(n) = f^*(n)$$

$f(n)$  is a lower bound on best cost through  $n$

- **Best-First Algorithms, (A\*)**
  - Expand nodes in OPEN list in order of  $\min f(n)$
  - Terminates with first full solution (for MAP)
- **Properties**
  - Optimal, if  $h(n) \leq v(n)$
  - Expands least set of nodes
  - exponential memory
  - **Not anytime solution for MAP**
  - **Yields lower bounds on value, anytime**

# Basic Heuristic Search; Depth-First



(UB) Upper Bound = best solution so far

- **Depth-First (B&B for MAP)**
  - Expand in dfs order
  - Update UB with each solution
  - Prunes if  $f(n) \geq UB$
- **Properties**
  - Can use only linear memory
  - Yields upper bounds anytime



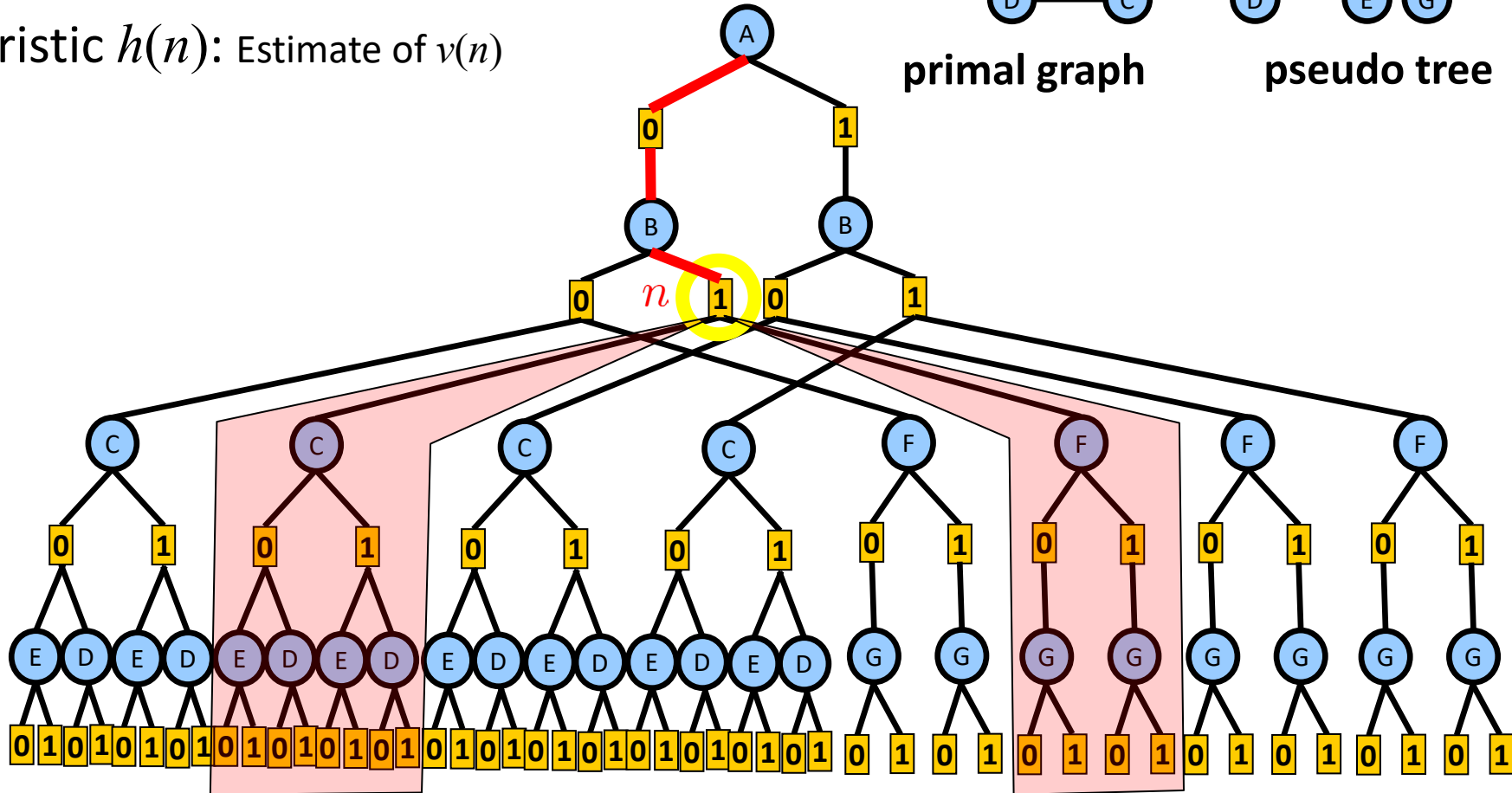
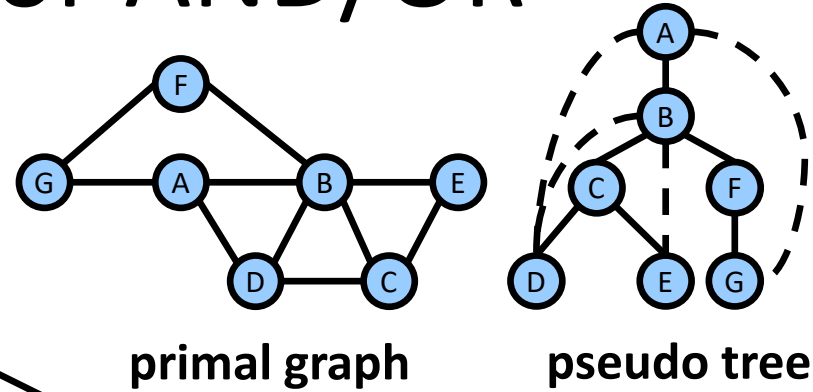
# Outline: Search



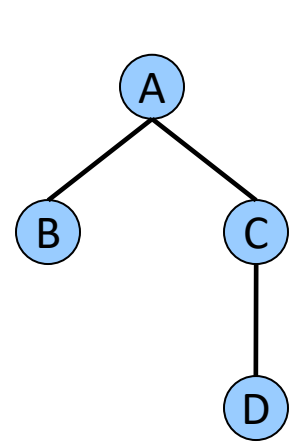
# Value and Heuristic for AND/OR

Value  $v(n)$ : answer of the subtree rooted at node  $n$

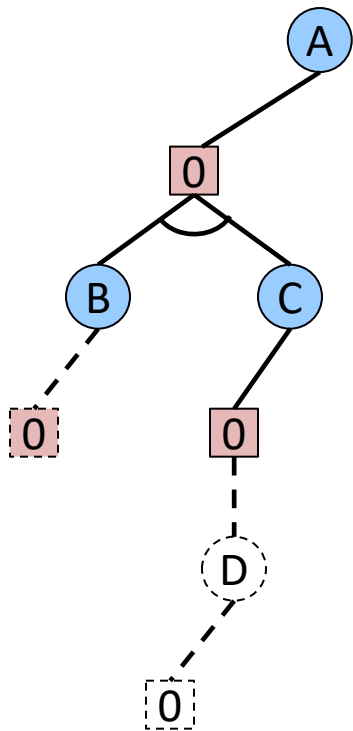
Heuristic  $h(n)$ : Estimate of  $v(n)$



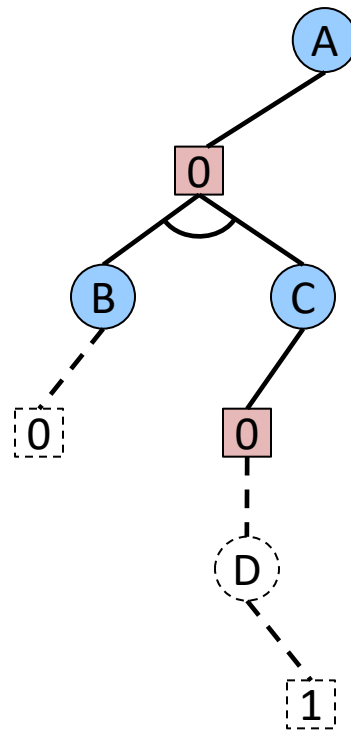
# Partial Solution Tree



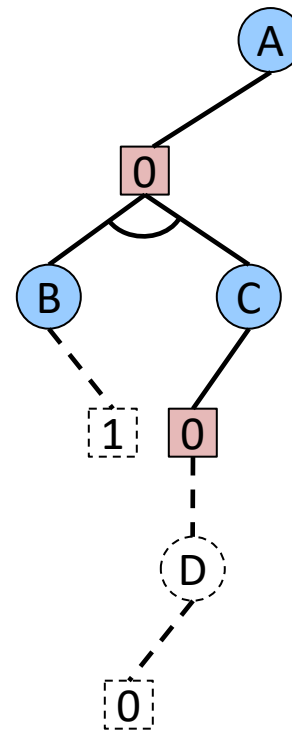
Pseudo tree



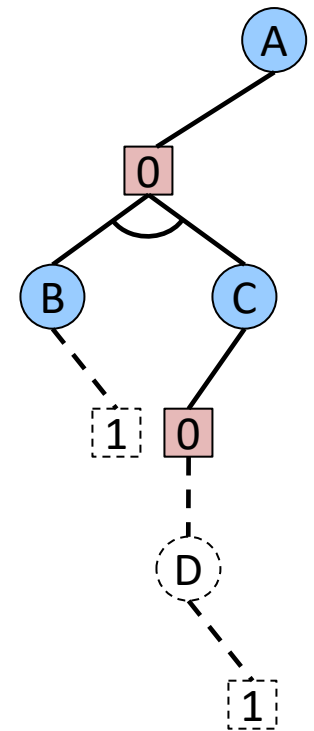
(A=0, B=0, C=0, D=0)



(A=0, B=0, C=0, D=1)



(A=0, B=1, C=0, D=0)



(A=0, B=1, C=0, D=1)

Extension( $T'$ ) – solution trees that extend  $T'$

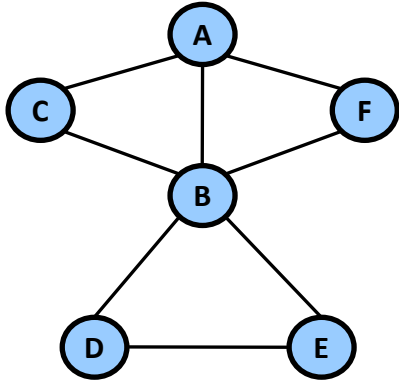
$g(T')$  = conditioned value of a node

$V(T')$  = the combined value below  $T'$

$f^*(T')$  = conditioned value through  $T'$



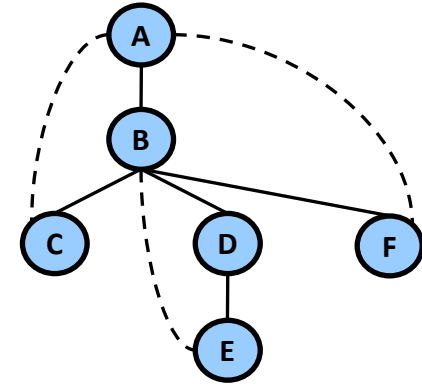
# Heuristic Evaluation Function



A	B	C	$f_1(ABC)$
0	0	0	2
0	0	1	5
0	1	0	3
0	1	1	5
1	0	0	9
1	0	1	3
1	1	0	7
1	1	1	2

A	B	F	$f_2(ABF)$
0	0	0	3
0	0	1	5
0	1	0	1
0	1	1	4
1	0	0	6
1	0	1	5
1	1	0	6
1	1	1	5

B	D	E	$f_3(BDE)$
0	0	0	6
0	0	1	4
0	1	0	8
0	1	1	5
1	0	0	9
1	0	1	3
1	1	0	7
1	1	1	4



OR

AND

OR

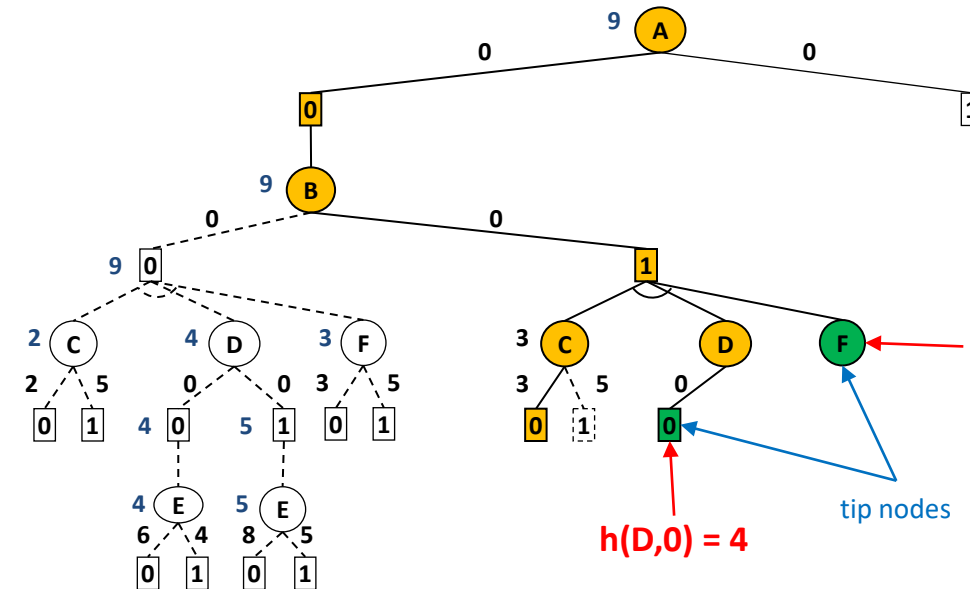
AND

OR

AND

OR

AND



$$h(n) \leq v(n)$$

$$h(F) = 5$$

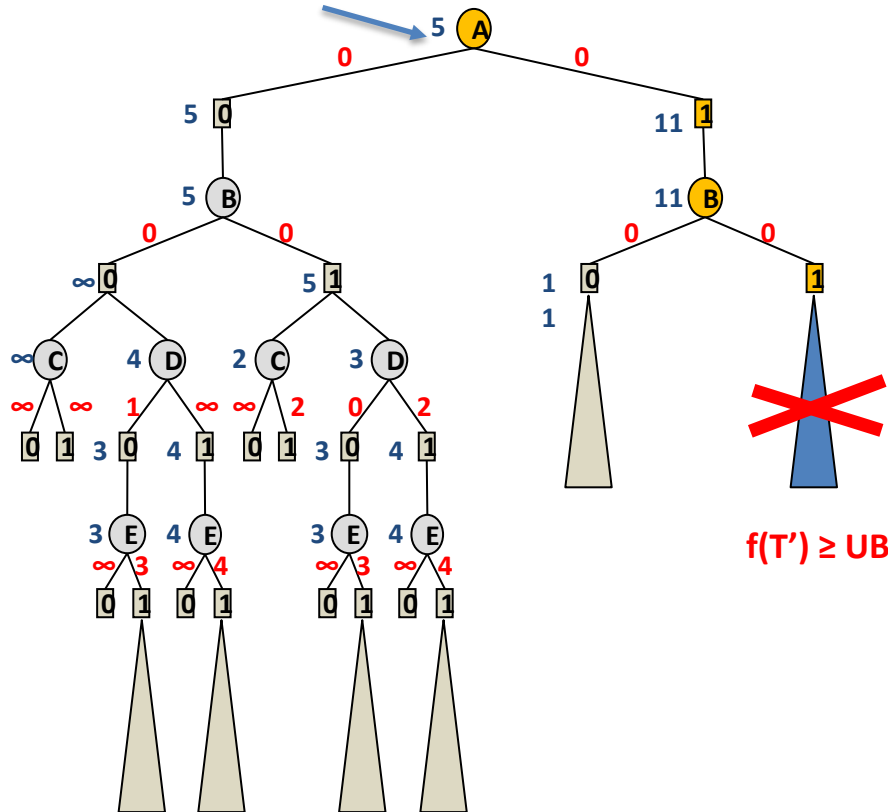
$$h(D,0) = 4$$

tip nodes

$$f(T') = w(A,0) + w(B,1) + w(C,0) + w(D,0) + h(D,0) + h(F) = 12 \leq f^*(T')$$

# Depth-First AND/OR Branch-and-Bound

UB (best solution so far)



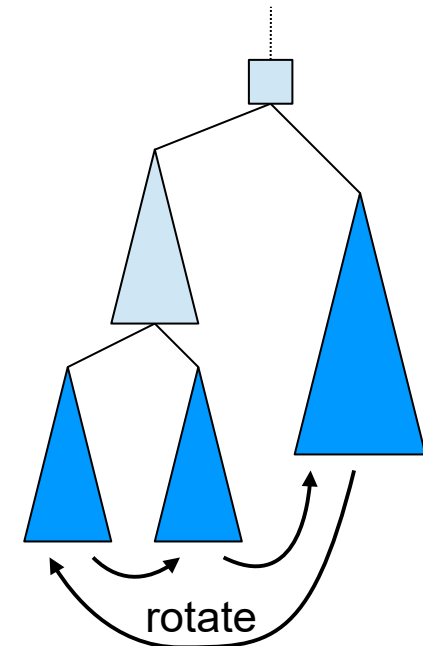
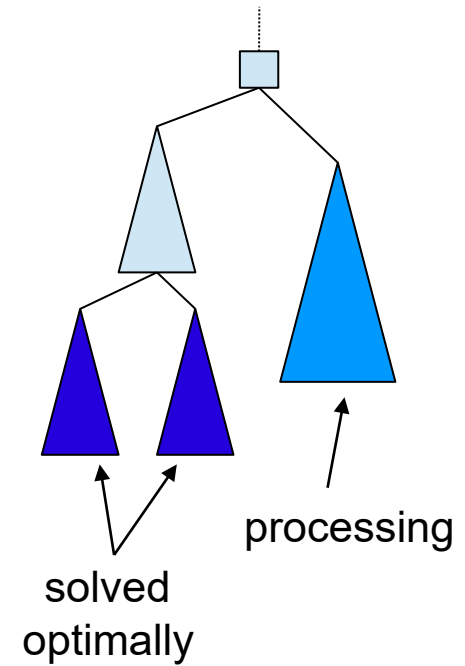
- Associate each node  $n$  with a heuristic lower bound  $h(n)$  on  $v(n)$

## Algorithm AOBB:

- **EXPAND** (top-down)
  - Evaluate  $f(T')$  and prune search if  $f(T') \geq UB$
  - If not in cache, generate successors of the tip node  $n$
- **PROPAGATE** (bottom-up)
  - Update value of the parent  $p$  of  $n$ 
    - OR nodes: **minimization**
    - AND nodes: **summation**
  - Cache value of  $n$  based on context

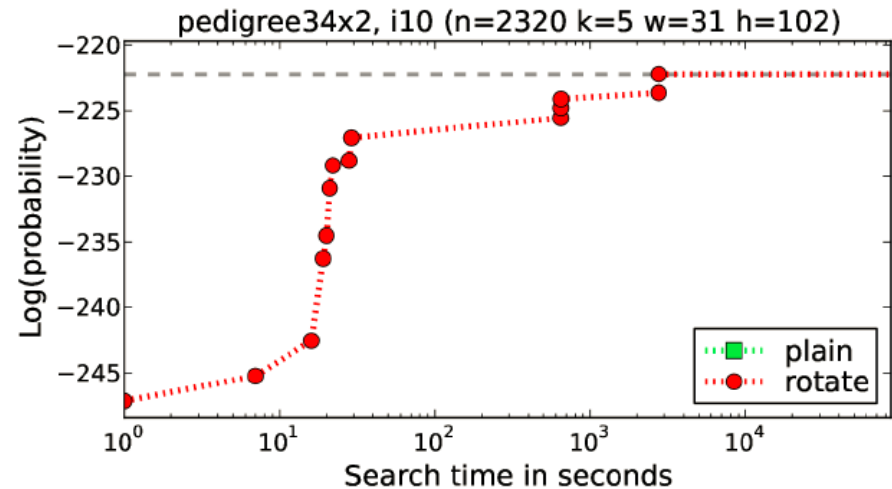
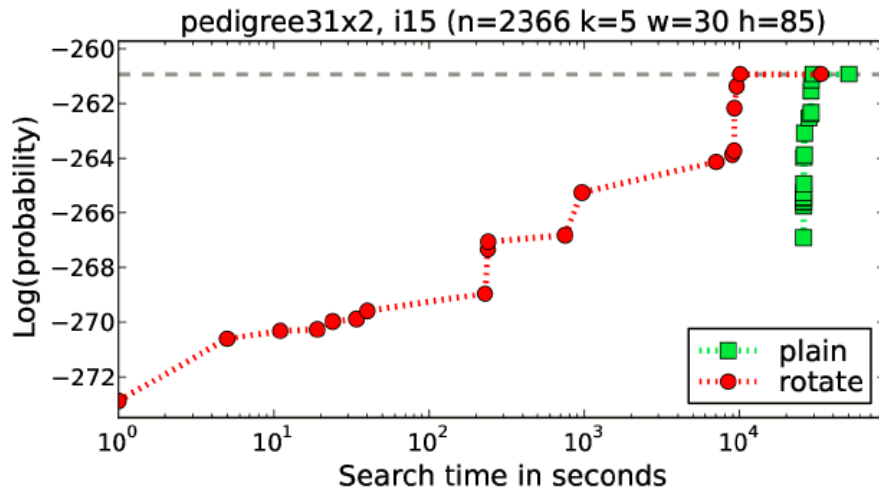
# Anytime Performance

- OR Branch-and-Bound is anytime
- But AND/OR breaks anytime behavior of depth-first scheme:
  - First anytime solution delayed until last sub-problem starts processing
- **Breadth-Rotating AOBB:**
  - Take turns processing sub-problems
    - Limit number of expansions per visit
  - Solve each sub-problem depth-first
    - Maintain favorable complexity bounds



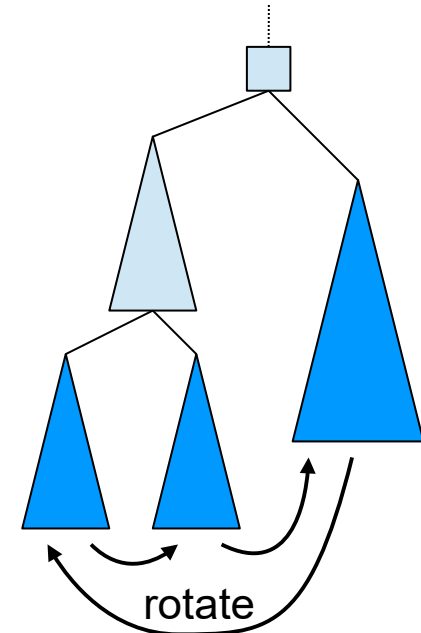
[Otten and Dechter, 2012]

# Anytime Performance



- **Breadth-Rotating AOBB:**

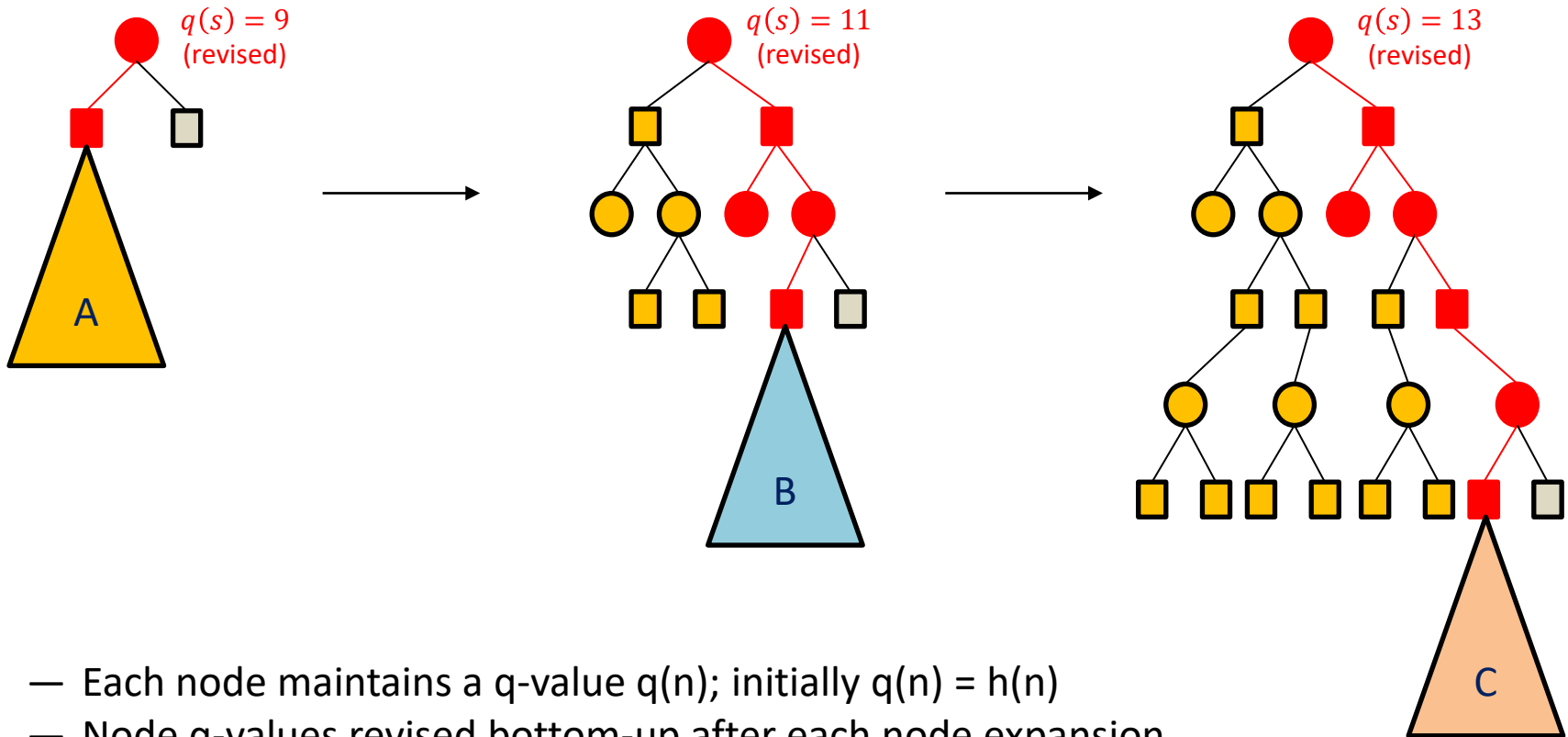
- Take turns processing sub-problems
  - Limit number of expansions per visit
- Solve each sub-problem depth-first
  - Maintain favorable complexity bounds



[Otten and Dechter, 2012]



# AOBF: Best-First AND/OR Search



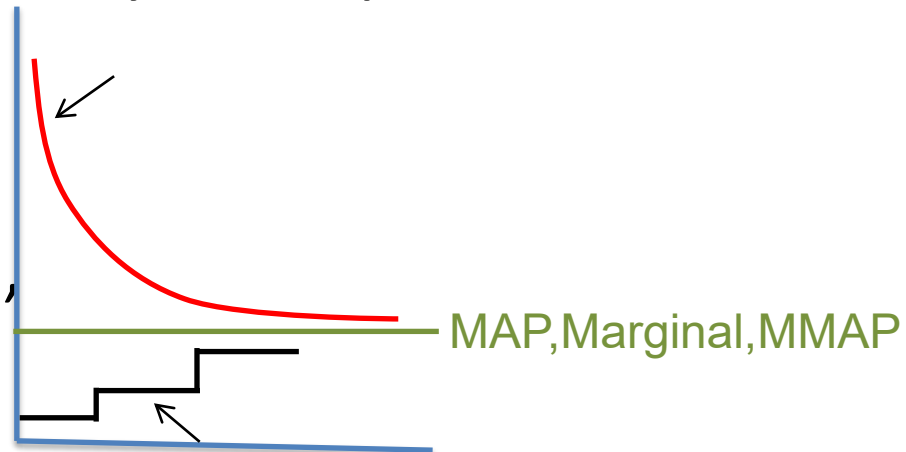
- Each node maintains a  $q$ -value  $q(n)$ ; initially  $q(n) = h(n)$
- Node  $q$ -values revised bottom-up after each node expansion
- Update current best partial solution subtree (a tip node expanded next)
- All expanded nodes are stored in memory
- Search terminates with optimal solution (cost)

# AOBF: Best-First AND/OR Search

- AO\*-traverses the context-minimal AND/OR graph
  - All nodes expanded are stored in memory
  - Each node maintains a q-value:  $q(n)$ , (Best lower bound below n)
- Node q-values are revised bottom-up after each expansion
  - OR: minimization:  $q(n) = \min_{n' \in \text{succ}(n)} (w(n, n') + q(n'))$
  - AND: summation:  $q(n) = \sum_{n' \in \text{succ}(n)} q(n')$ , (initially,  $q(n) = h(n)$ )

# AOBF versus AOBB

- **AOBF** expands a smallest subset of the AO search space
  - This translates into significant time savings
- **AOBB** can use far less memory by avoiding dead-caches, whereas **AOBF** keeps in memory the explicated search graph



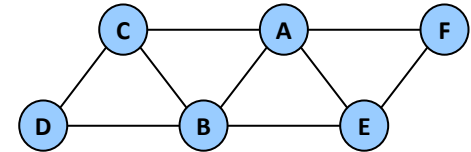
- **AOBB (BRAOBB)** is anytime,
- **AOBF** generates lower bounds anytime, but not anytime solutions (configuration)

# Outline: Search

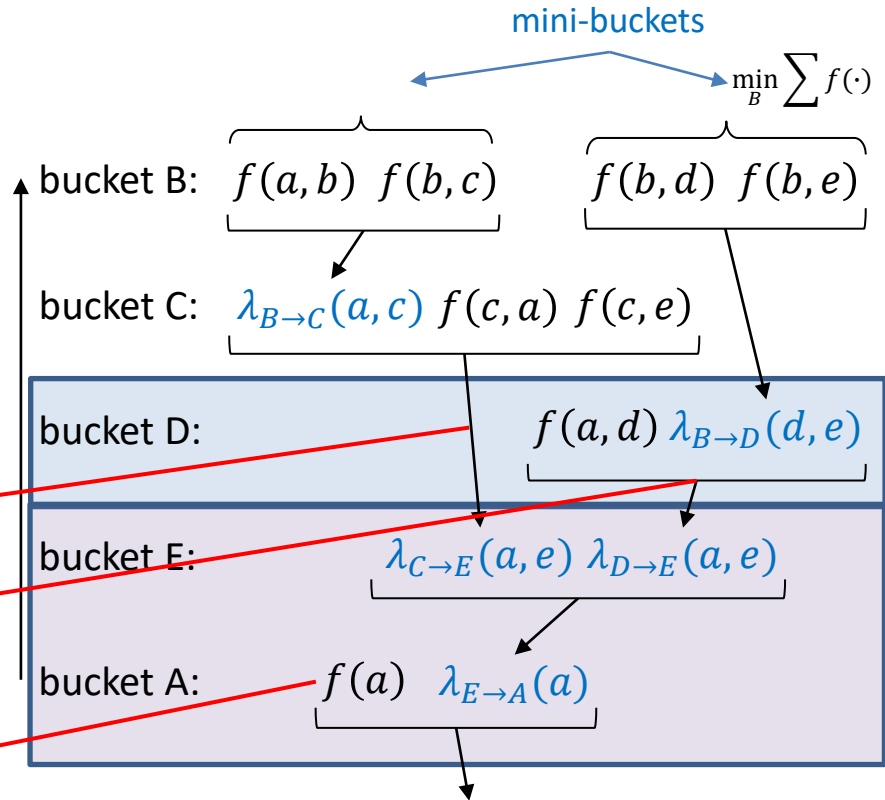
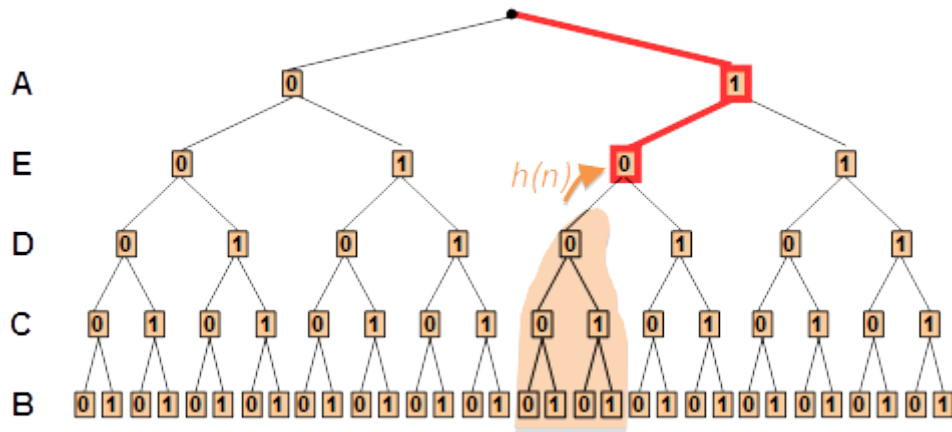




# Static Mini-Bucket Heuristics



Given a partial assignment,  $[\hat{a} = 1, \hat{e} = 0]$   
 (weighted) mini-bucket gives an admissible heuristic:



cost to go:

$$h(\hat{a}, \hat{e}, D) = \lambda_{C \rightarrow E}(\hat{a}, \hat{e}) + f(\hat{a}, D) + \lambda_{B \rightarrow D}(D, \hat{e})$$

(admissible:  $h(\hat{a}, \hat{e}, D) \leq h^*(\hat{a}, \hat{e}, D)$ )

cost so far:

$$g(\hat{a}, \hat{e}, D) = f(A = \hat{a})$$

**L = lower bound**

# Properties of the MBE Heuristics

- MBE heuristic is monotone, admissible
- Computed in linear time (during search)
- Important:
  - Heuristic strength can vary by  $MBE(i)$
  - Higher  $i$ -bound  $\rightarrow$  more pre-processing  $\rightarrow$  more accurate heuristic  $\rightarrow$  less search
- Allows controlled trade-off between pre-processing and search
- Can be computed **statically** or **dynamically** during search

# Review: Weighted Mini-bucket

[Liu & Ihler 2011]

## For Sum-Inference

$$\lambda_{B \rightarrow C} = \sum_b^{w_{B1}} f(a, b) \cdot f(b, c)$$

$$w_{B1} + w_{B2} = 1$$

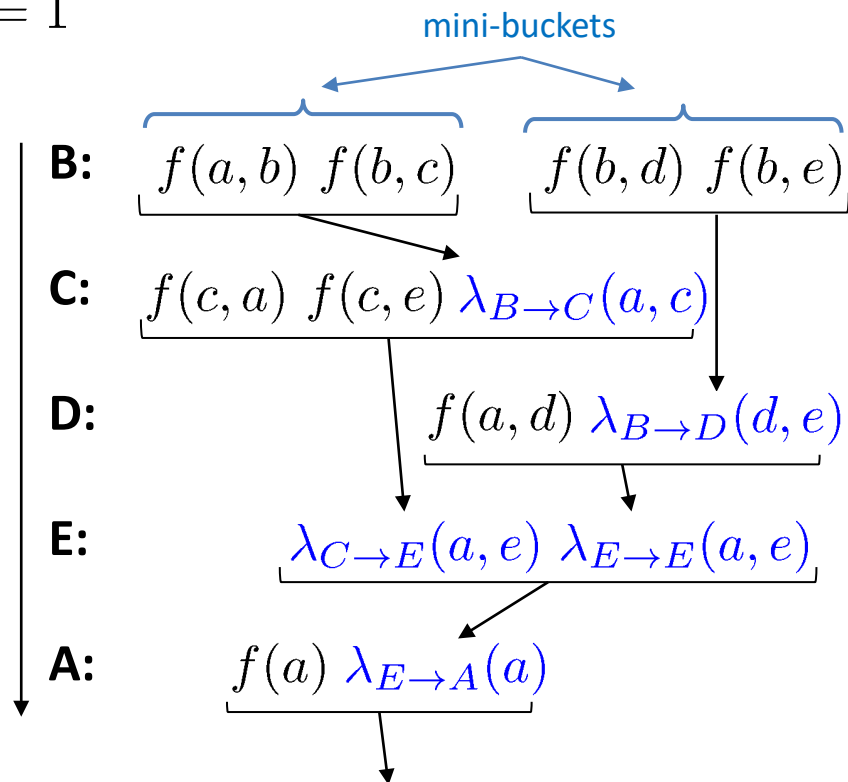
$$\lambda_{B \rightarrow D} = \sum_b^{w_{B2}} f(b, d) \cdot f(b, e)$$

$$\lambda_{C \rightarrow E} = \sum_c f(c, a) \cdot f(c, e) \cdot \lambda_{B \rightarrow C}$$

⋮

Compute downward messages  
using weighted sum

Upper bound if all weights positive  
(corresponding lower bound if only one positive, rest negative)



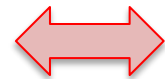
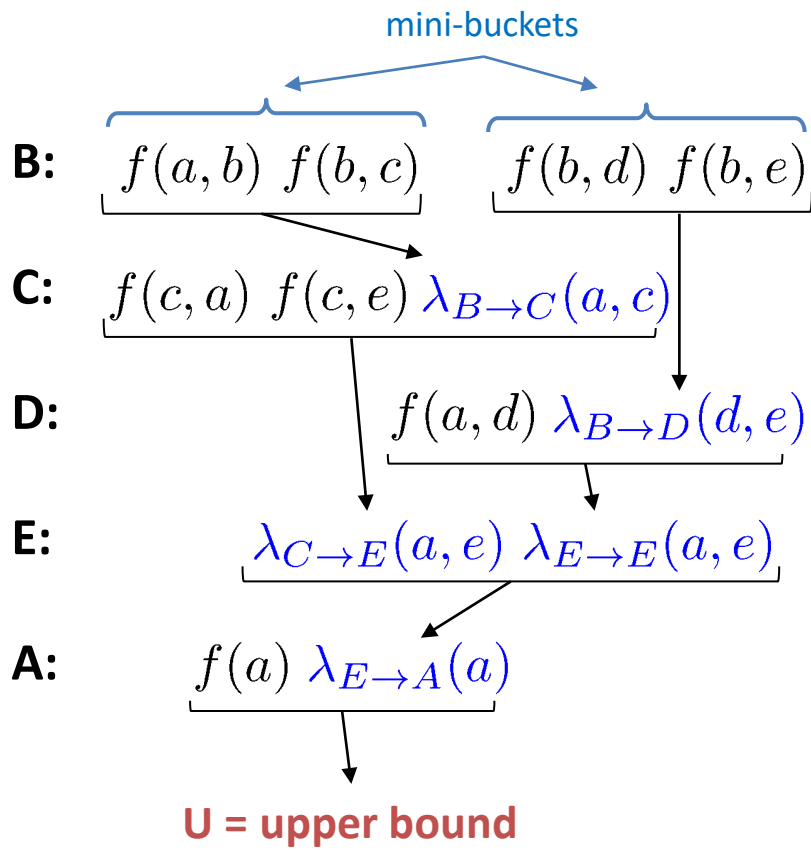
**U = upper bound**



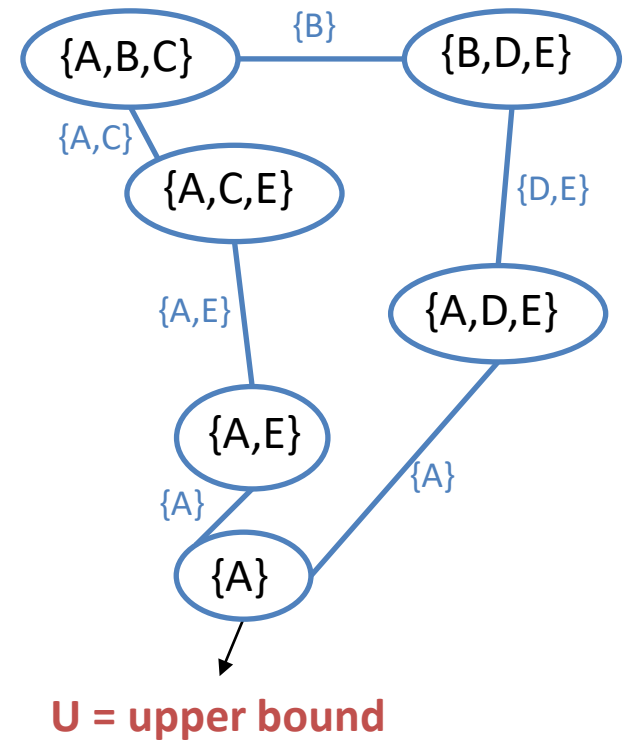
# Review: MBE+Moment Matching

For all queries

- Mini-bucket elimination defines regions with bounded complexity

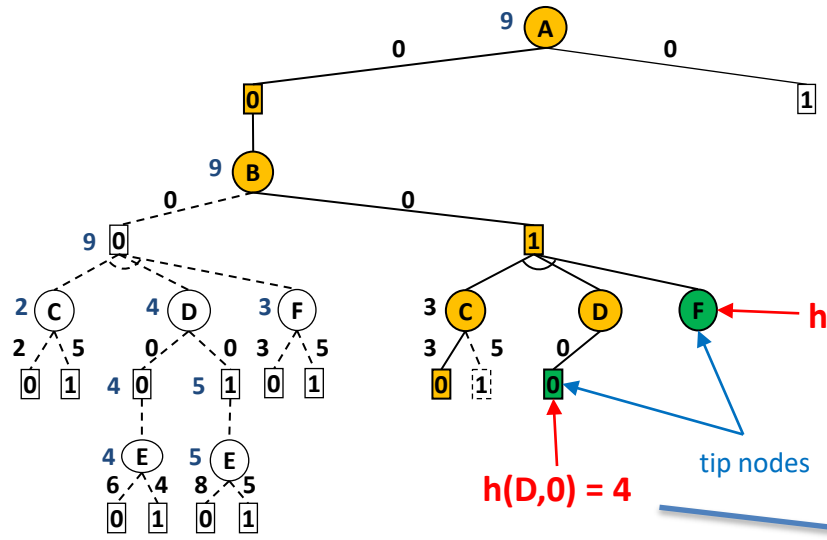


Join graph:

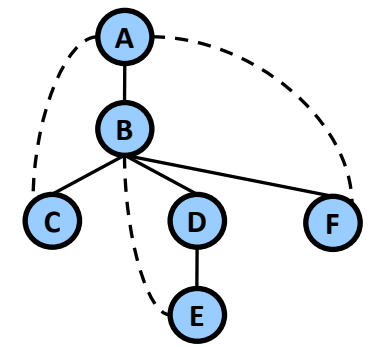


# MBE Heuristic Guides AO Search

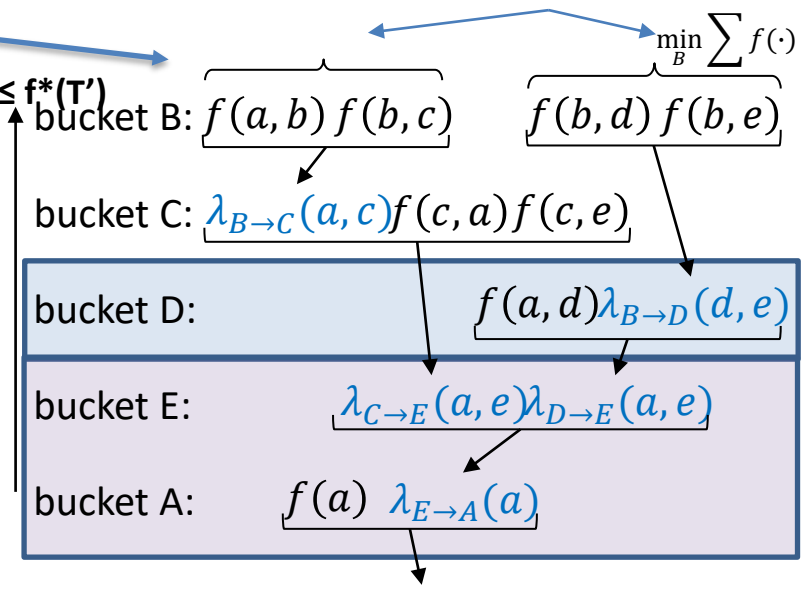
OR  
AND  
OR  
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OR  
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OR  
AND



$$h(n) \leq v(n)$$

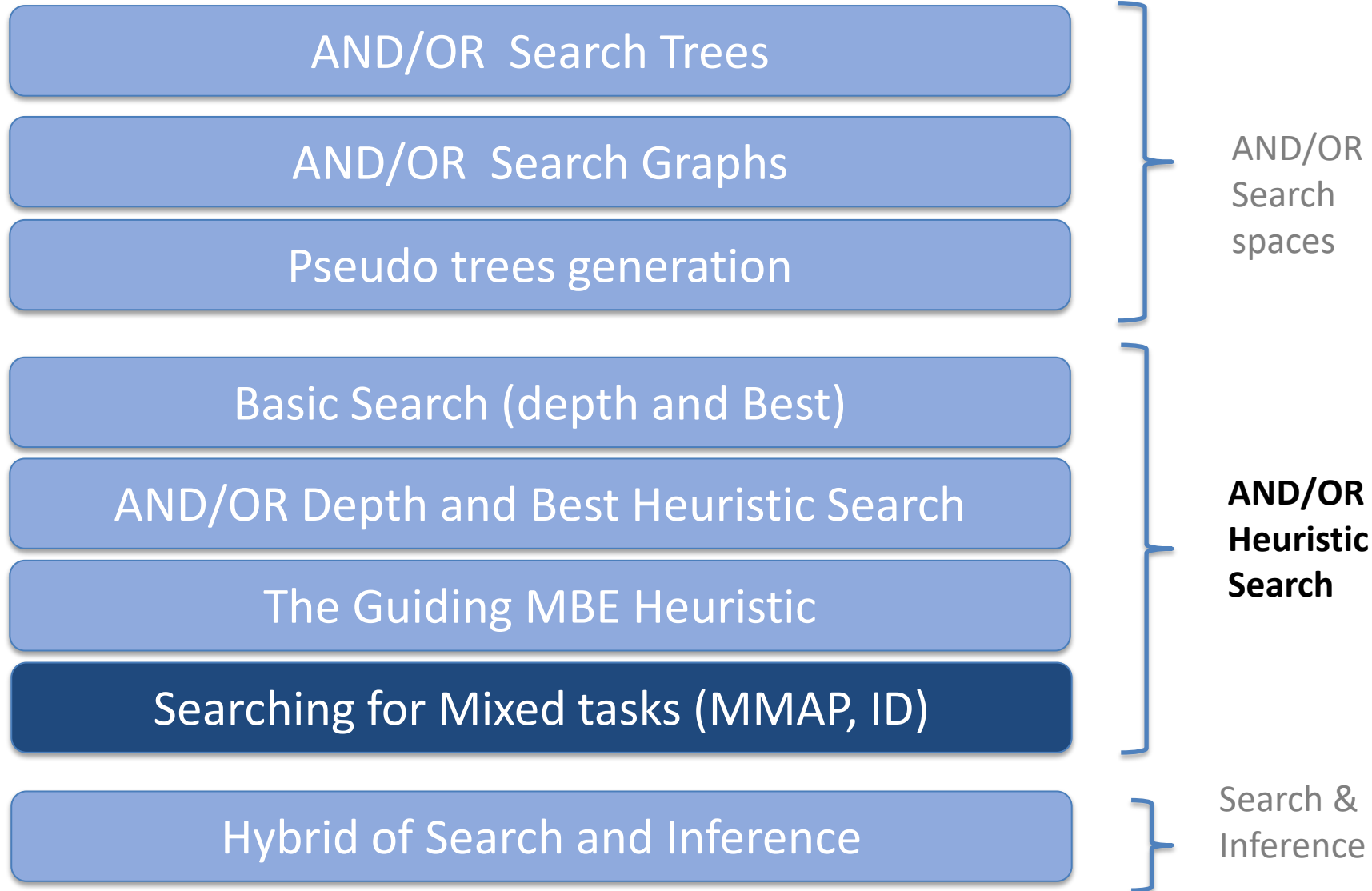


$$f(T') = w(A,0) + w(B,1) + w(C,0) + w(D,0) + h(D,0) + h(F) = 12 \leq f^*(T')$$



**L = lower bound**

# Outline: Search

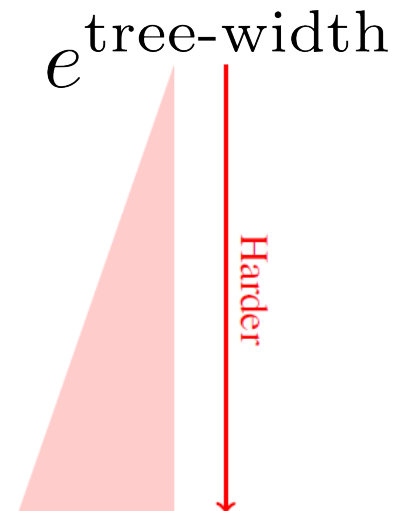


# Probabilistic Reasoning Problems

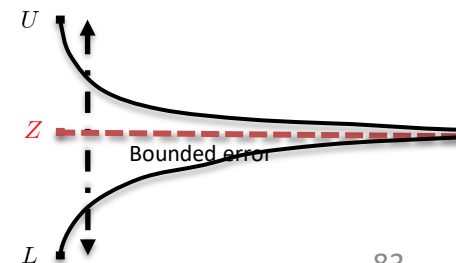
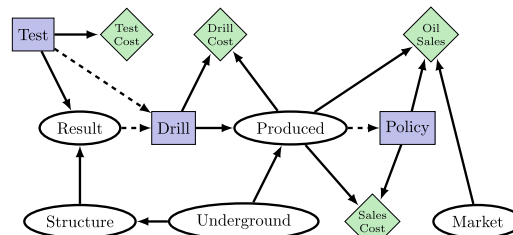
- Exact Inference by elimination or search
- Complexity:

Causal effects

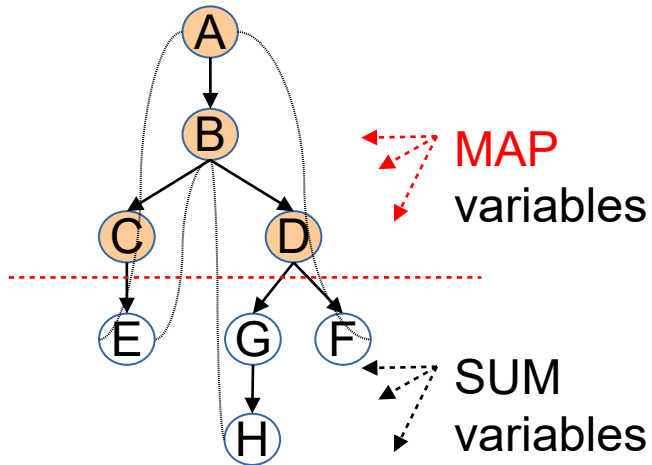
Max-Inference:	$f(x^*) = \max_x \prod_{\alpha} f_{\alpha}(x_{\alpha})$
Sum-Inference:	$Z = \sum_x \prod_{\alpha} f_{\alpha}(x_{\alpha})$
Mixed-Inference (MMAP):	$f_M(x_M^*) = \max_{x_M} \sum_{x_S} \prod_{\alpha} f_{\alpha}(x_{\alpha})$
Mixed-Inference (MEU):	$\text{MEU} = \max_{D_1, \dots, D_m} \sum_{X_1, \dots, X_n} \left( \prod_{P_i \in P} P_i \right) \times \left( \sum_{r_i \in R} r_i \right)$



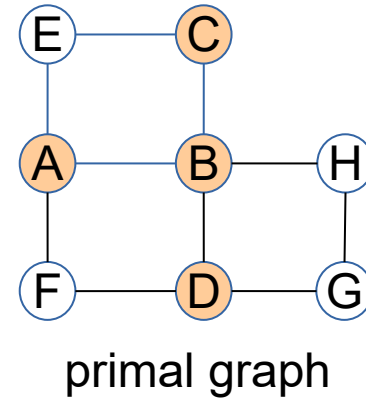
Influence diagrams & planning



# AND/OR Search for Marginal MAP

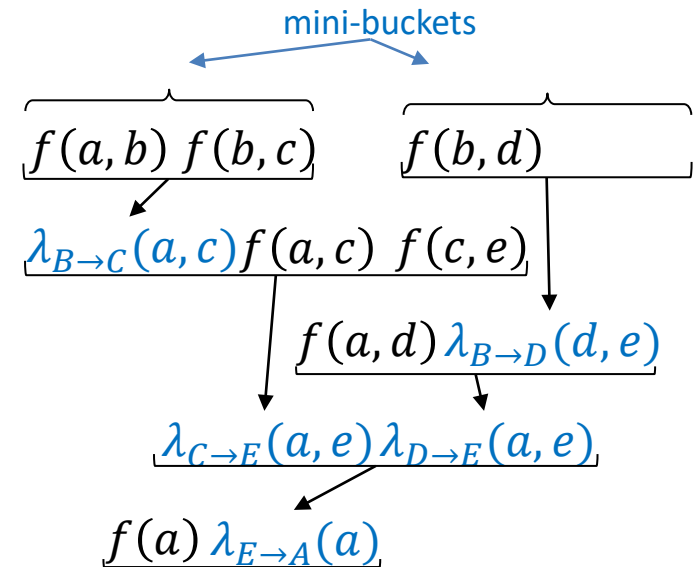
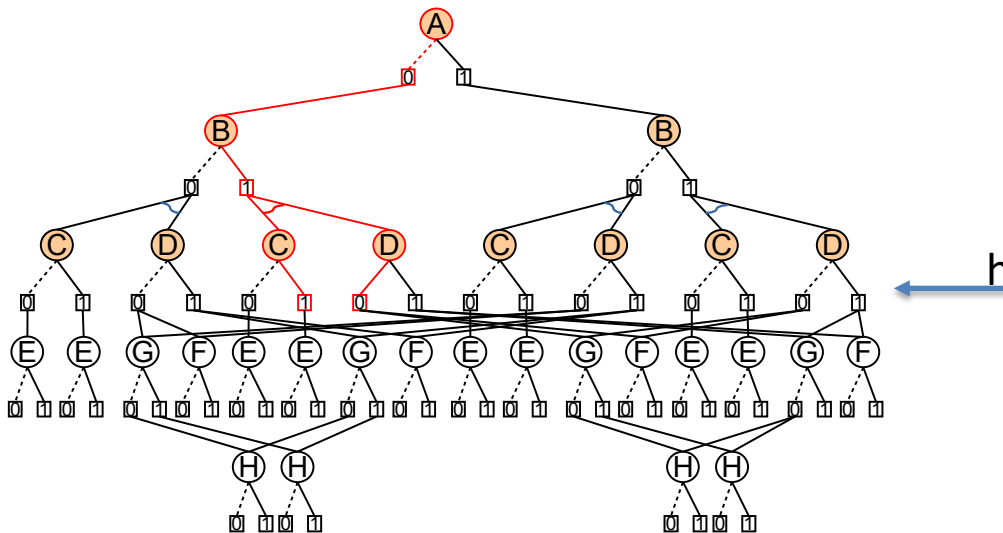


constrained pseudo tree



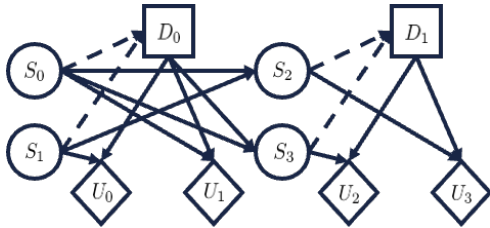
primal graph

**Anytime Depth+Best to yield upper and lower bounds**

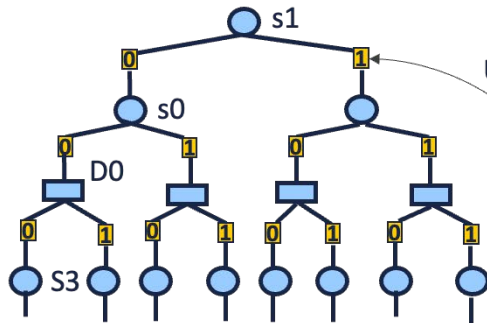


# AND/OR Search for Influence Diagrams

Heuristic AND/OR search with decomposition bounds

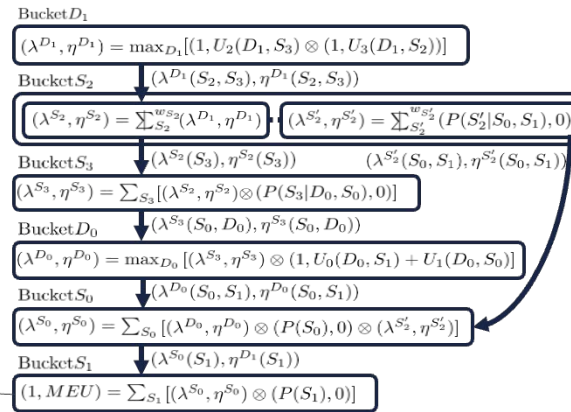


$$\sum_{S_0, S_1} \max_{D_0} \sum_{S_2, S_3} \max_{D_1} [\prod_{P_i \in \mathbf{P}} P_i] [\sum_{U_i \in \mathbf{U}} U_i] [\prod_{\Delta_i \in \mathbf{\Delta}} \Delta_i]$$



[Marinescu 2010]

AND/OR Search Graph for Influence Diagrams

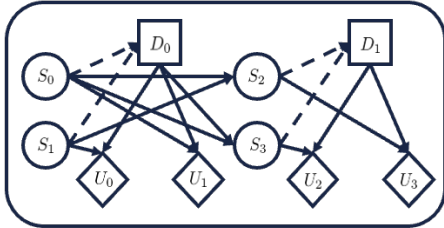


[Lee et.al 2019]

Weighted Mini-bucket elimination bound for Influence Diagrams

# AND/OR Search for Influence Diagrams

Influence Diagram [Howard and Matheson 1981]



[Jensen et al 1994; Maua et. al 2012]

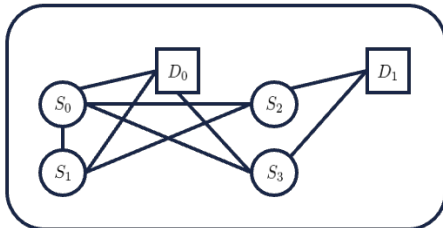
Valuation Algebra for Influence Diagrams

$$\Psi := \{(P_i, 0) \mid P_i \in \mathbf{P}\} \cup \{(1, U_i) \mid U_i \in \mathbf{U}\}$$

$$\Psi_1 \otimes \Psi_2 := (P_1 P_2, P_1 V_2 + P_2 V_1)$$

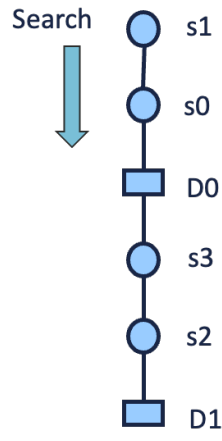
$$\sum_Y^w \Psi := (\sum_Y^w P, \sum_Y^w V)$$

Primal Graph



[Freuder and Quinn 1985]

Pseudo-tree

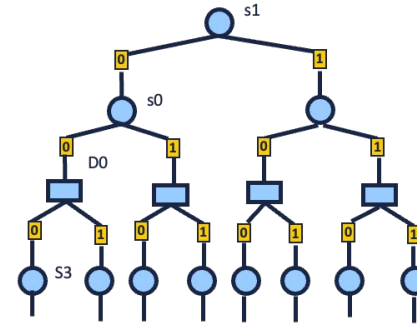


Inference [Dechter 1999]

By bucket elimination

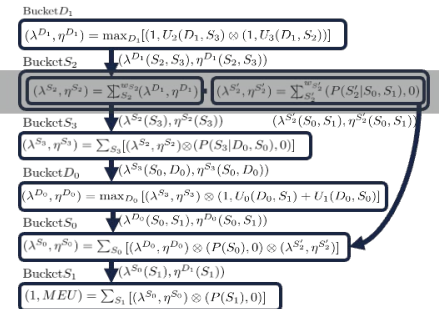
[Dechter et al 2007; Marinescu, et al 2010]

AND/OR Search Space



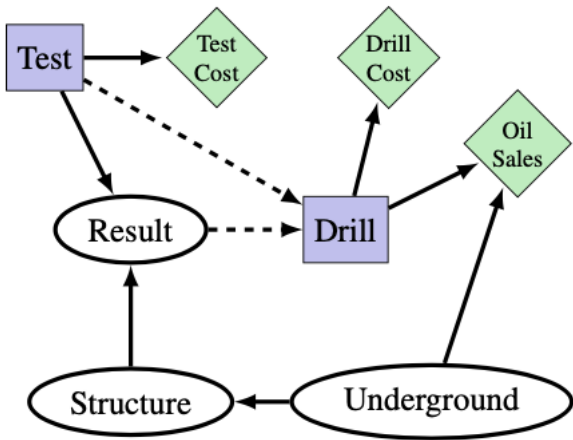
[Dechter et al 2003; Liu et al 2011; Lee et al 2019]

Decomposition bounds from bucket tree



# AND/OR Search for Influence Diagrams

- Use perfect recall order to enumerate search tree



$p(U)$	$U = 0$	$U = 1$	$U = 2$
	0.5	0.3	0.2

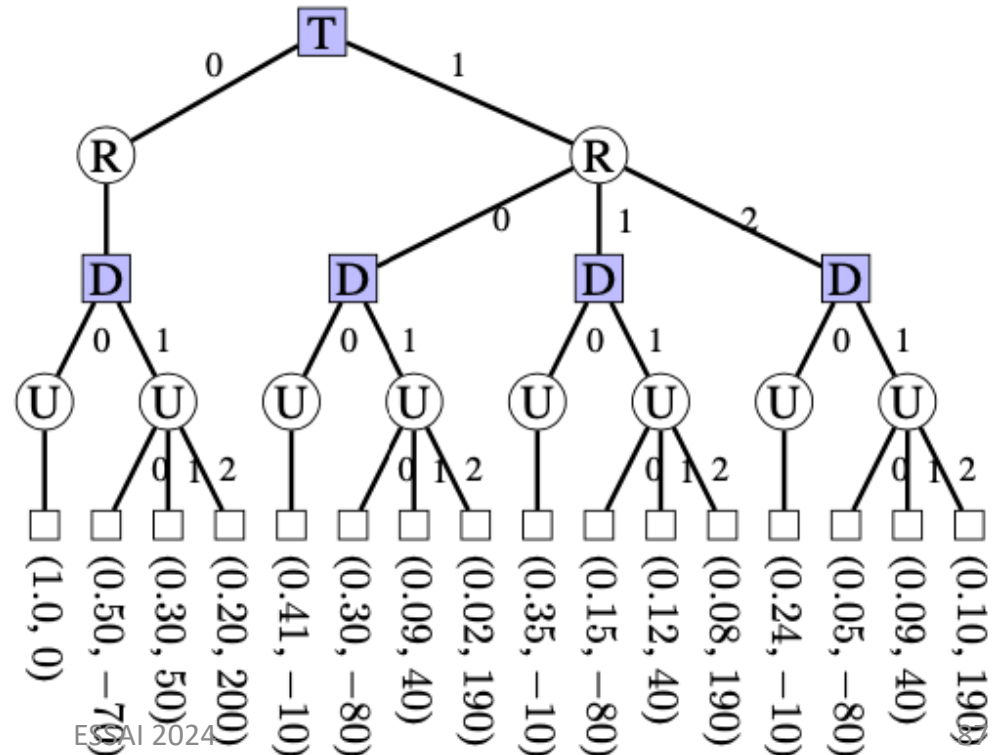
$p(S U)$	$S = 0$	$S = 1$	$S = 2$
$U = 0$	0.6	0.3	0.1
$U = 1$	0.3	0.4	0.3
$U = 2$	0.1	0.4	0.5

$$p(R|S, T) = \mathbb{1}[R = ST]$$

$u_1(T)$	$T = 0$	$T = 1$
	\$0	-\$10k

$u_2(D)$	$D = 0$	$D = 1$
	\$0	-\$70k

$u_3(D, U)$	
$D = 0$	\$0
$DU = 10$	\$0
$DU = 11$	\$120k
$DU = 12$	\$270k



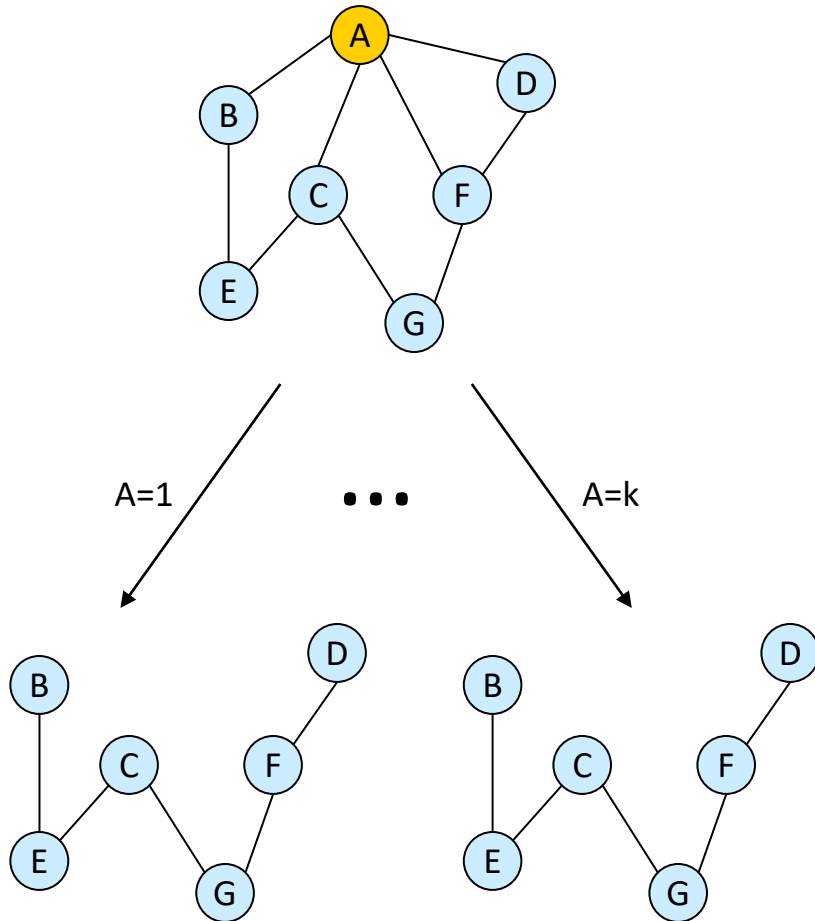


# Outline: Search



# Conditioning versus Elimination

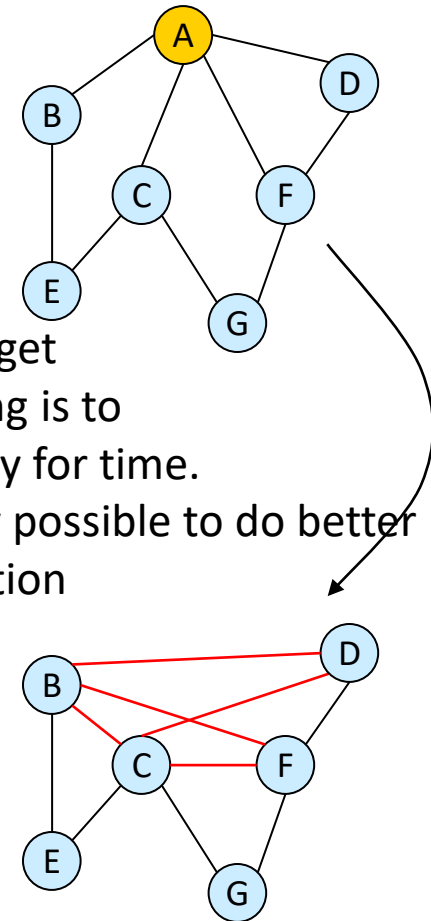
Conditioning (search)



k "sparser" problems

Elimination (inference)

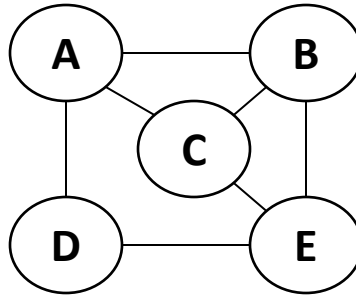
The main target  
In conditioning is to  
Trade memory for time.  
It is not really possible to do better  
Then elimination



1 "denser" problem

# Hybrid: Cutset-Conditioning

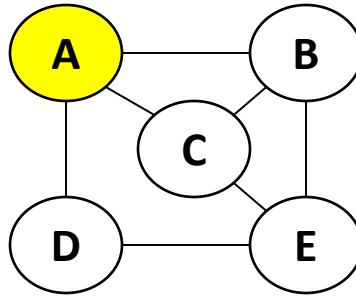
## Variable Branching by Conditioning



# Hybrid: Cutset-Conditioning

## Variable Branching by Conditioning

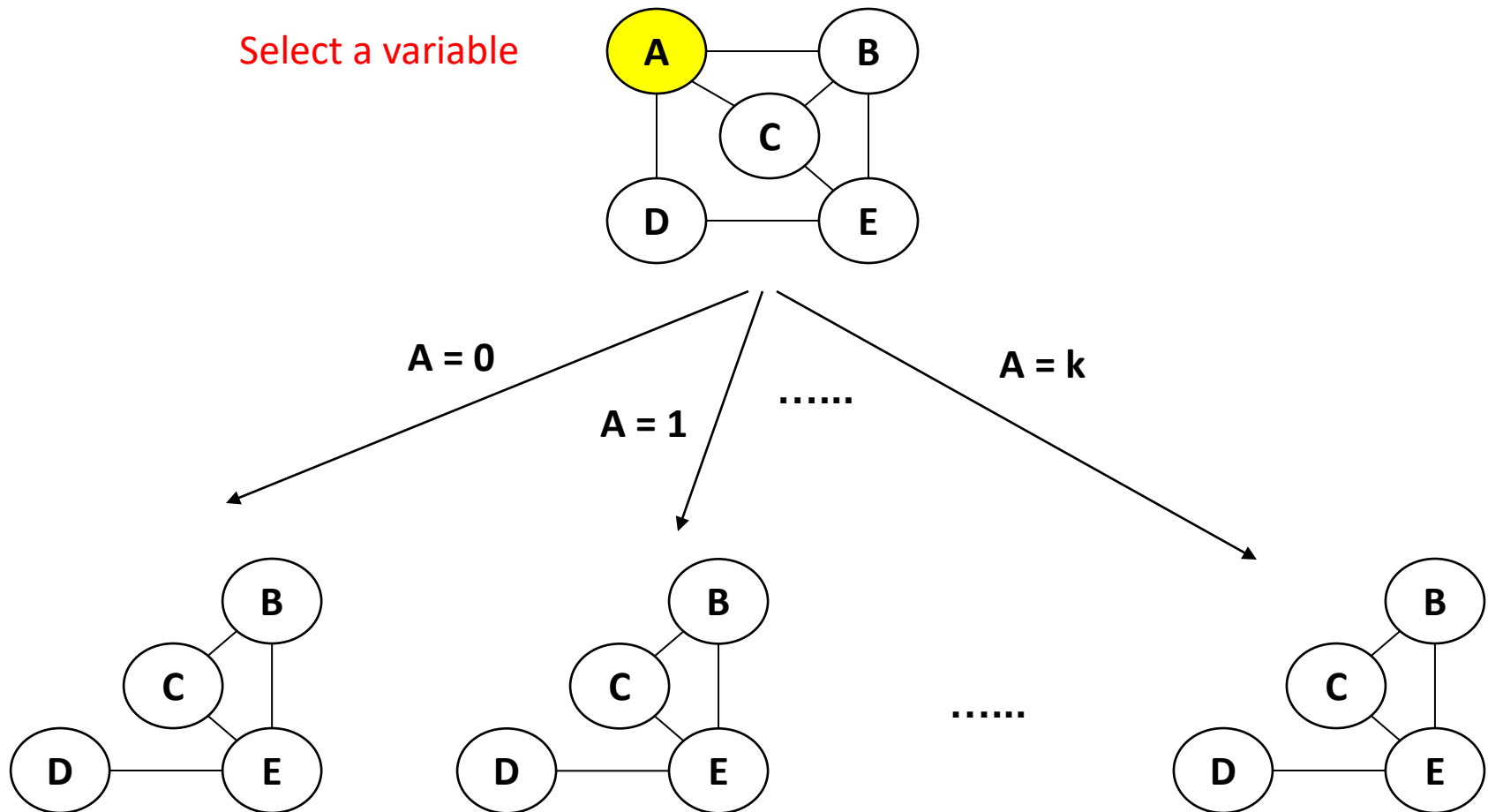
Select a variable



# Hybrid: Cutset-Conditioning

## Variable Branching by Conditioning

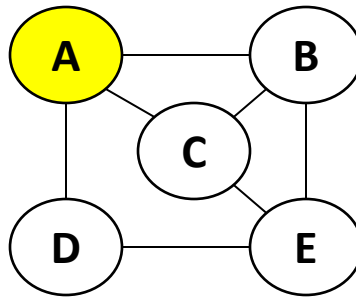
Select a variable



# Hybrid: Cutset-Conditioning

## Variable Branching by Conditioning

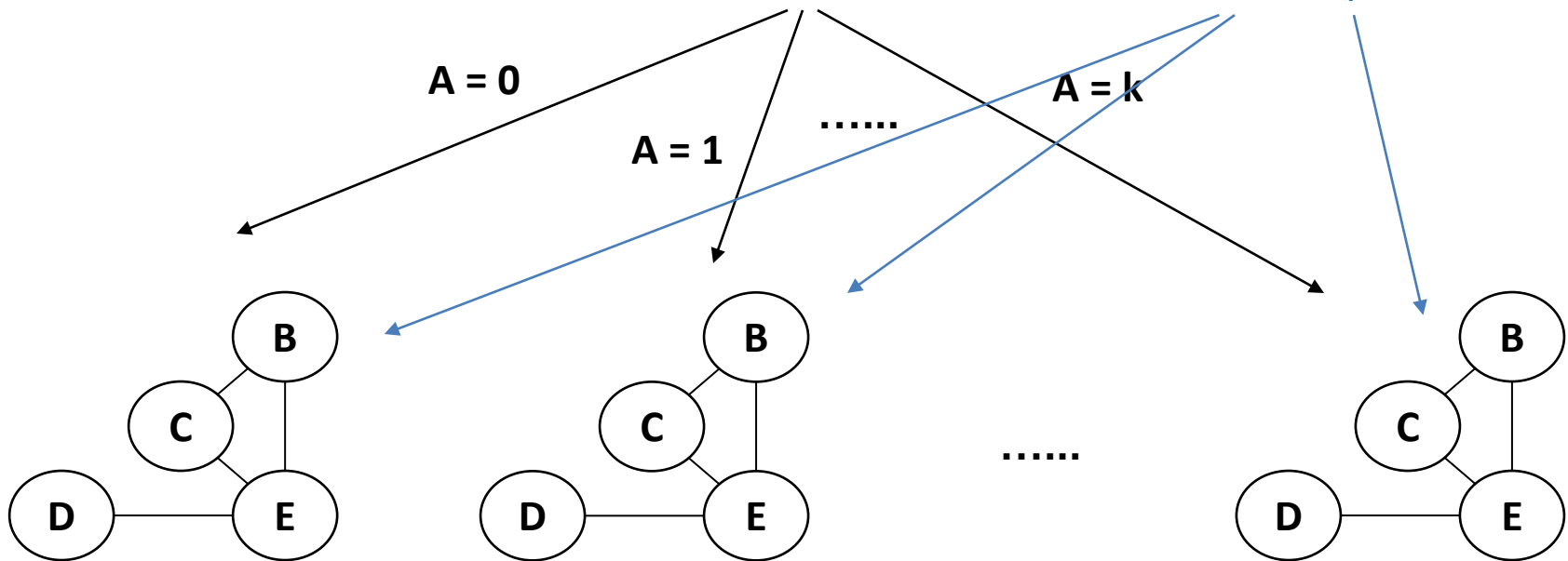
Select a variable



General principle:

Condition until tractable

Solve each sub-problem efficiently

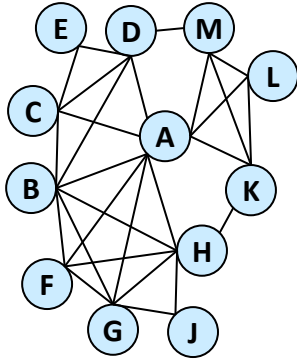


# Hybrids Variants

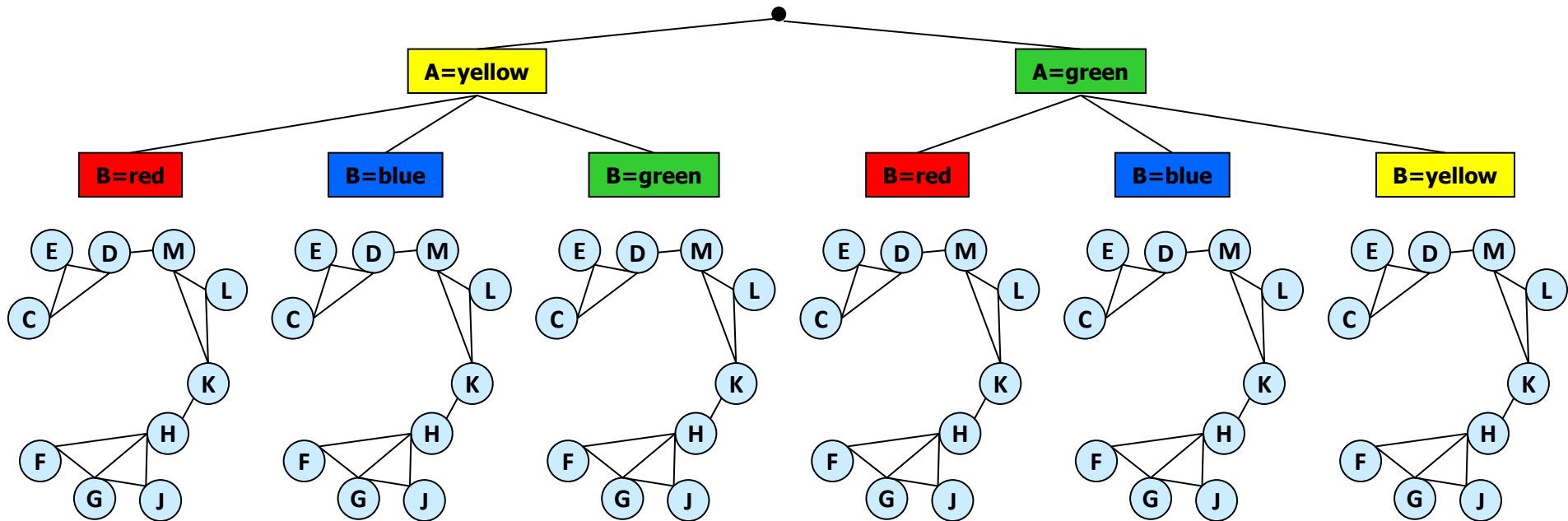
- **Condition, condition, condition**, ... and then only eliminate (w-cutset, cycle-cutset VEC(i))
- **Eliminate, eliminate, eliminate**, ... and then only search
- **Alternate** conditioning and elimination steps (elim-cond(i), ALT-VEC(i))

# OR w-Cutset

Graph  
Coloring  
problem

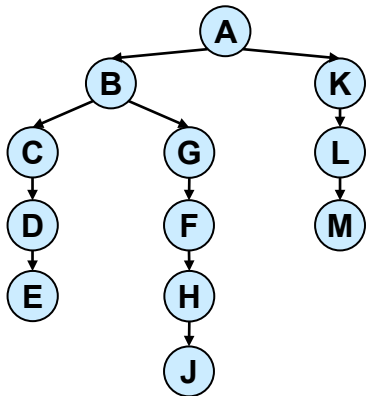
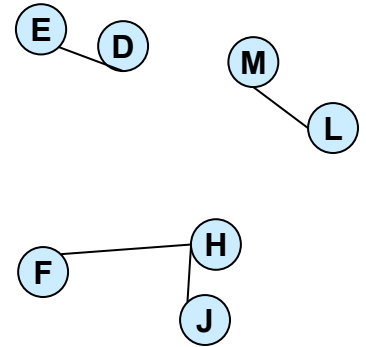
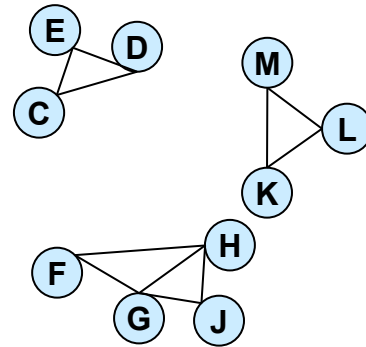
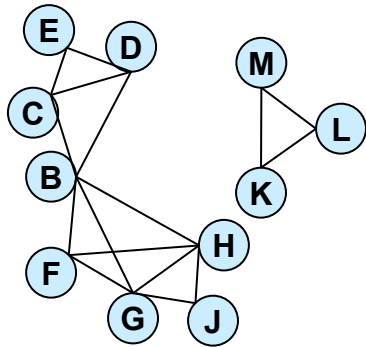
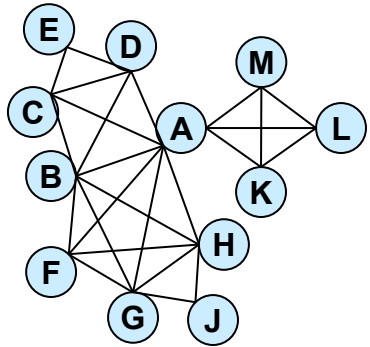


- Inference may require too much memory
- **Condition** on some of the variables

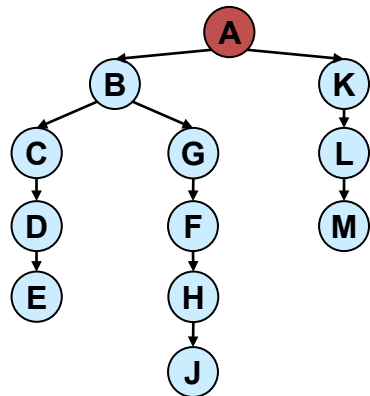




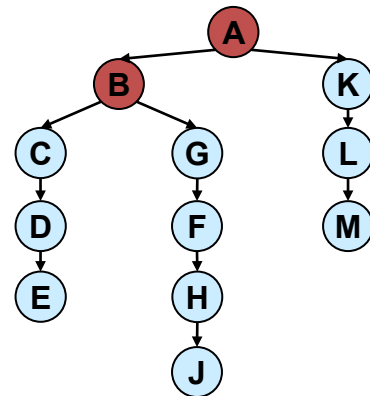
# AND/OR w-cutset



3-cutset



2-cutset



1-cutset

# Summary: Search methods

- **AND/OR search spaces** exploit the structure of the graphical model and create a far more compact search space.
  - **AND/OR Trees** are  $\exp(\text{height})$  of the pseudo-tree and can be traversed in linear memory
  - **AND/OR Graphs** are  $\exp(\text{induced-width})$  of the pseudo-tree and require  $\exp(w)$  memory when traversed.
  - **The pseudo-tree structure** is instrumental in facilitating effective search
- **The MBE heuristic** can guide heuristic search (depth-first, Best-first or hybrid) pruning search further.
- **Tasks:** this schemes are applicable to a large class of tasks:
  - Belief updating, marginal map and Influence diagram search
- **Mixed schemes of Inference and Search** like Cutset schemes facilitate tradeoff between memory and time.